

MATH 466/566
SPRING 2024

CHELSEA WALTON
RICE U.

LAST TIME

- CATEGORY ISOMORPHISM
- CATEGORY EQUIVALENCE
- MORITA EQUIVALENCE

LECTURE #10

TOPICS:

- I. ADJUNCTION (§2.5)
- II. UNIVERSALITY REVISITED (§2.6.1)
- III. YONEDA'S LEMMA (§2.6.2)

I. ADJUNCTION

$\mathcal{C} = \mathcal{D}$
EQUALITY
OF CATEGORIES

NOTION OF SAMENESS
FOR CATEGORIES

WEAKEN

WEAKEN

$\mathcal{C} \cong \mathcal{D}$
ISOMORPHISM
OF CATEGORIES

$\mathcal{C} \simeq \mathcal{D}$
EQUIVALENCE
OF CATEGORIES

I. ADJUNCTION

$\mathcal{C} = \mathcal{D}$
EQUALITY
OF CATEGORIES

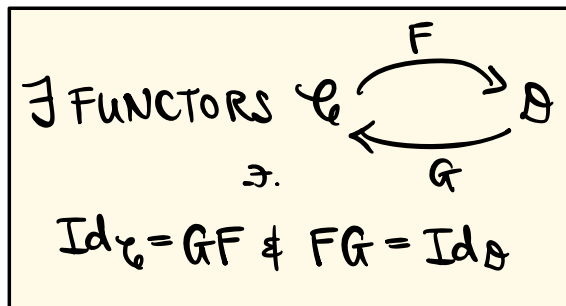
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INVOLVING
EQUALITIES
OF FUNCTORS

I. ADJUNCTION

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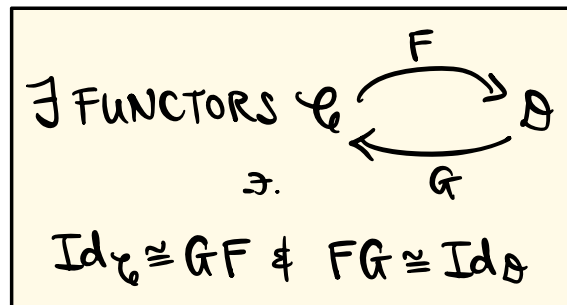
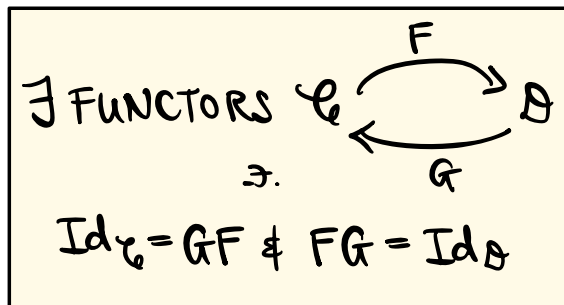
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INVOLVING
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I. ADJUNCTION

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EQUALITY
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NOTION OF SAMENESS
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RICH THEORY OF
FUNCTORS

WEAKEN

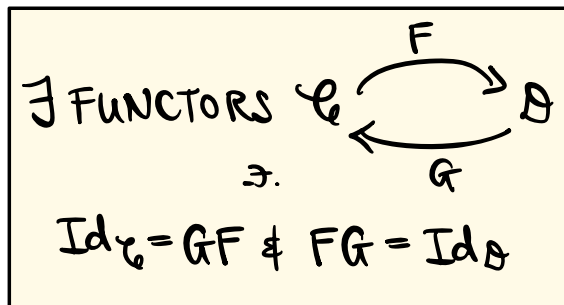
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ISOMORPHISM
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WEAKEN

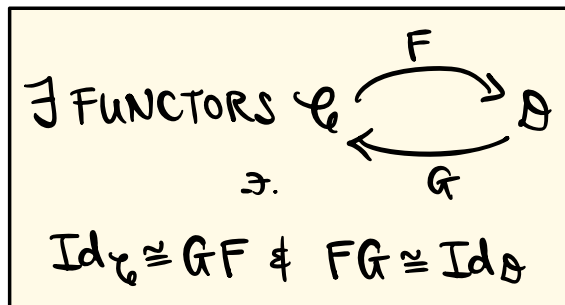
$\mathcal{C} \simeq \mathcal{D}$
EQUIVALENCE
OF CATEGORIES

WEAKEN

NOW:
ADJUNCTION



INVOLVING
EQUALITIES
OF FUNCTORS



INVOLVING
ISOMORPHISMS
OF FUNCTORS

I. ADJUNCTION

$\mathcal{C} = \mathcal{D}$
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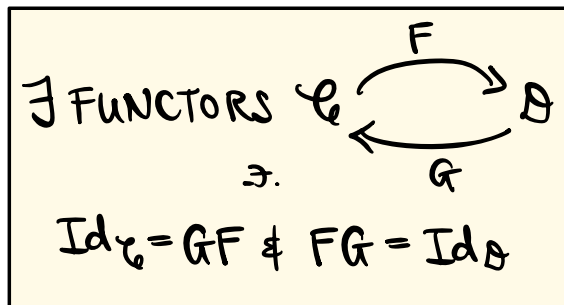
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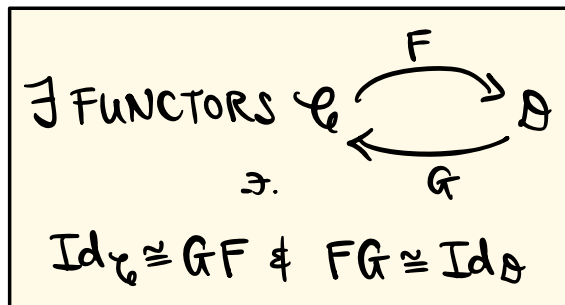
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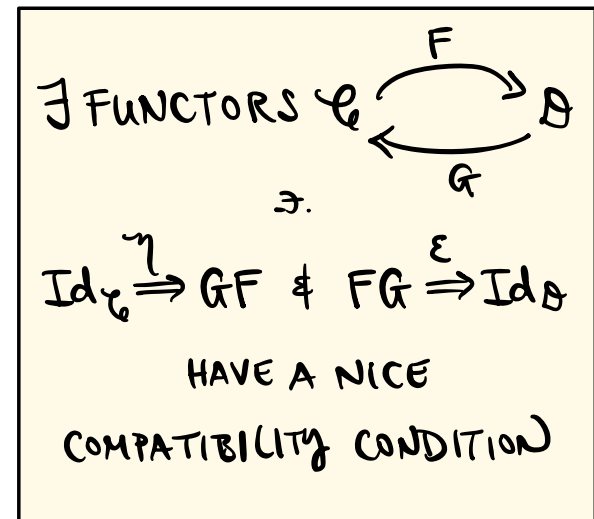
NOW:
ADJUNCTION



INVOLVING
EQUALITIES
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INVOLVING
ISOMORPHISMS
OF FUNCTORS



INVOLVES CERTAIN TRANSFORMS
OF FUNCTORS

I. ADJUNCTION

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN ADJUNCTION

IF \exists NAT'L TRANSFORMATIONS

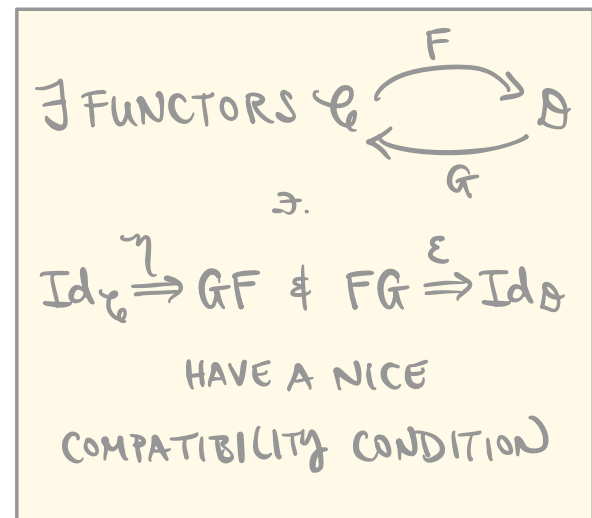
$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \quad \& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

\exists .

RICH THEORY OF
FUNCTORS

NOW:

ADJUNCTION



INVOLVES CERTAIN TRANSFORMS
OF FUNCTORS

I. ADJUNCTION

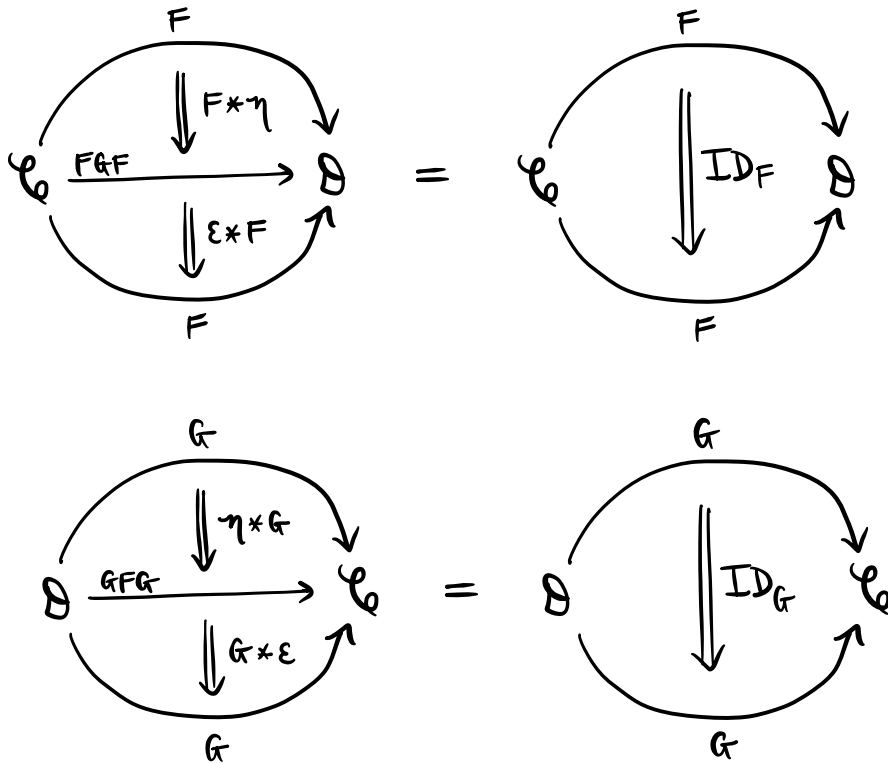
FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN **ADJUNCTION**

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$$\eta: Id_{\mathcal{C}} \Rightarrow GF \quad \& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

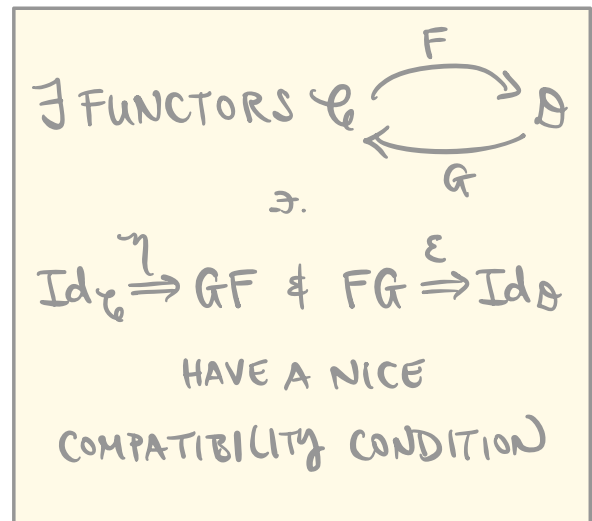
$\exists.$



RICH THEORY OF
FUNCTORS

NOW:

ADJUNCTION



INVOLVES CERTAIN TRANSFORMS
OF FUNCTORS

I. ADJUNCTION

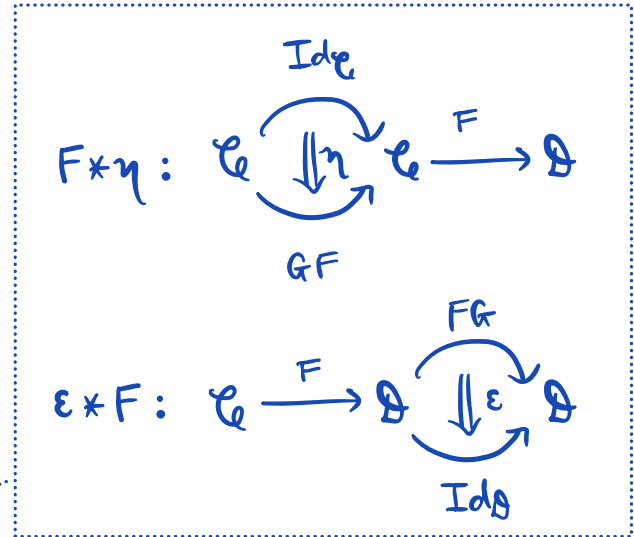
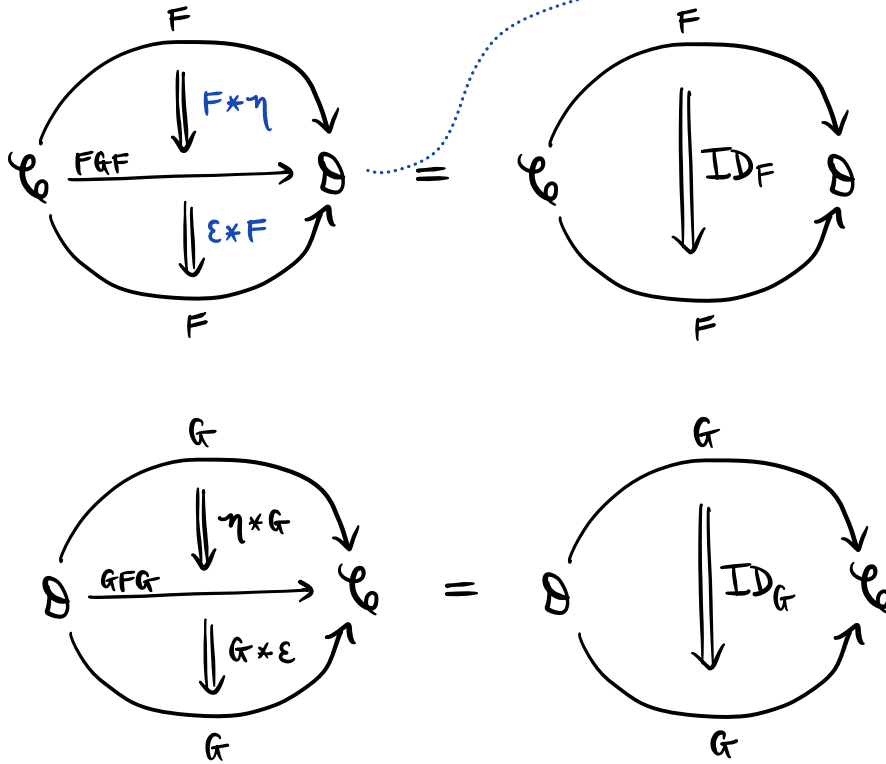
FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN **ADJUNCTION**

