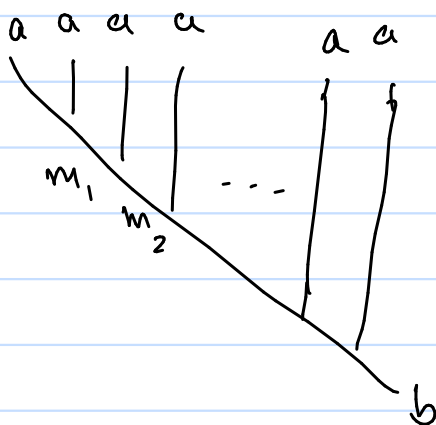


Last time:



Choose a basis & label it

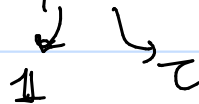


Example: Fib (Fibonacci theory)

$$r = 5$$

$$k = r - 2 = 3$$

We look at  $\{0, 2\} \subseteq \{0, 1, 2, 3\} = \mathcal{L}$



Example: Ising

$$r = 4, k = 2$$

$$\{0, 1, 2\}$$

$$\{\mathbb{1}, \sigma, \psi\}$$

$\psi$  Majorana Fermion

(non abelian anyon)

Setup:  $V_n = \text{Hom}^S(\mathbb{1}, \mathcal{Z}^{\otimes n})$

$n$  Fibonacci anyons  
in a disc with boundary  
labelled by  $\mathbb{0}$

These support qubits!

Thm: The <sup>(Universal)</sup> Quantum Circuit Model can be efficiently simulated on a TQC modelled on Fib.

(We need to be able to, given an arbitrary unitary operator on Hilbert space of size  $2^k$  (say), approximate it)

Recall: gates are braiding

Ingredients:  $B_n$  acts on  $V_n = \text{Hom}(1, \tau^{\otimes n})$  &  $V'_n = \text{Hom}(\tau, \tau^{\otimes n})$

one checks that

$(V_{n+1} \cong V'_n$   
as  $B_n$  reps)

→ it acts irreducibly on both  
(because  $\rho: B_n \rightarrow \text{TL}_n$  is surjective)

(look up  
Braid diagrams)

Notice:  $H(\tau^{\otimes n}, \tau^{\otimes n})$  is an algebra  
it has  $V_n$  &  $V'_n$  as simple modules

→ it acts unitarily

→ [Freedman, Larsen, Wang]

( $\rho$  is the  
above mentioned  
representation)

$$\rho(B_n) \subset U(V_n)$$

$$\rho'(B_n) \subset U(V'_n)$$

but

$$\overline{\rho(B_n)} \cong \text{SU}(V_n), \quad \overline{\rho'(B_n)} \cong \text{SU}(V'_n)$$

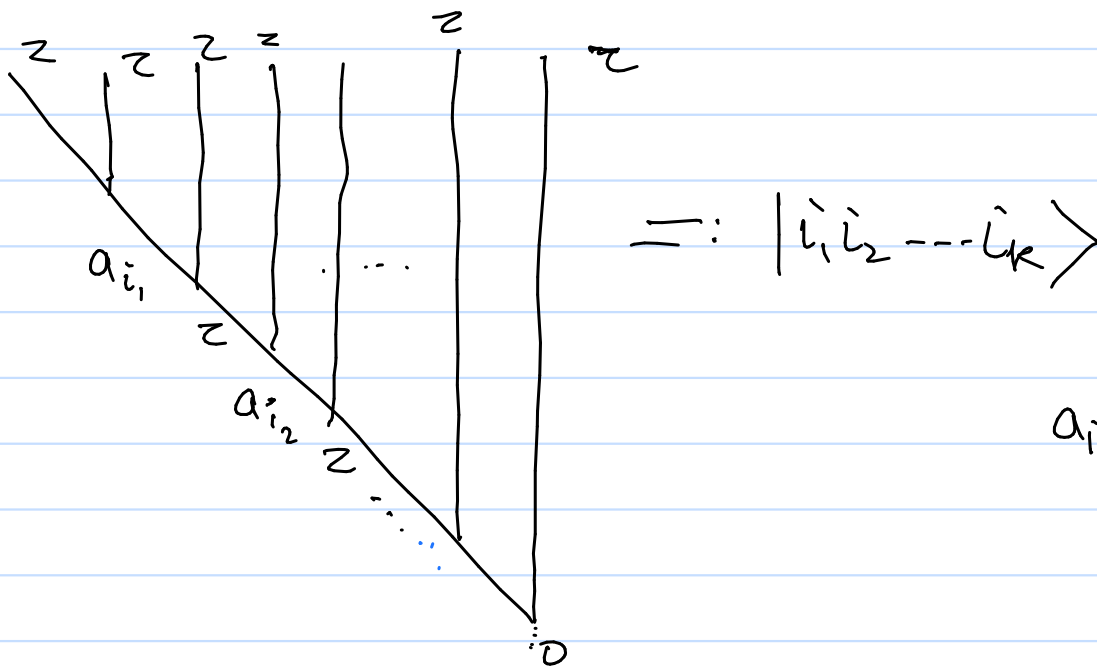
$$\rightarrow \dim(V_n) = \text{Fib}(n)$$