IF \exists NAT'L TRANSFORMATIONS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \quad \& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

\Rightarrow



I. ADJUNCTION

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FORM AN **ADJUNCTION**

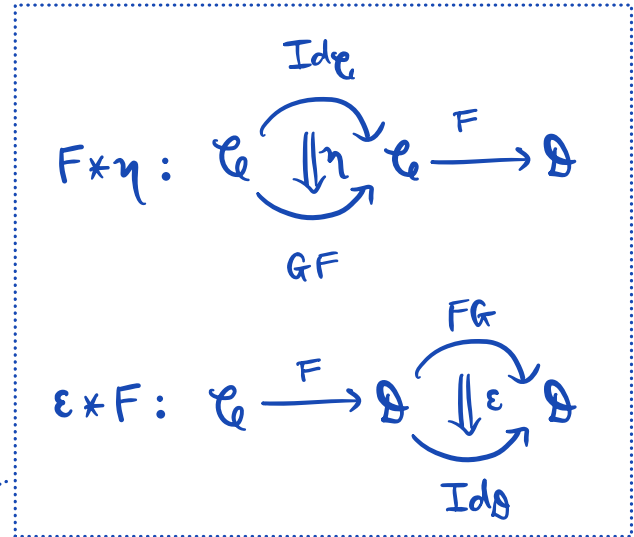
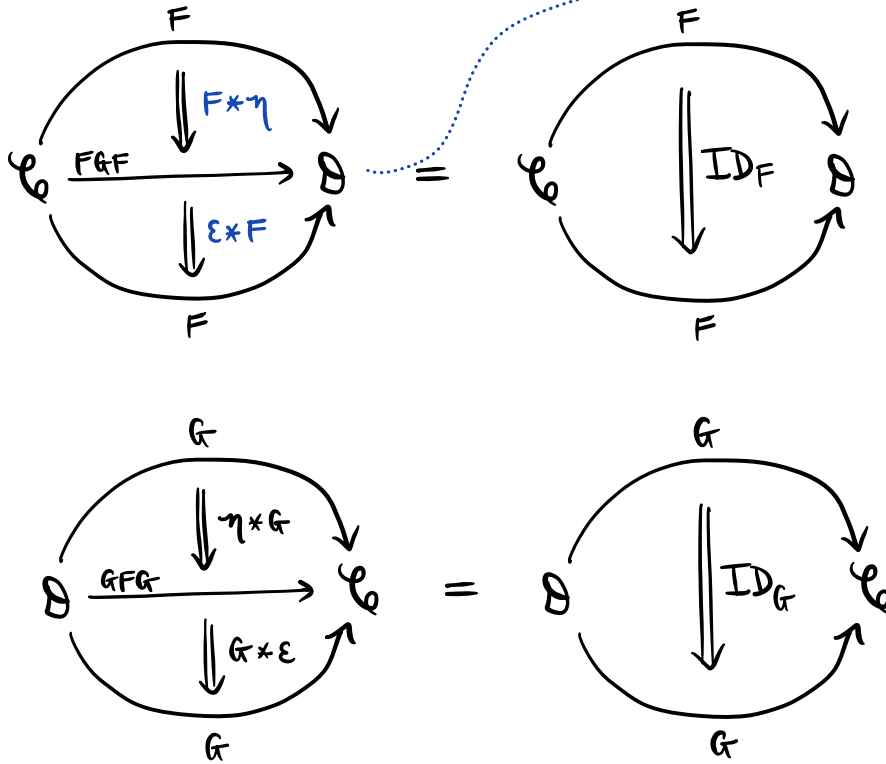
IF \exists NAT'L TRANSFORMATIONS

$\eta: Id_{\mathcal{C}} \Rightarrow GF$ & $\epsilon: FG \Rightarrow Id_{\mathcal{D}}$

ADJUNCTION UNIT

ADJUNCTION COUNIT

\Rightarrow



I. ADJUNCTION

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN **ADJUNCTION**

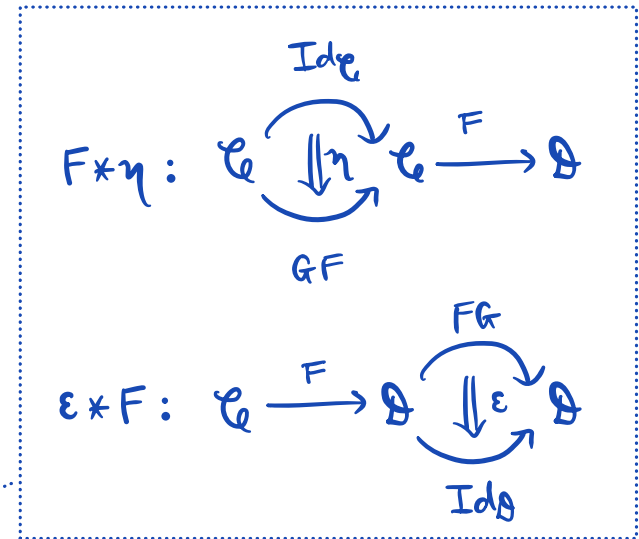
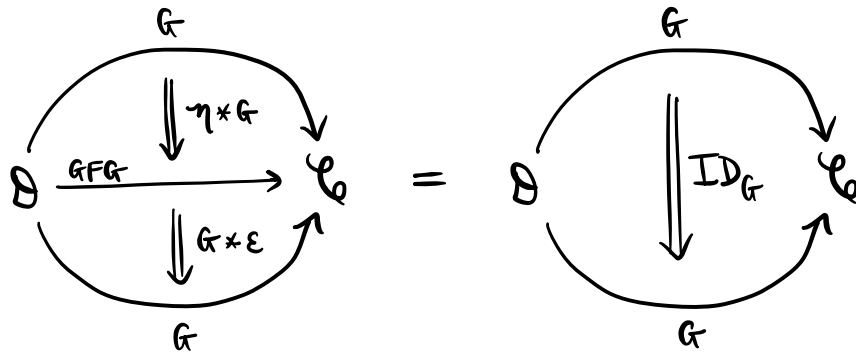
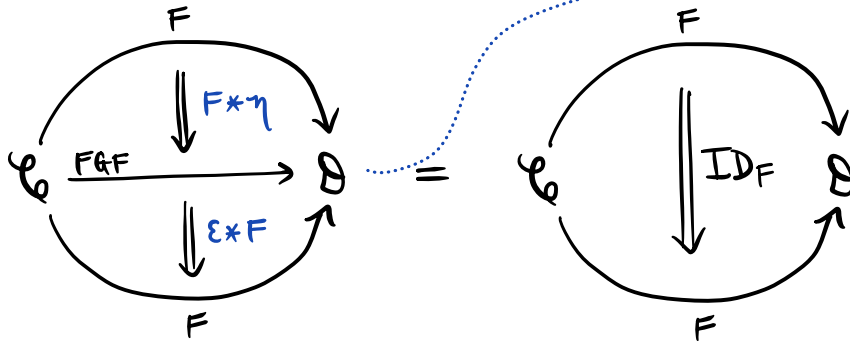
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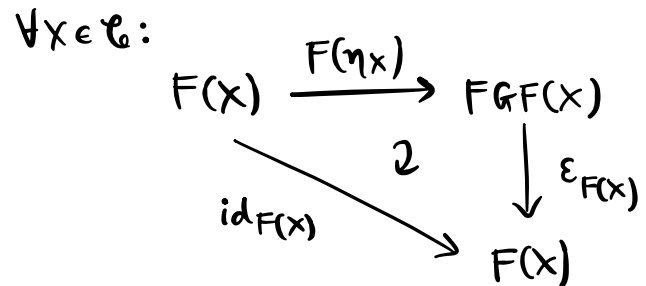
ADJUNCTION UNIT

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ON COMPONENTS



I. ADJUNCTION

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN **ADJUNCTION**

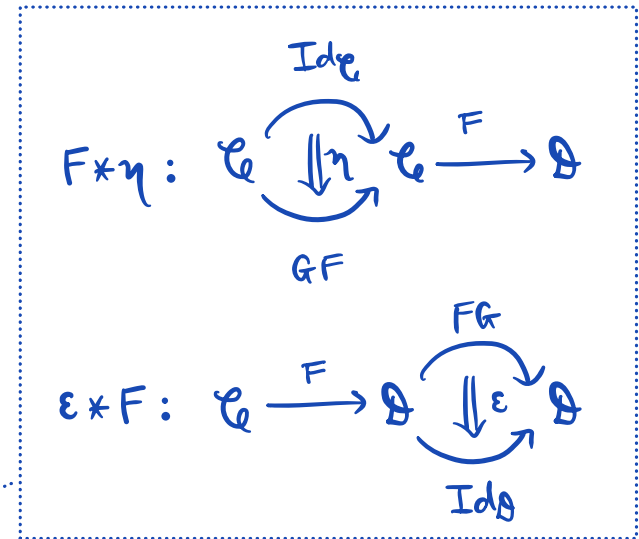
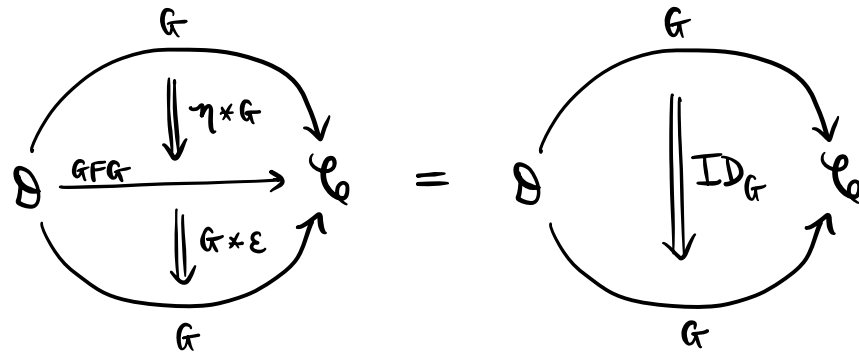
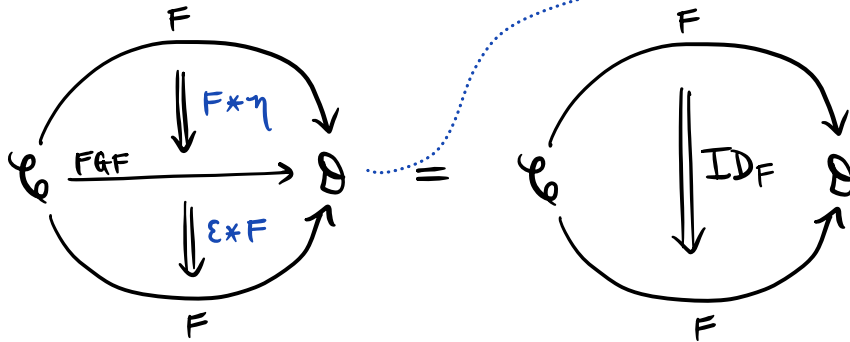
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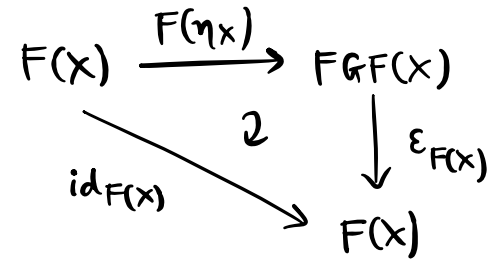
ADJUNCTION COUNIT

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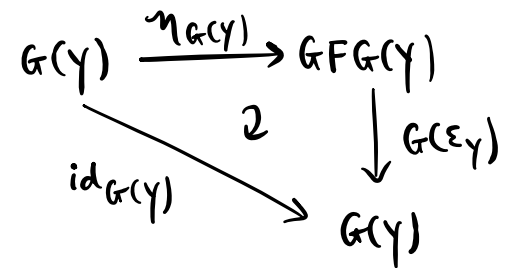


ON COMPONENTS

$\forall x \in \mathcal{C}$:



$\forall y \in \mathcal{D}$:



I. ADJUNCTION

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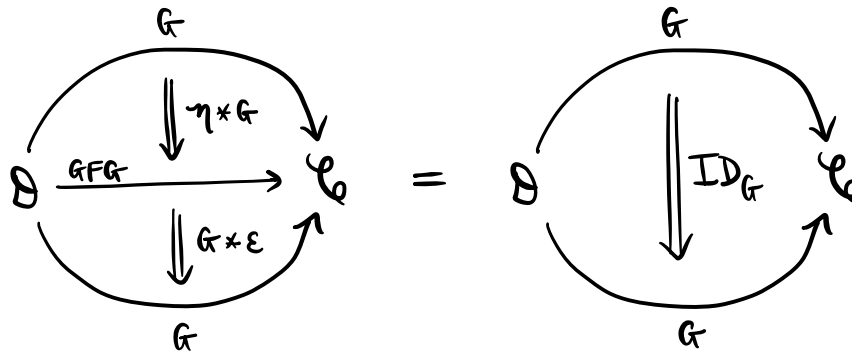
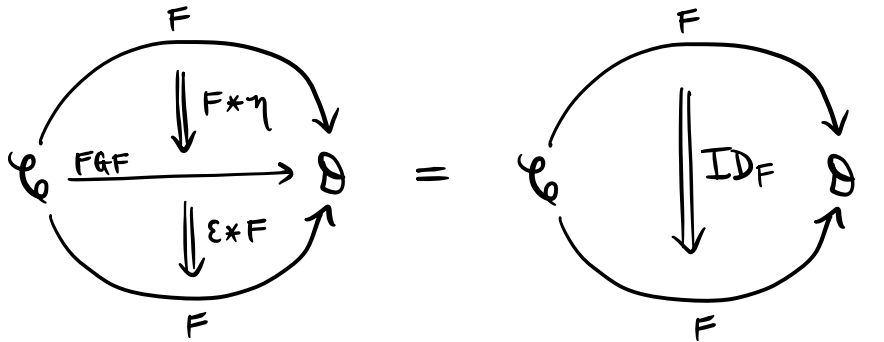
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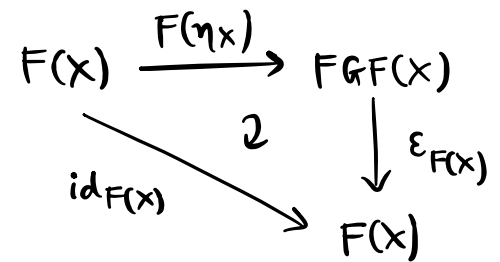
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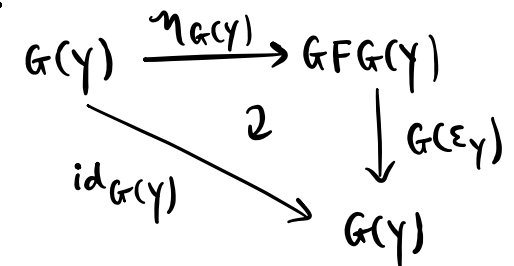
TRIANGLE IDENTITIES

ON COMPONENTS

$\forall x \in \mathcal{C}$:



$\forall y \in \mathcal{D}$:



I. ADJUNCTION

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FORM AN **ADJUNCTION**

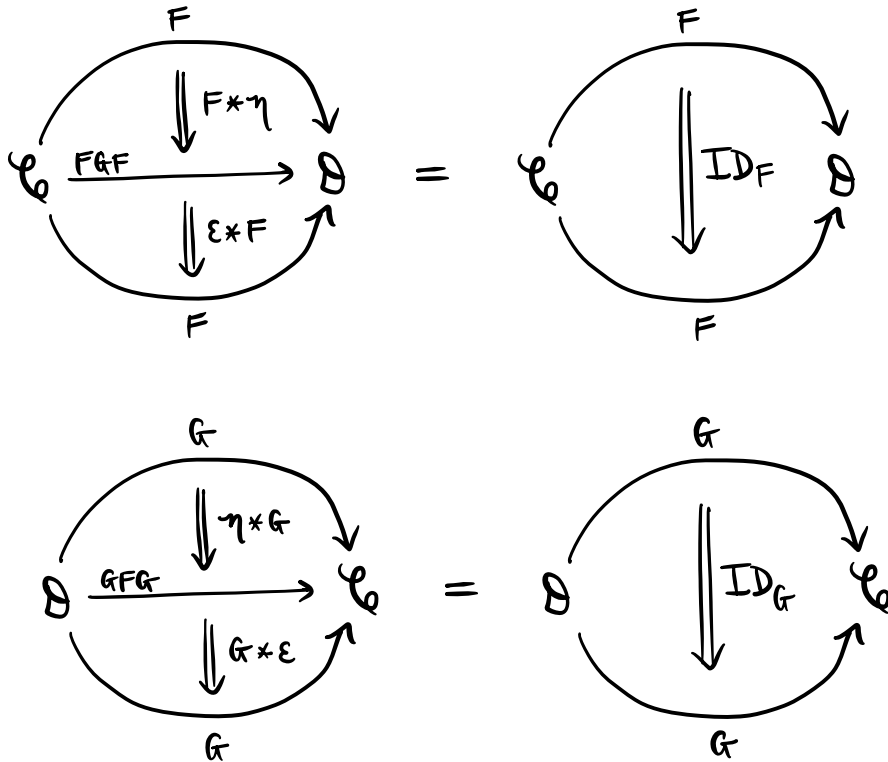
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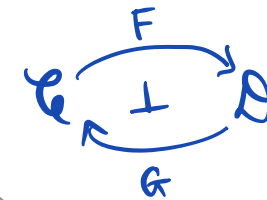
ADJUNCTION COUNIT

\exists .



WRITE $F \dashv G$

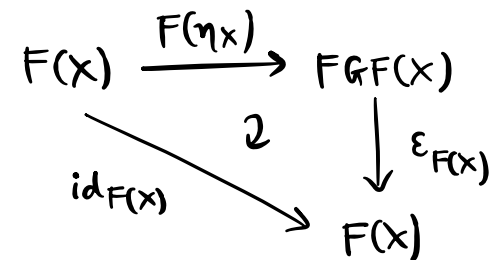
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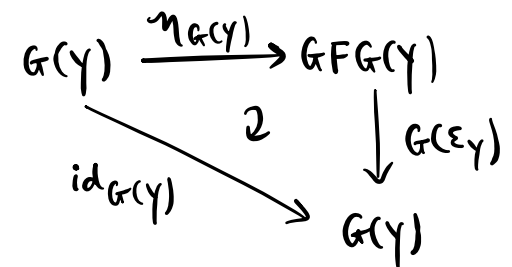
TRIANGLE IDENTITIES

ON COMPONENTS

$\forall x \in \mathcal{C}$:



$\forall y \in \mathcal{D}$:



I. ADJUNCTION \Leftrightarrow EQUIVALENCE (loosely)

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN ADJUNCTION

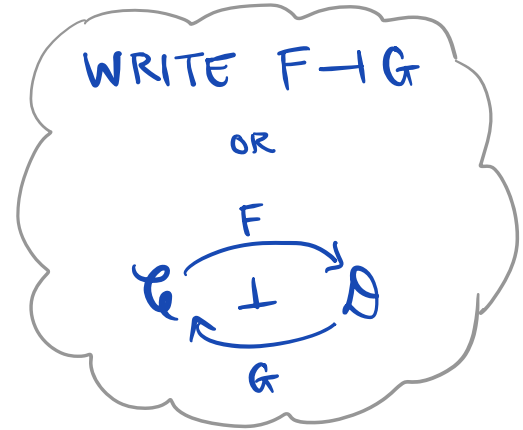
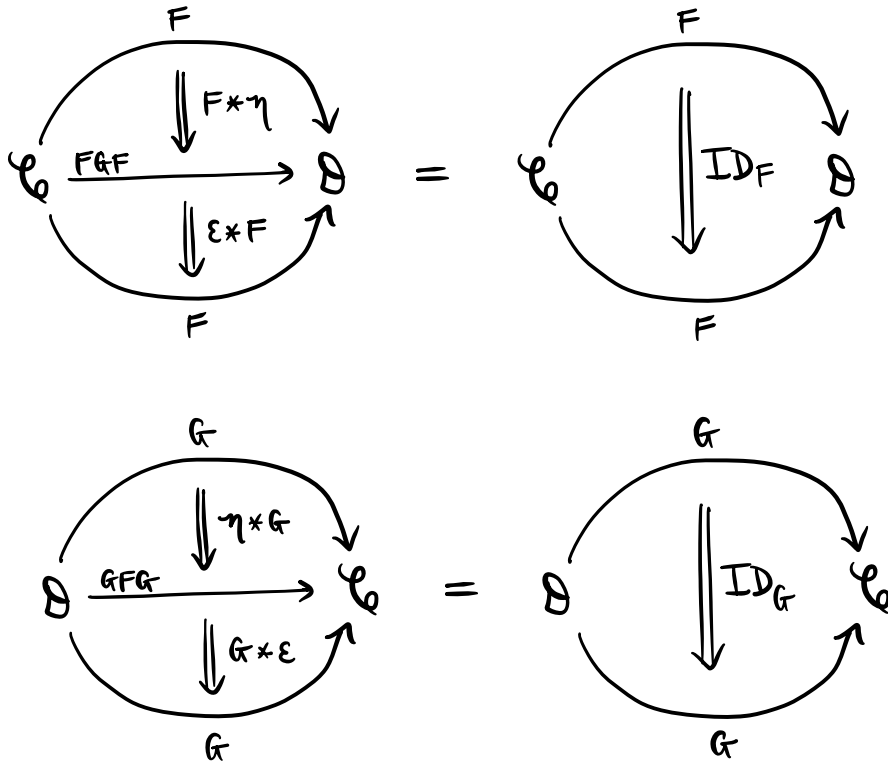
IF \exists NAT'L TRANSFORMATIONS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF \quad \& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

ADJUNCTION UNIT

ADJUNCTION COUNIT

\exists .



TRIANGLE IDENTITIES

ON COMPONENTS

$\forall x \in \mathcal{C}$:

$$\begin{array}{ccc}
 F(x) & \xrightarrow{F(\eta_x)} & FG F(x) \\
 \searrow id_{F(x)} & & \downarrow \varepsilon_{F(x)} \\
 & & F(x)
 \end{array}$$

$\forall y \in \mathcal{D}$:

$$\begin{array}{ccc}
 G(y) & \xrightarrow{\eta_{G(y)}} & GFG(y) \\
 \searrow id_{G(y)} & & \downarrow G(\varepsilon_y) \\
 & & G(y)
 \end{array}$$

I. ADJUNCTION \Leftarrow EQUIVALENCE (loosely)

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN ADJUNCTION

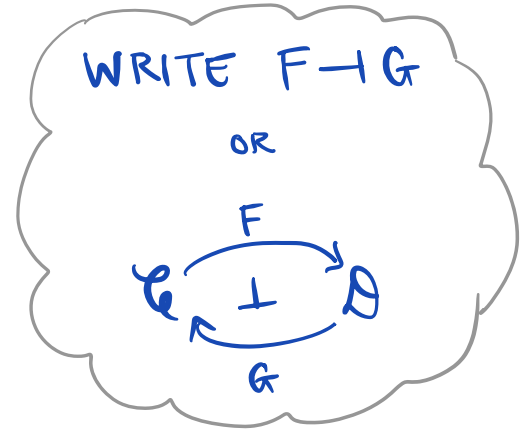
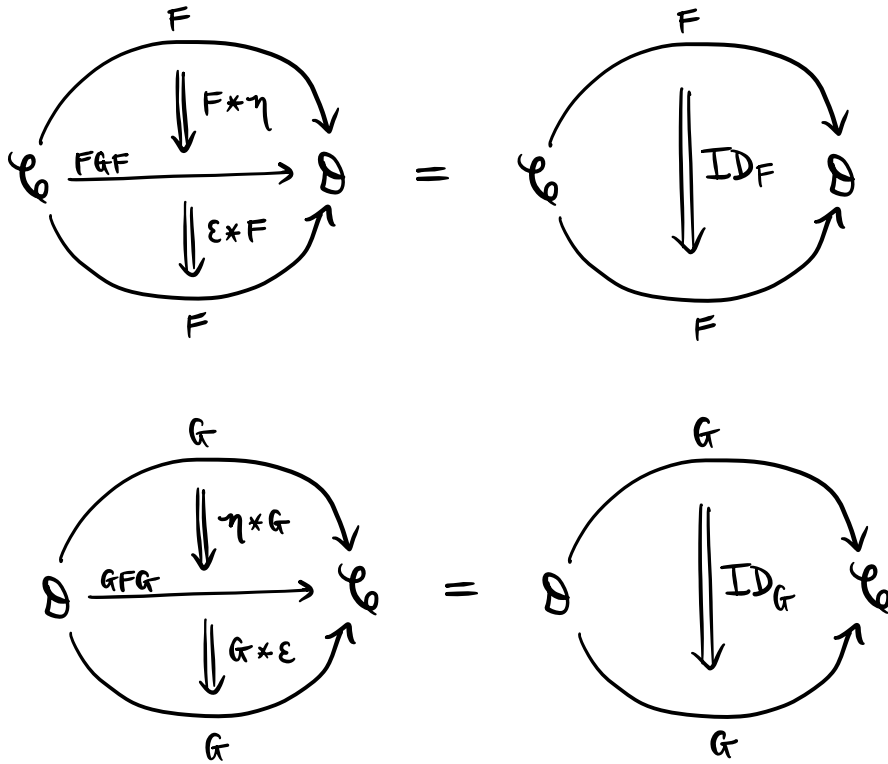
IF \exists NAT'L TRANSFORMATIONS

$$\eta: \text{Id}_{\mathcal{C}} \xrightarrow{\sim} GF \quad \& \quad \varepsilon: FG \xrightarrow{\sim} \text{Id}_{\mathcal{D}}$$

ADJUNCTION UNIT

ADJUNCTION COUNIT

\exists .



TRIANGLE IDENTITIES

ON COMPONENTS

$\forall x \in \mathcal{C}$:

$$\begin{array}{ccc}
 F(x) & \xrightarrow[\sim]{F(\eta_x)} & FG F(x) \\
 \searrow \text{id}_{F(x)} & & \downarrow \varepsilon_{F(x)} \\
 & & F(x)
 \end{array}$$

$\forall y \in \mathcal{D}$:

$$\begin{array}{ccc}
 G(y) & \xrightarrow[\sim]{\eta_{G(y)}} & GFG(y) \\
 \searrow \text{id}_{G(y)} & & \downarrow G(\varepsilon_y) \\
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 \end{array}$$

I. ADJUNCTION \Leftrightarrow EQUIVALENCE (loosely)

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN ADJUNCTION

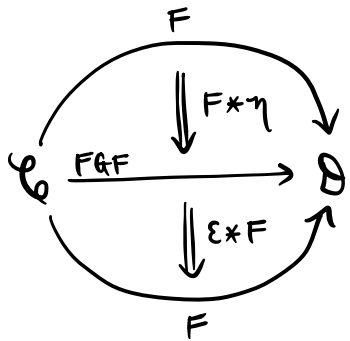
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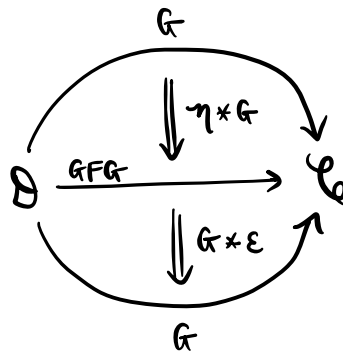
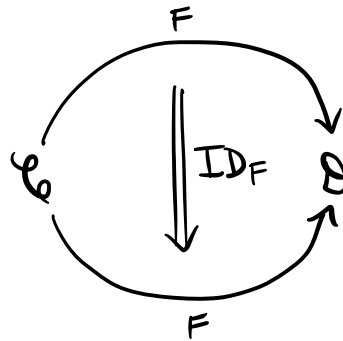
ADJUNCTION UNIT

ADJUNCTION COUNIT

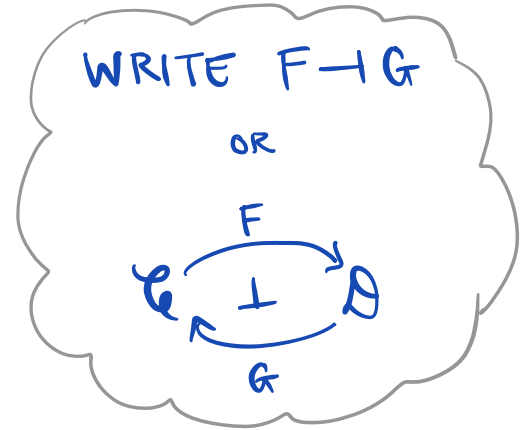
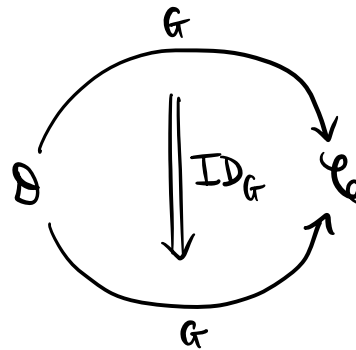
\exists .