$$\dim(V'_n) = \text{Fib}(n+1)$$

these dimensions grow exponentially

We are now going to realize  $k$ -qubits in  $V_{2k+2}$ .

$$\mathcal{H}_k = (\mathbb{C}^2)^k \hookrightarrow V_{2k+2}$$



$$a_{i_j} \in \{1, \tau\}$$

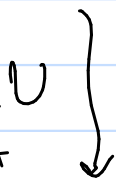
$$a_1 = \tau$$

$$a_0 = 1$$

$$i_j \in \{0, 1\}$$

$$\mathcal{H}_k \hookrightarrow V_{2k+2}$$

we want to simulate this operator



$$\mathcal{H}_k \hookrightarrow V_{2k+2}$$

$$p(\beta) = X$$

want this to approximate  $U$  in a certain way

$\exists$  poly time alg. to find this

choose  $\beta \in B_n$  s.t.

$$(n = 2k+2 \uparrow)$$

$$|X - U \oplus I_{\mathcal{H}_k}| < \epsilon$$

$\in U(V_{2k+2})$

this approximation leads to leakage

this proves the thm

OPEN Q: For given model, which ones can be realized in leakage free way?

Fib TQC is universal.

- Another way of proving the thm is that  $T, H, \text{CNOT}$  are realizable on this model.

- For the Ising model,  $P(B_n)$  turns out to be finite.

So  $\overline{P(B_n)} \not\cong \text{SU}(V_n)$

(What kind of computations can be done on the Ising model?)

Thm: Any TQC can be simulated on  ${}^{(U)}$ QCM efficiently

Pf: Let  $S_n = \mathcal{H} \left( \begin{array}{c} \bullet \quad \dots \quad \bullet \\ a \quad \dots \quad a \end{array} \right)_D$

In order to simulate a TQC, we want to

$\forall S_n \xrightarrow{P(\beta)} S_n \quad \beta \rightarrow \text{braid}$

want to approximate  $P(\beta)$

idea:  $S_n \xrightarrow{P(\beta)} S_n$  by  $U_\beta$

$$\begin{array}{ccc} \downarrow & & \downarrow \\ Y^{\otimes (n-1)} & \xrightarrow{U_\beta} & Y^{\otimes (n-1)} \end{array}$$

$Y = \bigoplus_{(a,b,c) \in \mathbb{Z}^3} \mathcal{H}(P, a, b, c)$

$\rightsquigarrow$  pair of pants with labels  $a, b, c$

Note:  $\dim(Y) \neq 2$


\* But qudit model is polynomially equivalent to qubit model.

$$\text{qudit} = \dim(Y)$$

$$Y^{\otimes n-1} = \left[ \bigoplus_{(a,b,c) \in \mathbb{F}^3} \mathcal{F}(P, a, b, c) \right]^{\otimes (n-1)}$$

distribute tensor product over direct sums

Example:

$$\mathcal{F} \left( \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \right) = \bigoplus_x \left[ \mathcal{F}(P, a, a, x) \otimes \mathcal{F}(P, x, a, 0) \right]$$


This shows that

$$Y^{\otimes (n-1)} \geq S_n$$

i.e.  $Y^{\otimes (n-1)} = S_n \oplus S_n^\perp$

non computational part

Find  $U_B$  s.t.  $U_B|_{S_n} \approx P(B)$

&  $U_B|_{S_n^\perp} \approx I_{S_n^\perp}$

Actual statement: Any 2-dim Topological Quantum functor can be simulated on the QCM.

- $\dim(Y_{\text{Ising}}) = 5$

- What if this theory comes from Hopf algebra? Then we have R-matrix for braiding  
In fact for Ising we can do this

$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & -1 & 1 & 0 \\ -1 & 0 & 0 & 1 \end{bmatrix} \text{ is } \begin{matrix} \text{unitary} \\ \text{solution to YBE} \end{matrix}$$

it can be used to simulate the Ising model.

- There's poly alg. to find  $U_B$ .