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TRIANGLE IDENTITIES

ON COMPONENTS

$\forall x \in \mathcal{C}$:

$$F(x) \xrightarrow{F(\eta_x)} FG(x) \xrightarrow{\varepsilon_{FG(x)}} F(x)$$

id_{F(x)} ↘

$\forall y \in \mathcal{D}$:

$$G(y) \xrightarrow{G(\eta_y)} GF(y) \xrightarrow{\varepsilon_{GF(y)}} G(y)$$

id_{G(y)} ↘

I. ADJUNCTION \Leftrightarrow EQUIVALENCE (WILL REVISIT...)

FUNCTORS $F: \mathcal{C} \rightarrow \mathcal{D}$ & $G: \mathcal{D} \rightarrow \mathcal{C}$

FORM AN ADJUNCTION

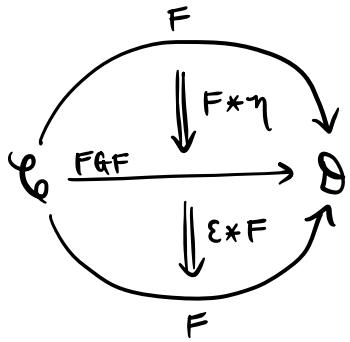
IF \exists NAT'L TRANSFORMATIONS

$$\eta: \text{Id}_{\mathcal{C}} \xrightarrow{\sim} GF \quad \& \quad \varepsilon: FG \xrightarrow{\sim} \text{Id}_{\mathcal{D}}$$

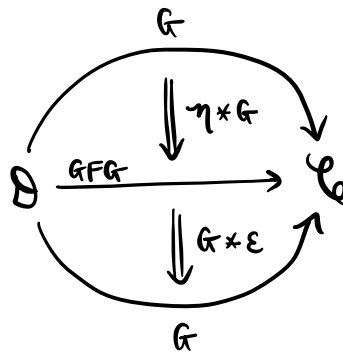
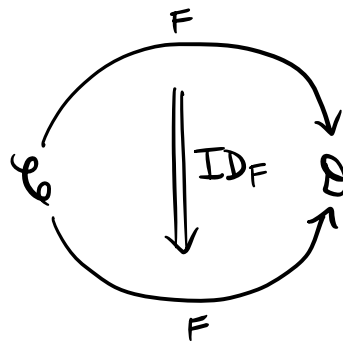
ADJUNCTION UNIT

ADJUNCTION COUNIT

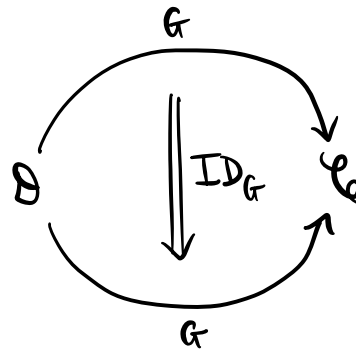
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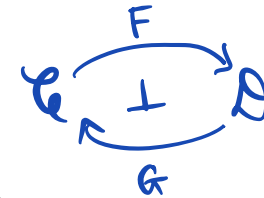


=



WRITE $F \dashv G$

OR



TRIANGLE IDENTITIES

ON COMPONENTS

$\forall x \in \mathcal{C}$:

$$F(x) \xrightarrow{F(\eta_x)} FG(x) \xrightarrow{\varepsilon_{FG(x)}} F(x)$$

$\downarrow \text{id}_{F(x)}$

$\forall y \in \mathcal{D}$:

$$G(y) \xrightarrow{G(\varepsilon_y)} GF(y) \xrightarrow{\eta_{GF(y)}} G(y)$$

$\downarrow \text{id}_{G(y)}$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(x) & \xrightarrow{F(\eta_x)} & FG F(x) \\ & \searrow \text{id}_{F(x)} & \downarrow \varepsilon_{F(x)} \\ & & F(x) \end{array} \quad \forall x \in \mathcal{C}$$

$$\begin{array}{ccc} G(y) & \xrightarrow{\eta_{G(y)}} & GFG(y) \\ & \searrow \text{id}_{G(y)} & \downarrow G(\varepsilon_y) \\ & & G(y) \end{array} \quad \forall y \in \mathcal{D}$$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

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$$\begin{array}{ccc} G(y) & \xrightarrow{\eta_{G(y)}} & GFG(y) \\ & \searrow \text{id}_{G(y)} & \downarrow G(\varepsilon_y) \\ & & G(y) \quad \forall y \in \mathcal{D} \end{array}$$

$F \equiv$ LEFT ADJOINT OF G

$G \equiv$ RIGHT ADJOINT OF F

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

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$$\begin{array}{ccc} G(y) & \xrightarrow{\eta_{G(y)}} & GFG(y) \\ & \searrow \text{id}_{G(y)} & \downarrow G(\varepsilon_y) \\ & & G(y) \quad \forall y \in \mathcal{D} \end{array}$$

$F \equiv$ LEFT ADJOINT OF G

$G \equiv$ RIGHT ADJOINT OF F

NEED NOT EXIST,

UNIQUE UP TO

NATURAL \cong

IF EXISTS

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

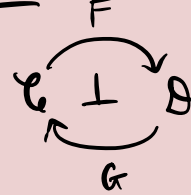
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

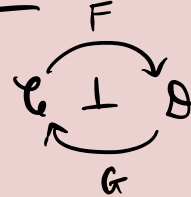
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

$\forall f: X' \rightarrow X''$ WITH Y FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(\bullet), Y) & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(f), Y)} & \text{Hom}_{\mathcal{D}}(F(\bullet), Y) \\ \downarrow \exists_{\bullet, Y} & \cong & \downarrow \exists_{\bullet, Y} \\ \text{Hom}_{\mathcal{C}}(\bullet, G(Y)) & \xrightarrow{\text{Hom}_{\mathcal{C}}(f, G(Y))} & \text{Hom}_{\mathcal{C}}(\bullet, G(Y)) \end{array}$$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

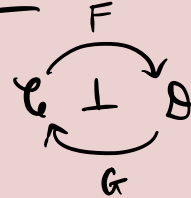
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists \zeta_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

$\forall f: X' \rightarrow X''$ WITH Y FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X''), Y) & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(f), Y)} & \text{Hom}_{\mathcal{D}}(F(X'), Y) \\ \zeta_{X'', Y} \downarrow & \cong & \downarrow \zeta_{X', Y} \\ \text{Hom}_{\mathcal{C}}(X'', G(Y)) & \xrightarrow{\text{Hom}_{\mathcal{C}}(f, G(Y))} & \text{Hom}_{\mathcal{C}}(X', G(Y)) \end{array}$$

RECALL $\text{Hom}_{\mathcal{C}}(-, Z)$ IS CONTRAVARIANT

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

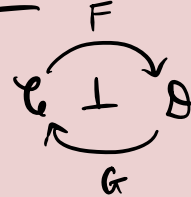
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists$ BIJECTION

$$\zeta_{X,Y}: \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

$\forall f: X' \rightarrow X''$ WITH Y FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X''), Y) & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(f), Y)} & \text{Hom}_{\mathcal{D}}(F(X'), Y) \\ \zeta_{X'', Y} \downarrow & \cong & \downarrow \zeta_{X', Y} \\ \text{Hom}_{\mathcal{C}}(X'', G(Y)) & \xrightarrow{\text{Hom}_{\mathcal{C}}(f, G(Y))} & \text{Hom}_{\mathcal{C}}(X', G(Y)) \end{array}$$

RECALL $\text{Hom}_{\mathcal{C}}(-, Z)$ IS CONTRAVARIANT

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

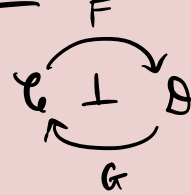
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists$ BIJECTION

$$\mathcal{S}_{X,Y}: \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

$\forall f: X' \rightarrow X''$ WITH Y FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X''), Y) & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(f), Y)} & \text{Hom}_{\mathcal{D}}(F(X'), Y) \\ \mathcal{S}_{X'', Y} \downarrow & \cong & \downarrow \mathcal{S}_{X', Y} \\ \text{Hom}_{\mathcal{C}}(X'', G(Y)) & \xrightarrow{\text{Hom}_{\mathcal{C}}(f, G(Y))} & \text{Hom}_{\mathcal{C}}(X', G(Y)) \end{array}$$

RECALL $\text{Hom}_{\mathcal{C}}(z, -)$ IS COVARIANT

$\forall g: Y' \rightarrow Y''$ WITH X FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X), Y') & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(X), g)} & \text{Hom}_{\mathcal{D}}(F(X), Y'') \\ \mathcal{S}_{X, Y'} \downarrow & \cong & \downarrow \mathcal{S}_{X, Y''} \\ \text{Hom}_{\mathcal{C}}(X, G(Y')) & \xrightarrow{\text{Hom}_{\mathcal{C}}(X, G(g))} & \text{Hom}_{\mathcal{C}}(X, G(Y'')) \end{array}$$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

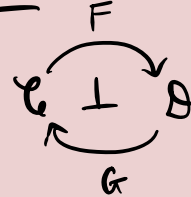
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists$ BIJECTION

$$\mathcal{S}_{X,Y}: \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

$\forall f: X' \rightarrow X''$ WITH Y FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X''), Y) & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(f), Y)} & \text{Hom}_{\mathcal{D}}(F(X'), Y) \\ \mathcal{S}_{X'', Y} \downarrow & \cong & \downarrow \mathcal{S}_{X', Y} \\ \text{Hom}_{\mathcal{C}}(X'', G(Y)) & \xrightarrow{\text{Hom}_{\mathcal{C}}(f, G(Y))} & \text{Hom}_{\mathcal{C}}(X', G(Y)) \end{array}$$

$\forall g: Y' \rightarrow Y''$ WITH X FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X), Y') & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(X), g)} & \text{Hom}_{\mathcal{D}}(F(X), Y'') \\ \mathcal{S}_{X, Y'} \downarrow & \cong & \downarrow \mathcal{S}_{X, Y''} \\ \text{Hom}_{\mathcal{C}}(X, G(Y')) & \xrightarrow{\text{Hom}_{\mathcal{C}}(X, G(g))} & \text{Hom}_{\mathcal{C}}(X, G(Y'')) \end{array}$$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

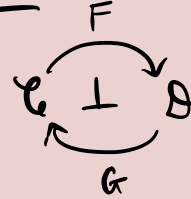
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow id_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow id_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



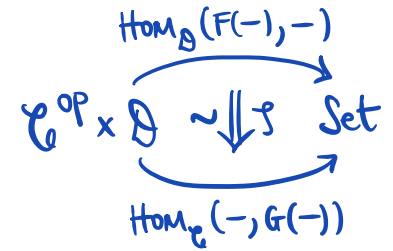
$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\Leftrightarrow \exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

$\forall f: X' \rightarrow X''$ WITH Y FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X''), Y) & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(f), Y)} & \text{Hom}_{\mathcal{D}}(F(X'), Y) \\ \downarrow \exists_{X'', Y} & \cong & \downarrow \exists_{X', Y} \\ \text{Hom}_{\mathcal{C}}(X'', G(Y)) & \xrightarrow{\text{Hom}_{\mathcal{C}}(f, G(Y))} & \text{Hom}_{\mathcal{C}}(X', G(Y)) \end{array}$$



$\forall g: Y' \rightarrow Y''$ WITH X FIXED GET:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(F(X), Y') & \xrightarrow{\text{Hom}_{\mathcal{D}}(F(X), g)} & \text{Hom}_{\mathcal{D}}(F(X), Y'') \\ \downarrow \exists_{X, Y'} & \cong & \downarrow \exists_{X, Y''} \\ \text{Hom}_{\mathcal{C}}(X, G(Y')) & \xrightarrow{\text{Hom}_{\mathcal{C}}(X, G(g))} & \text{Hom}_{\mathcal{C}}(X, G(Y'')) \end{array}$$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

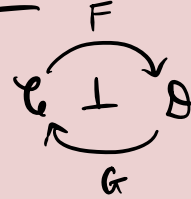
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(x) & \xrightarrow{F(\eta_x)} & FGf(x) \\ & \searrow id_{F(x)} & \downarrow \varepsilon_{F(x)} \\ & & F(x) \end{array} \quad \forall x \in \mathcal{C}$$

$$\begin{array}{ccc} G(y) & \xrightarrow{\eta_{G(y)}} & GFG(y) \\ & \searrow id_{G(y)} & \downarrow G(\varepsilon_y) \\ & & G(y) \end{array} \quad \forall y \in \mathcal{D}$$

PROP:



$$\forall x \in \mathcal{C} \quad \& \quad y \in \mathcal{D} : \exists \text{ BIJECTION}$$

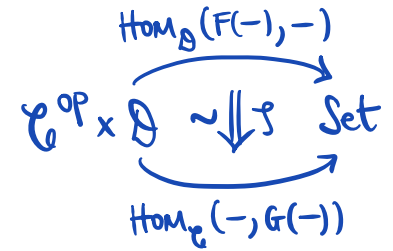
$$\int_{x,y} : \text{Hom}_{\mathcal{D}}(F(x), y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \quad \& \quad y$

PF/ (\Rightarrow) GIVEN $(F \dashv G, \eta, \varepsilon)$,

DEFINE:

$$\text{Hom}_{\mathcal{D}}(F(x), y) \xrightarrow{\int_{x,y}} \text{Hom}_{\mathcal{C}}(x, G(y))$$



I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

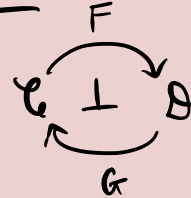
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\int_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

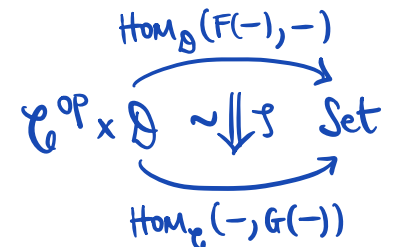
PF/ (\Rightarrow) GIVEN $(F \dashv G, \eta, \varepsilon)$,

DEFINE:

$$\text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\int_{X,Y}} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

$$[h: F(X) \rightarrow Y] \mapsto$$

$$\leftarrow [X \xrightarrow{d} G(Y)]$$



I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

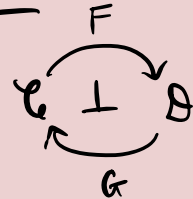
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow id_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \quad \forall X \in \mathcal{C} \end{array}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow id_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \quad \forall Y \in \mathcal{D} \end{array}$$

PROP:



$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists$ BIJECTION

$$\int_{X,Y} : Hom_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} Hom_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

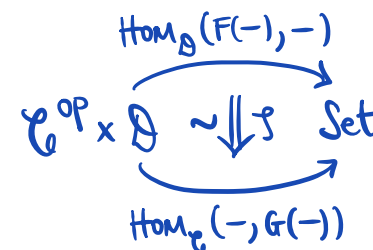
PF/ (\Rightarrow) GIVEN $(F \dashv G, \eta, \varepsilon)$,

DEFINE:

$$Hom_{\mathcal{D}}(F(X), Y) \xrightarrow{\int_{X,Y}} Hom_{\mathcal{C}}(X, G(Y))$$

$$\left[h: F(X) \rightarrow Y \right] \longmapsto \left[X \xrightarrow{\eta_X} GF(X) \xrightarrow{G(h)} G(Y) \right]$$

$$\left[F(X) \xrightarrow{F(\alpha)} FG(Y) \xrightarrow{\varepsilon_Y} Y \right] \longleftarrow \left[X \xrightarrow{\alpha} G(Y) \right]$$



I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

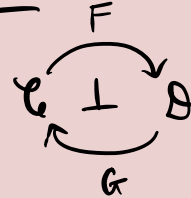
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow id_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{\eta_{G(Y)}} & GFG(Y) \\ & \searrow id_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:



$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\int_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Rightarrow) GIVEN $(F \dashv G, \eta, \varepsilon)$,

DEFINE:

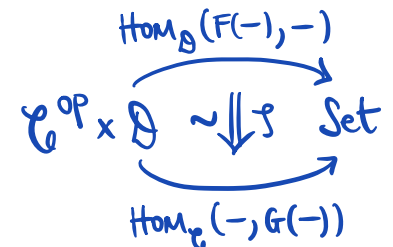
$$\text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\int_{X,Y}} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

$$\left[h: F(X) \rightarrow Y \right] \longmapsto \left[X \xrightarrow{\eta_X} GF(X) \xrightarrow{G(h)} G(Y) \right]$$

$$\left[F(X) \xrightarrow{F(h)} FG(Y) \xrightarrow{\varepsilon_Y} Y \right] \longleftarrow \left[X \xrightarrow{h} G(Y) \right]$$

CHECK ASSIGNMENTS ARE MUTUALLY INVERSE
& NATURAL IN X AND Y .

EXER. 2.40



I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

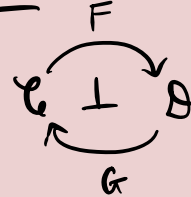
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc}
 F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\
 \searrow \text{id}_{F(X)} & \cong & \downarrow \varepsilon_{F(X)} \\
 & & F(X) \quad \forall X \in \mathcal{C}
 \end{array}$$

$$\begin{array}{ccc}
 G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\
 \searrow \text{id}_{G(Y)} & \cong & \downarrow G(\varepsilon_Y) \\
 & & G(Y) \quad \forall Y \in \mathcal{D}
 \end{array}$$

PROP:

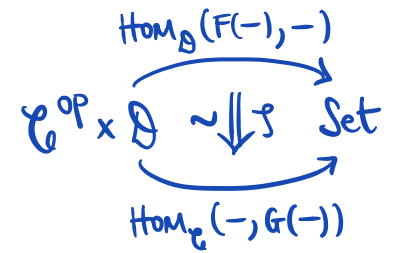


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{X,Y}$ ON RHS.



I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

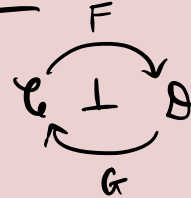
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\ & \searrow id_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\ & \searrow id_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:

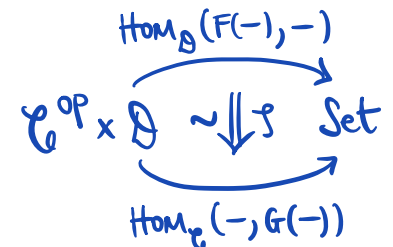


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{X,Y}$ ON RHS.



DEFINE: $\forall X \in \mathcal{C}$ AND $Y \in \mathcal{D}$:

$$X \xrightarrow{\eta_X} GF(X) \quad \& \quad FG(Y) \xrightarrow{\varepsilon_Y} Y$$

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

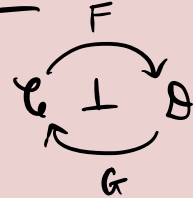
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:

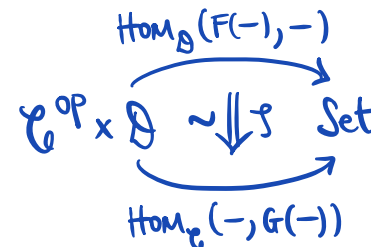


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{X,Y}$ ON RHS.



DEFINE: $\forall X \in \mathcal{C}$ AND $Y \in \mathcal{D}$:

$$X \xrightarrow{\eta_X} \underline{GF(X)} \quad \& \quad FG(Y) \xrightarrow{\varepsilon_Y} Y$$

TAKE Y ABOVE
TO BE $F(X)$ HERE

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

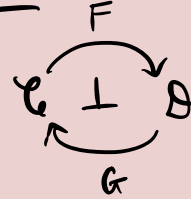
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow id_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \end{array} \quad \forall X \in \mathcal{C}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\ & \searrow id_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \end{array} \quad \forall Y \in \mathcal{D}$$

PROP:

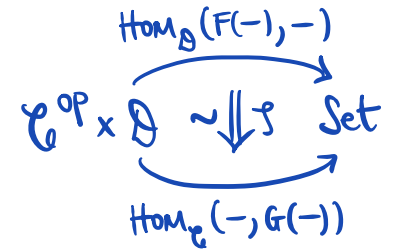


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{X,Y}$ ON RHS.



DEFINE: $\forall X \in \mathcal{C}$ AND $Y \in \mathcal{D}$:

$$X \xrightarrow{\eta_X} GF(X) \quad \& \quad FG(Y) \xrightarrow{\varepsilon_Y} Y$$

$$\exists_{X, F(X)} (id_{F(X)})$$

TAKE Y ABOVE
TO BE $F(X)$ HERE

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$$

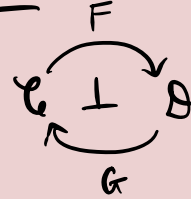
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc}
 F(X) & \xrightarrow{F(\eta_X)} & FG F(X) \\
 \searrow id_{F(X)} & \lrcorner & \downarrow \varepsilon_{F(X)} \\
 & & F(X) \quad \forall X \in \mathcal{C}
 \end{array}$$

$$\begin{array}{ccc}
 G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\
 \searrow id_{G(Y)} & \lrcorner & \downarrow G(\varepsilon_Y) \\
 & & G(Y) \quad \forall Y \in \mathcal{D}
 \end{array}$$

PROP:

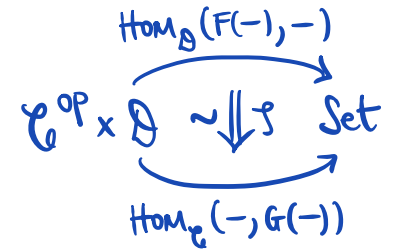


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{X,Y}$ ON RHS.



DEFINE: $\forall X \in \mathcal{C}$ AND $Y \in \mathcal{D}$:

$$X \xrightarrow{\eta_X} GF(X) \quad \& \quad FG(Y) \xrightarrow{\varepsilon_Y} Y$$

$$\exists_{X, F(X)} (id_{F(X)})$$

TAKE Y ABOVE
TO BE $F(X)$ HERE

I. ADJUNCTION : CHARACTERIZATION

$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN **ADJUNCTION**
 $F \dashv G$

IF \exists NAT'L TRANSF'NS

$\eta: Id_{\mathcal{C}} \Rightarrow GF$
ADJ'N UNIT

$\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$
ADJ'N COUNIT

\Rightarrow **Δ IDENTITIES** HOLD:

$$\begin{array}{ccc}
 F(x) & \xrightarrow{F(\eta_x)} & FG F(x) \\
 \searrow id_{F(x)} & \quad \quad \quad \downarrow \varepsilon_{F(x)} & \\
 & & F(x) \quad \forall x \in \mathcal{C}
 \end{array}$$

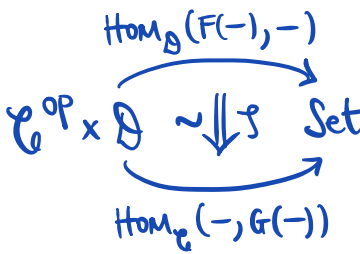
$$\begin{array}{ccc}
 G(y) & \xrightarrow{G(\varepsilon_y)} & GFG(y) \\
 \searrow id_{G(y)} & \quad \quad \quad \downarrow G(\varepsilon_y) & \\
 & & G(y) \quad \forall y \in \mathcal{D}
 \end{array}$$

PROP:

$$\begin{array}{c}
 \mathcal{C} \quad \mathcal{D} \\
 \begin{array}{ccc}
 & \xrightarrow{F} & \\
 \circlearrowleft & \perp & \circlearrowright \\
 & \xleftarrow{G} &
 \end{array}
 \end{array}
 \iff \exists_{x,y} : Hom_{\mathcal{D}}(F(x), y) \xrightarrow{\sim} Hom_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \& y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{x,y}$ ON RHS.



DEFINE: $\forall x \in \mathcal{C}$ AND $y \in \mathcal{D}$:

$$X \xrightarrow{\eta_x} GF(x) \quad \& \quad FG(y) \xrightarrow{\varepsilon_y} Y$$

$$\exists_{x, F(x)} (id_{F(x)})$$

TAKE y ABOVE TO BE $F(x)$ HERE

TAKE x ABOVE TO BE $G(y)$ HERE

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$$F \dashv G$$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF$$

ADJ'N UNIT

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}}$$

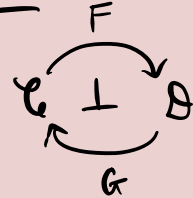
ADJ'N COUNIT

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FGFX \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{FX} \\ & & F(X) \quad \forall X \in \mathcal{C} \end{array}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \quad \forall Y \in \mathcal{D} \end{array}$$

PROP:

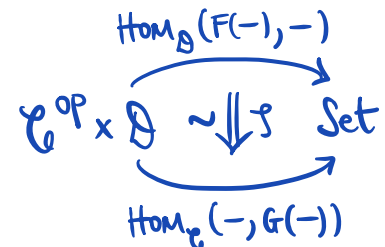


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\exists_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\exists_{X,Y}$ ON RHS.



DEFINE: $\forall X \in \mathcal{C}$ AND $Y \in \mathcal{D}$:

$$\begin{array}{ccc} X & \xrightarrow{\eta_X} & GF(X) \\ \parallel & & \parallel \\ \exists_{X, F(X)} & & (\text{id}_{F(X)}) \end{array}$$

TAKE Y ABOVE
TO BE $F(X)$ HERE

$$\begin{array}{ccc} FG(Y) & \xrightarrow{\varepsilon_Y} & Y \\ \parallel & & \parallel \\ \exists_{G(Y), Y}^{-1} & & (\text{id}_{G(Y)}) \end{array}$$

TAKE X ABOVE
TO BE $G(Y)$ HERE

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$

IF \exists NAT'L TRANSF'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \quad \text{ADJ'N UNIT}$$

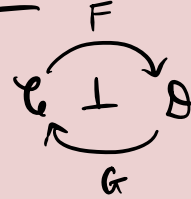
$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \quad \text{ADJ'N COUNIT}$$

\Rightarrow Δ IDENTITIES HOLD:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(\eta_X)} & FG(F(X)) \\ & \searrow \text{id}_{F(X)} & \downarrow \varepsilon_{F(X)} \\ & & F(X) \quad \forall X \in \mathcal{C} \end{array}$$

$$\begin{array}{ccc} G(Y) & \xrightarrow{G(\varepsilon_Y)} & GFG(Y) \\ & \searrow \text{id}_{G(Y)} & \downarrow G(\varepsilon_Y) \\ & & G(Y) \quad \forall Y \in \mathcal{D} \end{array}$$

PROP:

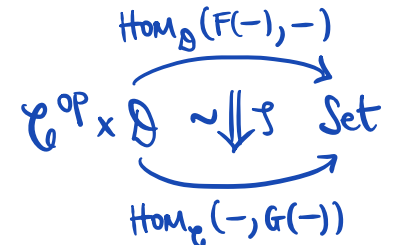


$$\forall X \in \mathcal{C} \quad \& \quad Y \in \mathcal{D} : \exists \text{ BIJECTION}$$

$$\mathcal{J}_{X,Y} : \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$$

NATURAL IN $X \quad \& \quad Y$

PF/ (\Leftarrow) GIVEN BIJECTIONS $\mathcal{J}_{X,Y}$ ON RHS.



DEFINE: $\forall X \in \mathcal{C}$ AND $Y \in \mathcal{D}$:

$$\begin{array}{ccc} X & \xrightarrow{\eta_X} & GF(X) \\ \parallel & & \parallel \\ \mathcal{J}_{X, F(X)}(\text{id}_{F(X)}) & & \mathcal{J}_{G(Y), Y}^{-1}(\text{id}_{G(Y)}) \end{array} \quad \& \quad \begin{array}{ccc} FG(Y) & \xrightarrow{\varepsilon_Y} & Y \\ \parallel & & \parallel \\ \mathcal{J}_{G(Y), Y}^{-1}(\text{id}_{G(Y)}) & & \mathcal{J}_{X, F(X)}(\text{id}_{F(X)}) \end{array}$$

TAKE Y ABOVE
TO BE $F(X)$ HERE

TAKE X ABOVE
TO BE $G(Y)$ HERE

CHECK THAT THESE FORM COMPONENTS OF
THE DESIRED NAT'L TRANS'NS $\eta \quad \& \quad \varepsilon$

EXER. 2.40

I. ADJUNCTION : CHARACTERIZATION

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \quad \text{UNIT}$$

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \quad \text{COUNIT}$$

\Rightarrow Δ IDENTITIES HOLD:

$$\varepsilon_{F(x)} \circ F(\eta_x) = \text{id}_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\varepsilon_y) \circ \eta_{G(y)} = \text{id}_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nabla G: \mathcal{D} \rightarrow \mathcal{C}$
FORM AN **ADJUNCTION**
 $F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\nabla \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \text{ COUNIT}$$

\Rightarrow Δ IDENTITIES HOLD:

$$\varepsilon_{F(x)} \circ F(\eta_x) = \text{id}_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\varepsilon_y) \circ \eta_{G(y)} = \text{id}_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{J}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \nabla y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ IS AN EQUIV. OF CATEGS,
THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nabla G \dashv F$
WHERE UNIT ∇ COUNIT ARE NAT'L ISOMS.

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
FORM AN ADJUNCTION
 $F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\nmid \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \text{ COUNIT}$$

$\Rightarrow \Delta$ IDENTITIES HOLD:

$$\varepsilon_{F(x)} \circ F(\eta_x) = \text{id}_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\varepsilon_y) \circ \eta_{G(y)} = \text{id}_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ IS AN EQUIV. OF CATEGORIES,
THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
WHERE UNIT \nmid COUNIT ARE NAT'L ISOMS.

PF/ $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ EQUIVALENCE

$\Rightarrow \exists$ QUASI-INVERSE $G: \mathcal{D} \rightarrow \mathcal{C}$.