## Back to Jones Polynomial:

Thm:  $\left| \frac{J(L, q)}{d^n} \right| \leftarrow \begin{matrix} \text{computation of this is} \\ \text{Bounded error Quantum Poly} \\ \text{B. Q. P.} \end{matrix}$

$$L = \hat{\beta}, \beta \in B_n$$

that for  $q = A^{-4} = e^{-8\pi i / 2\alpha}$   $\alpha = \begin{cases} 2 & \text{if even} \\ 1 & \text{if odd} \end{cases}$

$$\alpha = 5 : q = e^{-4\pi i / 5}$$

---


$$\beta \in B_{2n} \quad \hat{\beta}^{\text{tr}} = \begin{array}{|c|} \hline \alpha \alpha \dots \alpha \\ \hline \beta \\ \hline \alpha \alpha \dots \alpha \\ \hline \end{array}$$

Alexander thm still holds

- Every link is part closure of some braid.

Algorithm:

① initialize:  $U \dots U =: |\text{cup}\rangle$

$$|\text{cup}\rangle \in \text{Hom}(0, 1^{\otimes 2n})$$

②  $P(\beta) |\text{cup}\rangle = \begin{array}{|c|} \hline \text{---} \\ \hline \beta \\ \hline \text{---} \\ \hline \end{array}$

③ Measure against  $\langle \text{cap} | := \overleftarrow{n} \dots \overleftarrow{n}$

(the above is one run of a QC)

Probability of the outcome =  $\overleftarrow{n} \dots \overleftarrow{n}$

is  $\| \langle \overleftarrow{\text{cap}} | P(\beta) |\overrightarrow{\text{cup}} \rangle \|^2$

Output : random variable  $\hat{Z}(\beta) \in \{0, 1\}$   
for each run

(interpretation) output 0 if  $\hat{Z}(\beta) \neq 0$   
(true)

1 else

$$\text{prob}(0) = \| \langle \overleftarrow{\text{cap}} | P(\beta) |\overrightarrow{\text{cup}} \rangle \|^2$$

$$L = \hat{\beta}^{\text{opt}} = \left| \frac{J(L, \beta)}{d^n} \right|^2$$

Let  $Z(\beta) =$  average of  $\hat{Z}(\beta)$  for some  
in  $\delta \leftarrow$  poly nb. of tries

then  $Z(\beta) \in [0, 1]$

Then:  $\text{Prob} \left( \left| \frac{J(L, \alpha)}{d^n} - z(\beta) \right| < \delta \right) \geq \frac{3}{4}$

$q = e^{\pi i/l}$

$l = 1, 2, 3, 4, 6 \xrightarrow{\mathbb{Z}_2} \text{Ising}$   
 $\xrightarrow{\text{metaplectic cat. obj of dim } \{1, 2, 5, 3\}}$

there is classical easy algorithm

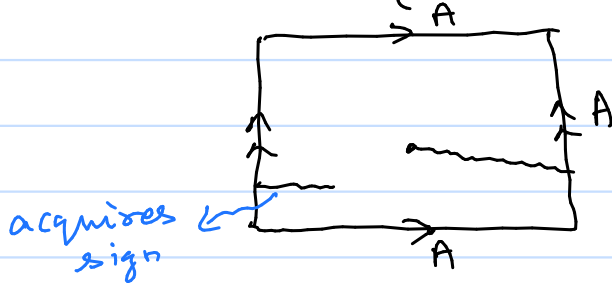
( $l=3$ ) Ising theory  $\rightarrow$  gives link invariant called Arf invariant

( $l=2$ )  $\rightarrow$  counts rank of homology of some cover

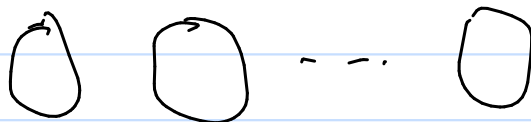
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Beyond 2D bosonic theories

• Fermionic (non-local)



• 3D: pointlike particle: boson/fermion  
 can consider looplike particles



(3+1) TQFTs



Invariant associated with the double of a finite group  $D(G)$

→ it counts no. of homomorphisms

$$\pi_1(\tilde{K}) \rightarrow G / \sim$$

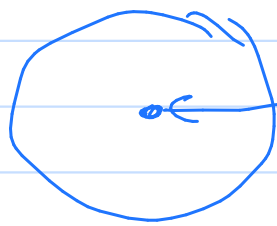
↑  
Knot complement

(eg  $\tilde{K} = B_3$  for  $K$  trefoil)

It's open if there is a poly-time alg. for counting such homo.

One expects to have such an alg.

## Gapped phases of matter



instead of just simple object, we put an algebra in the category

→ Iris Cong