THAT IS: \exists NATURAL ISOMORPHISMS:

$$\phi: \text{Id}_{\mathcal{C}} \xrightarrow{\sim} GF \quad \nmid \quad \psi: FG \xrightarrow{\sim} \text{Id}_{\mathcal{D}}.$$

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN ADJUNCTION
 $F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\nmid \varepsilon: FG \Rightarrow Id_{\mathcal{D}} \text{ COUNIT}$$

$\Rightarrow \Delta$ IDENTITIES HOLD:

$$\varepsilon_{F(x)} \circ F(\eta_x) = id_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\varepsilon_y) \circ \eta_{G(y)} = id_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ IS AN EQUIV. OF CATEGS,
 THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
 WHERE UNIT \nmid COUNIT ARE NAT'L ISOMS.

PF/ $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ EQUIVALENCE

$\Rightarrow \exists$ QUASI-INVERSE $G: \mathcal{D} \rightarrow \mathcal{C}$.

THAT IS: \exists NATURAL ISOMORPHISMS:

$$\phi: Id_{\mathcal{C}} \xrightarrow{\sim} GF \quad \nmid \quad \psi: FG \xrightarrow{\sim} Id_{\mathcal{D}}$$

NOW DEFINE:

$$\text{Hom}_{\mathcal{D}}(F(x), y) \longrightarrow \text{Hom}_{\mathcal{C}}(GF(x), G(y)) \longrightarrow \text{Hom}_{\mathcal{C}}(x, G(y))$$

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN **ADJUNCTION**
 $F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\nmid \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \text{ COUNIT}$$

\Rightarrow **Δ IDENTITIES** HOLD:

$$\varepsilon_{F(x)} \circ F(\eta_x) = \text{id}_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\varepsilon_y) \circ \eta_{G(y)} = \text{id}_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{J}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{D}}(x, G(y))$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ IS AN EQUIV. OF CATEGORIES,
 THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
 WHERE UNIT \nmid COUNIT ARE NAT'L ISOMS.

PF/ $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ EQUIVALENCE

$\Rightarrow \exists$ QUASI-INVERSE $G: \mathcal{D} \rightarrow \mathcal{C}$.

THAT IS: \exists NATURAL ISOMORPHISMS:

$$\phi: \text{Id}_{\mathcal{C}} \xrightarrow{\sim} GF \quad \nmid \quad \psi: FG \xrightarrow{\sim} \text{Id}_{\mathcal{D}}$$

NOW DEFINE:

$$\text{Hom}_{\mathcal{D}}(F(x), y) \longrightarrow \text{Hom}_{\mathcal{C}}(GF(x), G(y)) \longrightarrow \text{Hom}_{\mathcal{C}}(x, G(y))$$

$$[f: F(x) \rightarrow y] \longmapsto [G(f): GF(x) \rightarrow G(y)]$$

$$[h: GF(x) \rightarrow G(y)] \longmapsto [x \xrightarrow{\phi_x} GF(x) \xrightarrow{h} G(y)]$$

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN **ADJUNCTION**
 $F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: Id_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\nmid \epsilon: FG \Rightarrow Id_{\mathcal{D}} \text{ COUNIT}$$

\Rightarrow **Δ IDENTITIES** HOLD:

$$\epsilon_{F(x)} \circ F(\eta_x) = id_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\epsilon_y) \circ \eta_{G(y)} = id_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: Hom_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\cong} Hom_{\mathcal{D}}(x, G(y))$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\cong} \mathcal{D}$ IS AN EQUIV. OF CATEGORIES,
 THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
 WHERE UNIT \nmid COUNIT ARE NAT'L ISOMS.

PF/ $F: \mathcal{C} \xrightarrow{\cong} \mathcal{D}$ EQUIVALENCE

$\Rightarrow \exists$ QUASI-INVERSE $G: \mathcal{D} \rightarrow \mathcal{C}$.

THAT IS: \exists NATURAL ISOMORPHISMS:

$$\phi: Id_{\mathcal{C}} \xrightarrow{\cong} GF \quad \nmid \quad \psi: FG \xrightarrow{\cong} Id_{\mathcal{D}}$$

NOW DEFINE:

$$\begin{array}{c}
 \xrightarrow{\mathcal{I}_{x,y} \text{ AS COMPOSITION}} \\
 \xrightarrow{\cong \text{ FULLY FAITHFUL}} \\
 Hom_{\mathcal{D}}(F(x), y) \longrightarrow Hom_{\mathcal{D}}(GF(x), G(y)) \longrightarrow Hom_{\mathcal{D}}(x, G(y)) \\
 [f: F(x) \rightarrow y] \longmapsto [G(f): GF(x) \rightarrow G(y)] \\
 [h: GF(x) \rightarrow G(y)] \longmapsto [x \xrightarrow{\phi_x} GF(x) \xrightarrow{h} G(y)]
 \end{array}$$

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN **ADJUNCTION**
 $F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\nmid \epsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \text{ COUNIT}$$

\exists **Δ IDENTITIES** HOLD:

$$\epsilon_{F(x)} \circ F(\eta_x) = \text{id}_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\epsilon_y) \circ \eta_{G(y)} = \text{id}_{G(y)} \quad \forall y \in \mathcal{D}$$

\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\begin{aligned} \mathcal{S}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y) \\ \xrightarrow{\cong} \text{Hom}_{\mathcal{D}}(x, G(y)) \end{aligned}$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\cong} \mathcal{D}$ IS AN EQUIV. OF CATEGS,
 THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
 WHERE UNIT \nmid COUNIT ARE NAT'L ISOMS.

PF/ $F: \mathcal{C} \xrightarrow{\cong} \mathcal{D}$ EQUIVALENCE

$\Rightarrow \exists$ QUASI-INVERSE $G: \mathcal{D} \rightarrow \mathcal{C}$.

THAT IS: \exists NATURAL ISOMORPHISMS:

$$\phi: \text{Id}_{\mathcal{C}} \xrightarrow{\cong} GF \quad \nmid \quad \psi: FG \xrightarrow{\cong} \text{Id}_{\mathcal{D}}$$

NOW DEFINE:

$$\begin{aligned} & \xrightarrow[\text{AS COMPOSITION}]{\mathcal{S}_{x,y}} \\ \text{Hom}_{\mathcal{D}}(F(x), y) & \xrightarrow[\cong \text{ FULLY FAITHFUL}]{\mathcal{S}_{x,y}} \text{Hom}_{\mathcal{D}}(GF(x), G(y)) \rightarrow \text{Hom}_{\mathcal{D}}(x, G(y)) \\ [f: F(x) \rightarrow y] & \mapsto [G(f): GF(x) \rightarrow G(y)] \\ [h: GF(x) \rightarrow G(y)] & \mapsto [x \xrightarrow[\cong]{\phi_x} GF(x) \xrightarrow{h} G(y)] \end{aligned}$$

DEFINE $\mathcal{S}_{x,y}^{-1}$ USING $\psi \dots \equiv$

EXER. 2.41

I. ADJUNCTION \Leftarrow EQUIVALENCE

$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
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\equiv OR EQUIVALENTLY \equiv

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\cong} \text{Hom}_{\mathcal{D}}(x, G(y))$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\cong} \mathcal{D}$ IS AN EQUIV. OF CATEGS,
 THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
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PF/ $F: \mathcal{C} \xrightarrow{\cong} \mathcal{D}$ EQUIVALENCE

$\Rightarrow \exists$ QUASI-INVERSE $G: \mathcal{D} \rightarrow \mathcal{C}$.

THAT IS: \exists NATURAL ISOMORPHISMS:

$$\phi: \text{Id}_{\mathcal{C}} \xrightarrow{\cong} GF \quad \nmid \quad \psi: FG \xrightarrow{\cong} \text{Id}_{\mathcal{D}}$$



NOT THE
 CO/UNIT OF
 $F \dashv G$

NOW DEFINE:

$$\text{Hom}_{\mathcal{D}}(F(x), y) \xrightarrow[\cong]{\mathcal{I}_{x,y} \text{ AS COMPOSITION}} \text{Hom}_{\mathcal{C}}(GF(x), G(y)) \xrightarrow[\cong]{\text{G FULLY FAITHFUL}} \text{Hom}_{\mathcal{D}}(x, G(y))$$

$$[f: F(x) \rightarrow y] \mapsto [G(f): GF(x) \rightarrow G(y)]$$

$$[h: GF(x) \rightarrow G(y)] \mapsto [x \xrightarrow[\cong]{\phi_x} GF(x) \xrightarrow{h} G(y)]$$

DEFINE $\mathcal{I}_{x,y}^{-1}$ USING $\psi \dots //$

EXER. 2.41

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$F: \mathcal{C} \rightarrow \mathcal{D} \nmid G: \mathcal{D} \rightarrow \mathcal{C}$
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\equiv OR EQUIVALENTLY \equiv

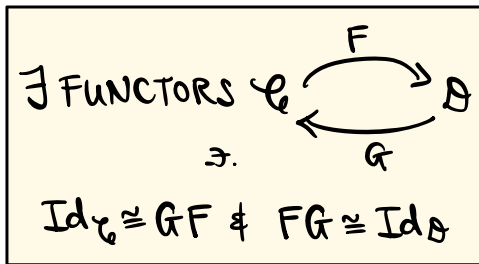
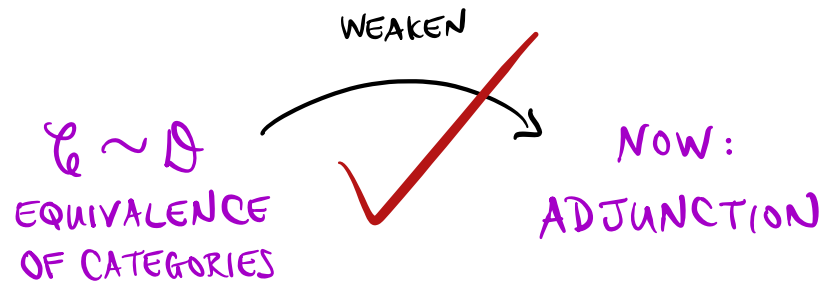
\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

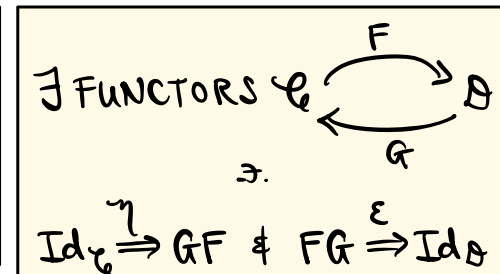
$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \nmid y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

PROP IF $F: \mathcal{C} \xrightarrow{\sim} \mathcal{D}$ IS AN EQUIV. OF CATEGS,
 THEN $\exists G: \mathcal{D} \rightarrow \mathcal{C}$ WITH $F \dashv G \nmid G \dashv F$
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INVOLVING
ISOMORPHISMS
OF FUNCTORS



HAVE A NICE
COMPATIBILITY CONDITION

INVOLVES CERTAIN TRANS'NS
OF FUNCTORS

I. ADJUNCTION : EXAMPLES

≡ TENSOR-HOM ADJUNCTION ≡ (LECT #2)
FOR \mathbb{R} -VECTOR SPACES U, V, W

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \quad \text{UNIT}$$

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \quad \text{COUNIT}$$

\Rightarrow Δ IDENTITIES HOLD:

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≡ OR EQUIVALENTLY ≡

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

I. ADJUNCTION : EXAMPLES

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≡ OR EQUIVALENTLY ≡

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 NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

$$\text{Hom}_{\mathbb{R}}(U \otimes_{\mathbb{R}} V, W) \cong \text{Hom}_{\mathbb{R}}(U, \text{Hom}_{\mathbb{R}}(V, W))$$

$$f \mapsto \left[\begin{array}{ccc} U & \longrightarrow & \text{Hom}_{\mathbb{R}}(V, W) \\ u \mapsto & & [v \mapsto f(u \otimes v)] \end{array} \right]$$

$$\left[\begin{array}{ccc} U \otimes V & \longrightarrow & W \\ u \otimes v \mapsto & & g(u)(v) \end{array} \right] \longleftarrow g$$

I. ADJUNCTION : EXAMPLES

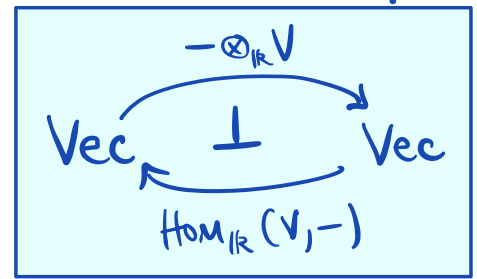
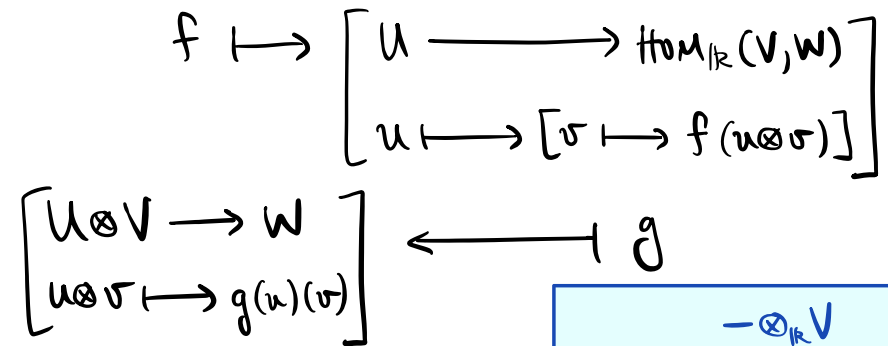
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$$Hom_{\mathbb{R}}(U \otimes_{\mathbb{R}} V, W) \cong Hom_{\mathbb{R}}(U, Hom_{\mathbb{R}}(V, W))$$



I. ADJUNCTION : EXAMPLES

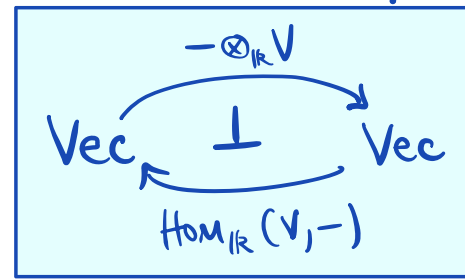
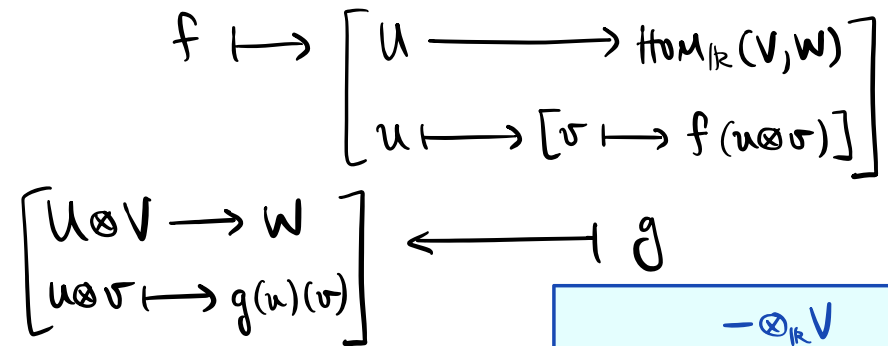
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 NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

$$\text{Hom}_{\mathbb{R}}(U \otimes_{\mathbb{R}} V, W) \cong \text{Hom}_{\mathbb{R}}(U, \text{Hom}_{\mathbb{R}}(V, W))$$



ALSO GET:

$$\text{Hom}_{\mathbb{R}}(U \otimes_{\mathbb{R}} V, W) \cong \text{Hom}_{\mathbb{R}}(V, \text{Hom}_{\mathbb{R}}(U, W))$$

I. ADJUNCTION : EXAMPLES

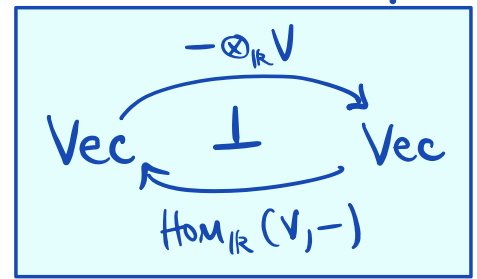
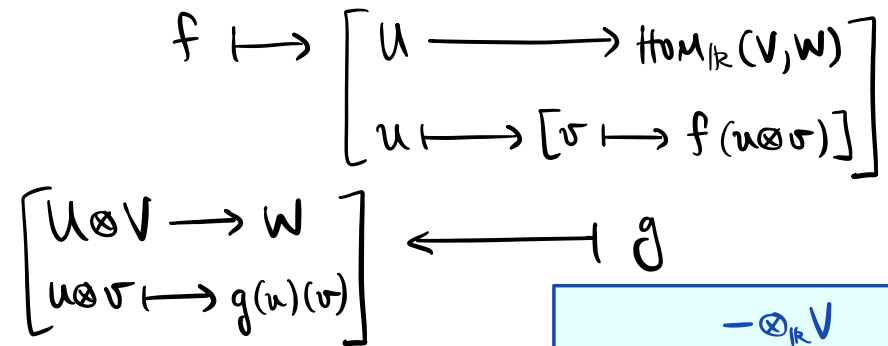
\equiv TENSOR-HOM ADJUNCTION \equiv (LECT #2)
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\equiv OR EQUIVALENTLY \equiv

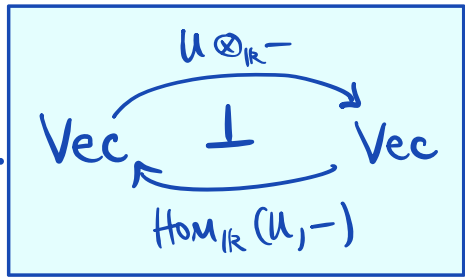
\exists BIJECTIONS
 $\mathcal{I}_{x,y}: Hom_{\mathcal{D}}(F(x), y) \xrightarrow{\cong} Hom_{\mathcal{C}}(x, G(y))$
 NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

$$\text{Hom}_{\mathbb{R}}(U \otimes_{\mathbb{R}} V, W) \cong \text{Hom}_{\mathbb{R}}(U, \text{Hom}_{\mathbb{R}}(V, W))$$



ALSO GET:

$$\text{Hom}_{\mathbb{R}}(U \otimes_{\mathbb{R}} V, W) \cong \text{Hom}_{\mathbb{R}}(V, \text{Hom}_{\mathbb{R}}(U, W))$$



I. ADJUNCTION : EXAMPLES

≡ BIMODULE TENSOR-HOM ADJUN ≡ (LECT #4)

$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \text{ UNIT}$$

$$\& \quad \varepsilon: FG \Rightarrow \text{Id}_{\mathcal{D}} \text{ COUNIT}$$

\Rightarrow Δ IDENTITIES HOLD:

$$\varepsilon_{F(x)} \circ F(\eta_x) = \text{id}_{F(x)} \quad \forall x \in \mathcal{C}$$

$$G(\varepsilon_y) \circ \eta_{G(y)} = \text{id}_{G(y)} \quad \forall y \in \mathcal{D}$$

≡ OR EQUIVALENTLY ≡

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

I. ADJUNCTION : EXAMPLES

≡ BIMODULE TENSOR-HOM ADJ'N ≡ (LECT #4)

$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$
FORM AN ADJUNCTION
 $F \dashv G$ IF \exists NAT. TRANS'NS
 $\eta: Id_{\mathcal{C}} \Rightarrow GF$ UNIT
 $\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$ COUNIT
 $\Rightarrow \Delta$ IDENTITIES HOLD:
 $\varepsilon_{F(x)} \circ F(\eta_x) = id_{F(x)} \quad \forall x \in \mathcal{C}$
 $G(\varepsilon_y) \circ \eta_{G(y)} = id_{G(y)} \quad \forall y \in \mathcal{D}$

≡ OR EQUIVALENTLY ≡

\exists BIJECTIONS
 $\mathcal{I}_{x,y}: Hom_{\mathcal{D}}(F(x), y)$
 $\xrightarrow{\sim} Hom_{\mathcal{C}}(x, G(y))$
NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

$$\text{Hom}_{B_1\text{-mod}}(U \otimes_A V, W) \cong \text{Hom}_{A\text{-mod}}(V, \text{Hom}_{B_1\text{-mod}}(U, W))$$
$$U = {}_{B_1}U_A \quad V = {}_A V \quad W = {}_{B_1}W$$

I. ADJUNCTION : EXAMPLES

≡ BIMODULE TENSOR-HOM ADJUN ≡ (LECT #4)

$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN ADJUNCTION
 $F \dashv G$ IF \exists NAT. TRANS'NS
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≡ OR EQUIVALENTLY ≡

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 $\mathcal{I}_{x,y}: Hom_{\mathcal{D}}(F(x), y)$
 $\xrightarrow{\sim} Hom_{\mathcal{C}}(x, G(y))$
 NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

$$Hom_{B_1\text{-mod}}(U \otimes_A V, W) \cong Hom_{A\text{-mod}}(V, Hom_{B_1\text{-mod}}(U, W))$$

$$U = {}_{B_1}U_A \quad V = {}_A V \quad W = {}_{B_1}W$$

$$F := U \otimes_A - : A\text{-Mod} \longrightarrow B_1\text{-Mod}$$

$$G := Hom_{B_1\text{-mod}}(U, -) : B_1\text{-Mod} \longrightarrow A\text{-Mod}$$

I. ADJUNCTION : EXAMPLES

≡ BIMODULE TENSOR-HOM ADJ'N ≡ (LECT #4)

$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN **ADJUNCTION**
 $F \dashv G$ IF \exists NAT. TRANS'NS
 $\eta: Id_{\mathcal{C}} \Rightarrow GF$ **UNIT**
 $\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$ **COUNIT**
 $\Rightarrow \Delta$ **IDENTITIES** HOLD:
 $\varepsilon_{F(x)} \circ F(\eta_x) = id_{F(x)} \quad \forall x \in \mathcal{C}$
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≡ OR EQUIVALENTLY ≡

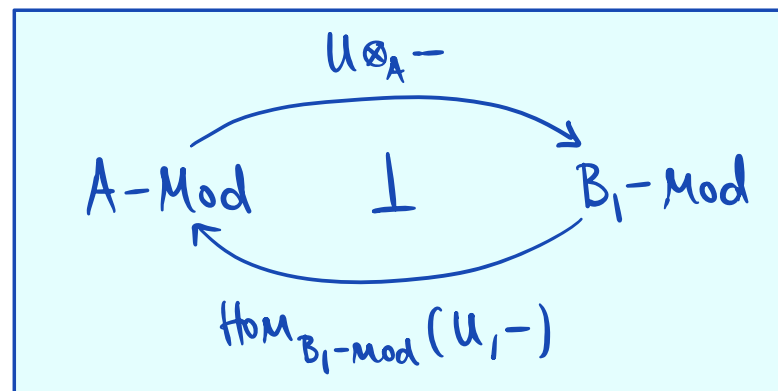
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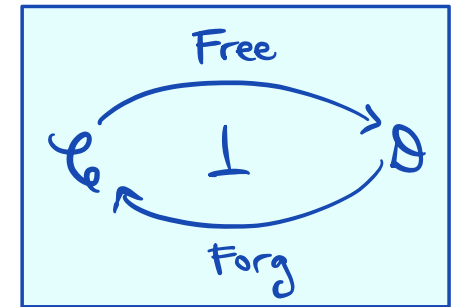
$$F := U \otimes_A - : A\text{-Mod} \longrightarrow B_1\text{-Mod}$$

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I. ADJUNCTION : EXAMPLES

≡ FREE - FORGET ADJUNCTION ≡



$$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$$

FORM AN ADJUNCTION

$F \dashv G$ IF \exists NAT. TRANS'NS

$$\eta: \text{Id}_{\mathcal{C}} \Rightarrow GF \quad \text{UNIT}$$

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≡ OR EQUIVALENTLY ≡

\exists BIJECTIONS

$$\mathcal{I}_{x,y}: \text{Hom}_{\mathcal{D}}(F(x), y)$$

$$\xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(x, G(y))$$

NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

I. ADJUNCTION : EXAMPLES

≡ FREE-FORGET ADJUNCTION ≡

$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$
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≡ OR EQUIVALENTLY ≡

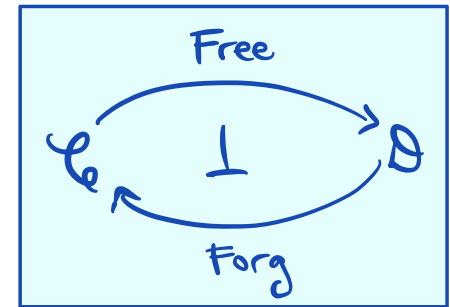
\exists BIJECTIONS
 $\mathcal{I}_{x,y}: Hom_{\mathcal{D}}(F(x), y)$
 $\xrightarrow{\sim} Hom_{\mathcal{C}}(x, G(y))$
NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

A FREE OBJECT ON $x \in \mathcal{C}$

IS AN OBJECT $Free_x \in \mathcal{D}$

EQUIPPED WITH

A MONO $\alpha_x: X \rightarrow Forg(Free_x) \in \mathcal{C}$



I. ADJUNCTION : EXAMPLES

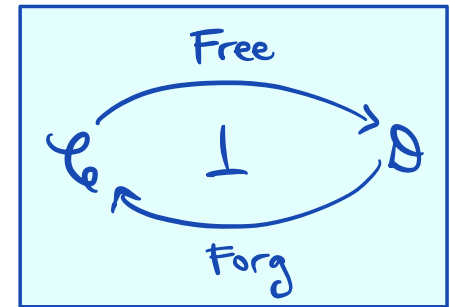
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≡ OR EQUIVALENTLY ≡

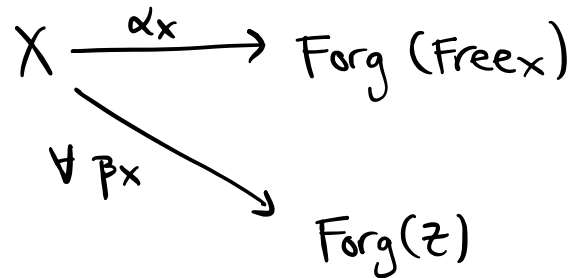
\exists BIJECTIONS
 $\mathcal{I}_{x,y}: Hom_{\mathcal{D}}(F(x), y)$
 $\xrightarrow{\sim} Hom_{\mathcal{C}}(x, G(y))$
 NATURAL IN $x \& y \quad \forall x \in \mathcal{C}, y \in \mathcal{D}$

A FREE OBJECT ON $x \in \mathcal{C}$
 IS AN OBJECT $Free_x \in \mathcal{D}$
 EQUIPPED WITH



A MONO $\alpha_x: X \rightarrow Forg(Free_x) \in \mathcal{C}$

$\exists \quad \forall \beta_x: X \rightarrow Forg(z) \in \mathcal{C}$ WITH $z \in \mathcal{D}$



I. ADJUNCTION : EXAMPLES

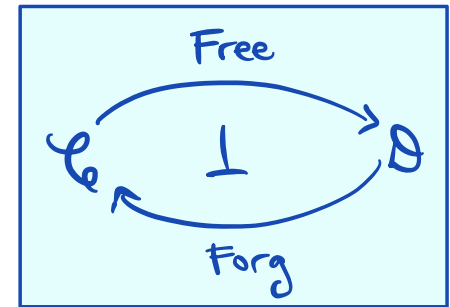
≡ FREE-FORGET ADJUNCTION ≡

$F: \mathcal{C} \rightarrow \mathcal{D} \quad \& \quad G: \mathcal{D} \rightarrow \mathcal{C}$
 FORM AN ADJUNCTION
 $F \dashv G$ IF \exists NAT. TRANS'NS
 $\eta: Id_{\mathcal{C}} \Rightarrow GF$ UNIT
 $\& \quad \varepsilon: FG \Rightarrow Id_{\mathcal{D}}$ COUNIT
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 $\varepsilon_{F(x)} \circ F(\eta_x) = id_{F(x)} \quad \forall x \in \mathcal{C}$
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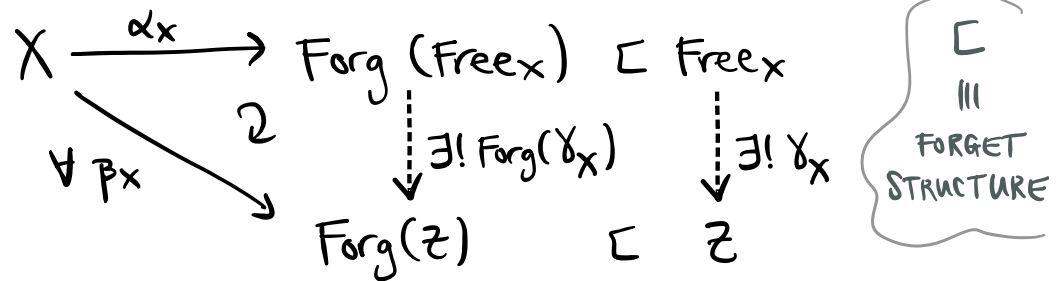
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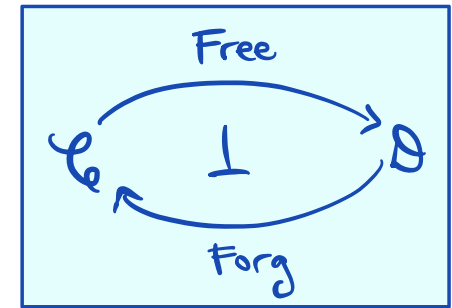
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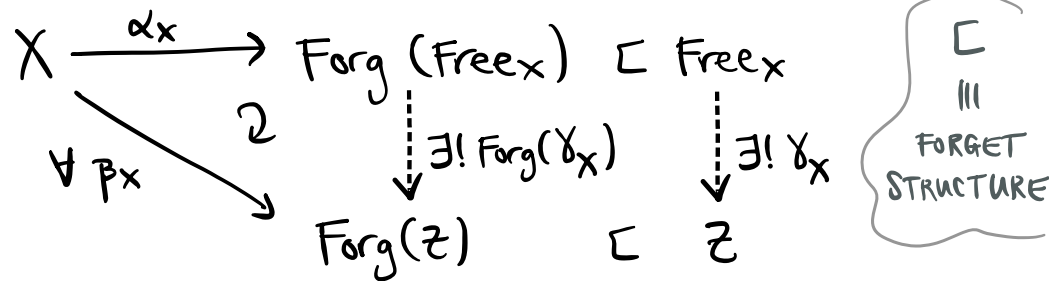
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Free: $\mathcal{C} \rightarrow \mathcal{D}$ IS A **FREE FUNCTOR**

$X \mapsto Free_x$ (IF IT INDEED IS A FUNCTOR)

I. ADJUNCTION : EXAMPLES

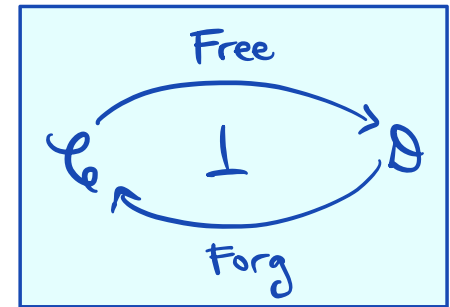
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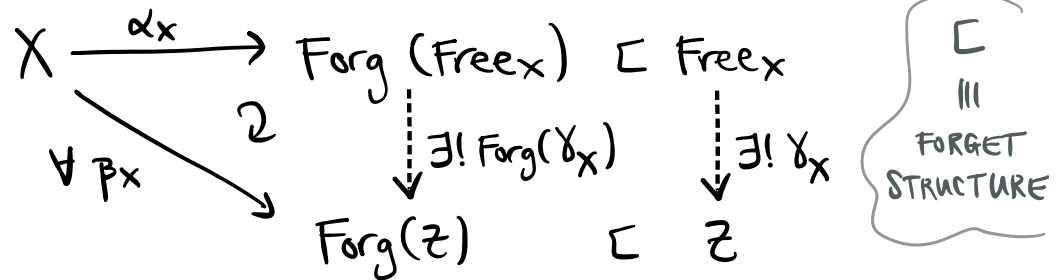
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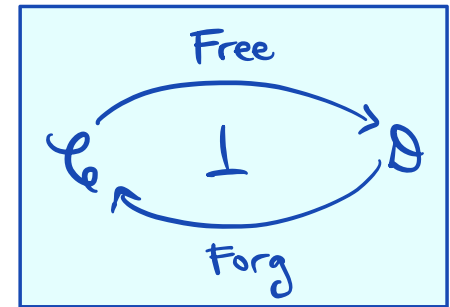
GET: $\gamma_x \xleftarrow{!} \beta_x$

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I. ADJUNCTION : EXAMPLES

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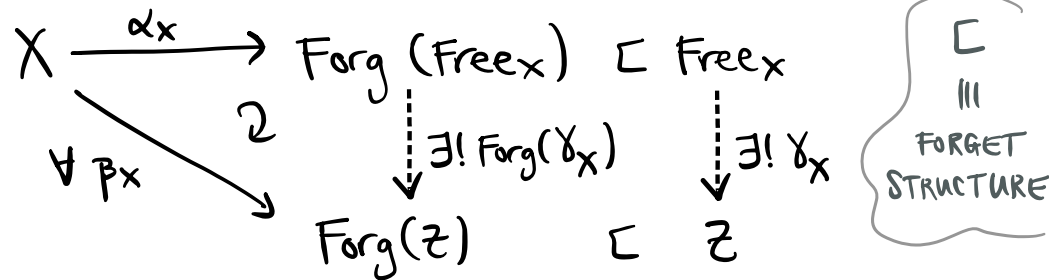
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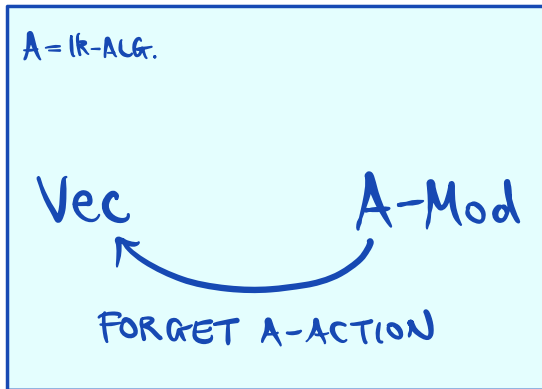
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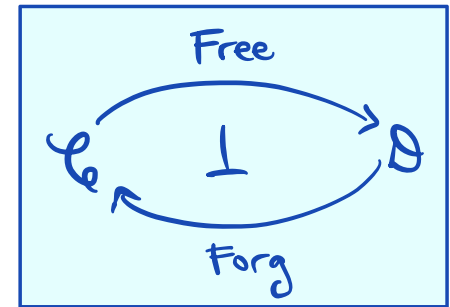
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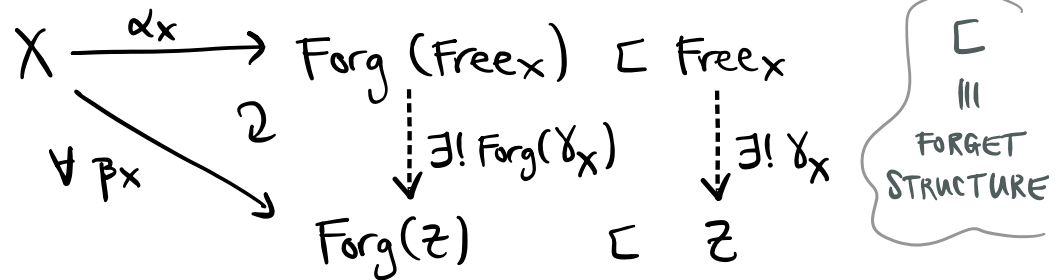
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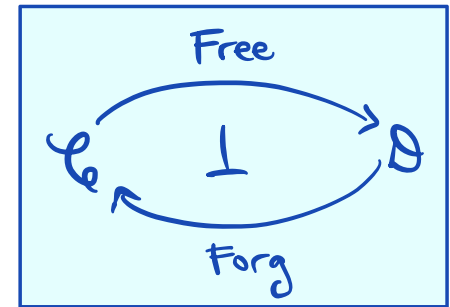
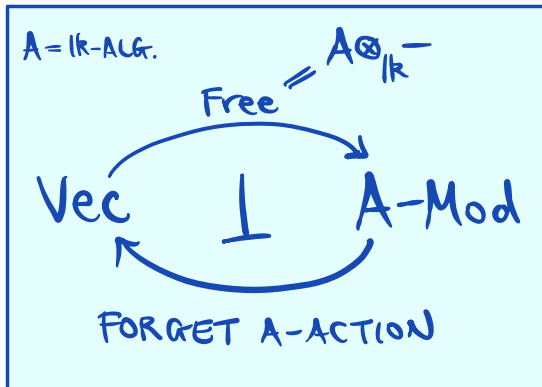
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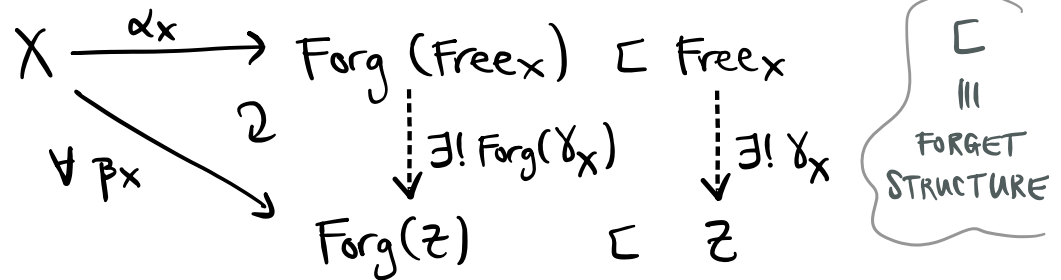


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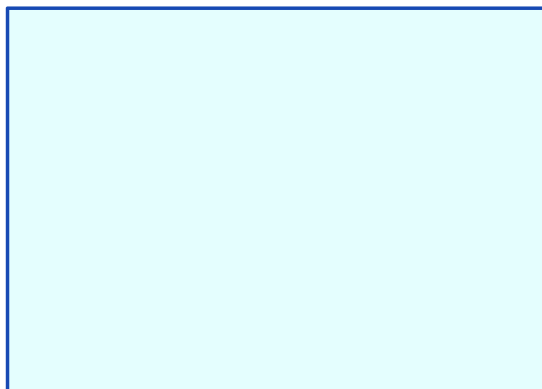
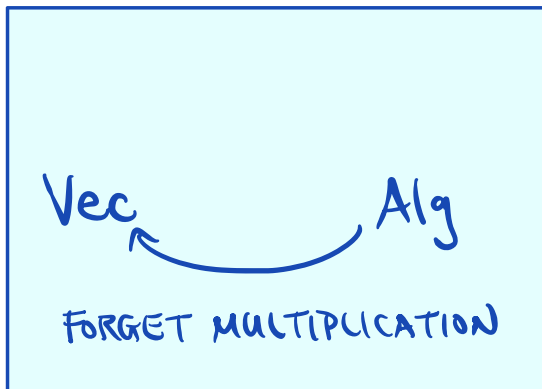
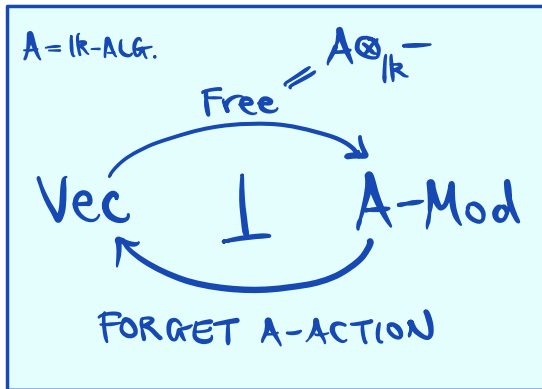
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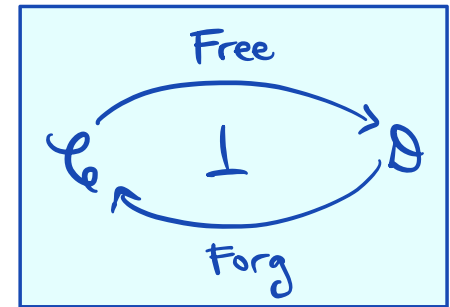
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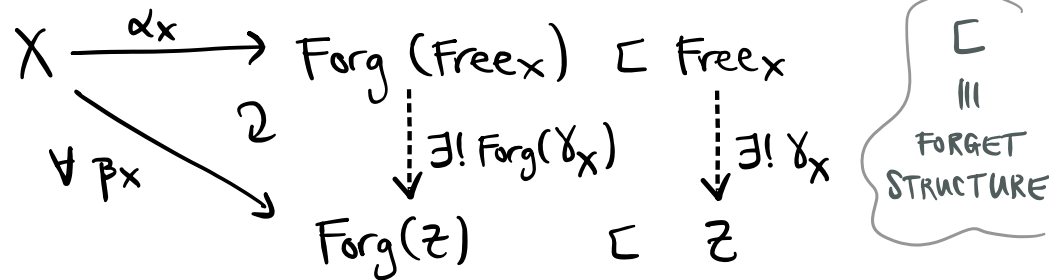
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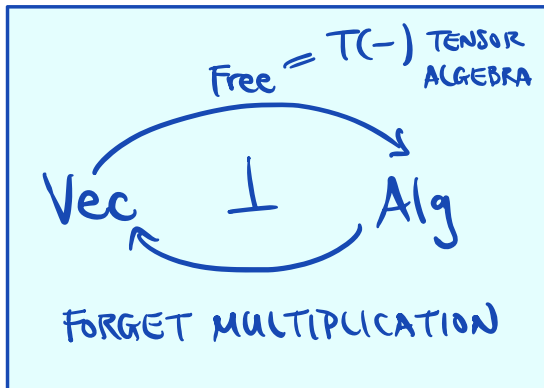
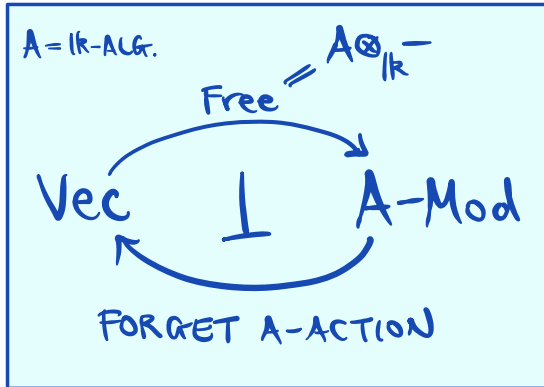
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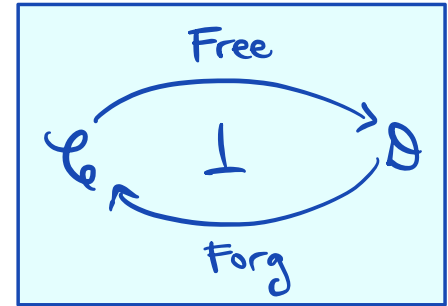
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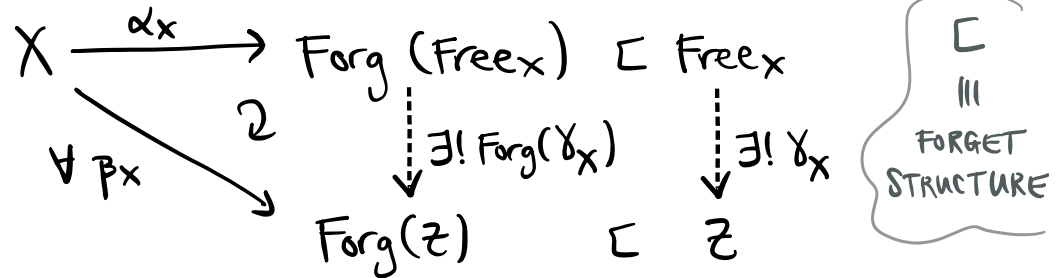
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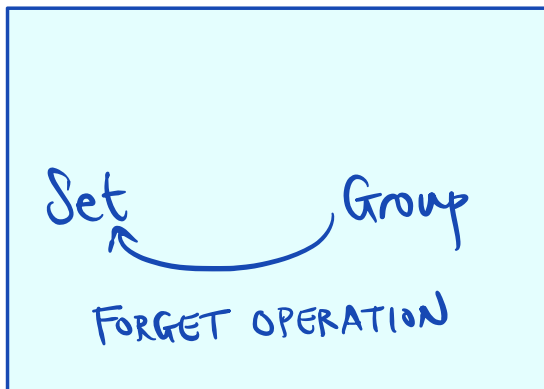
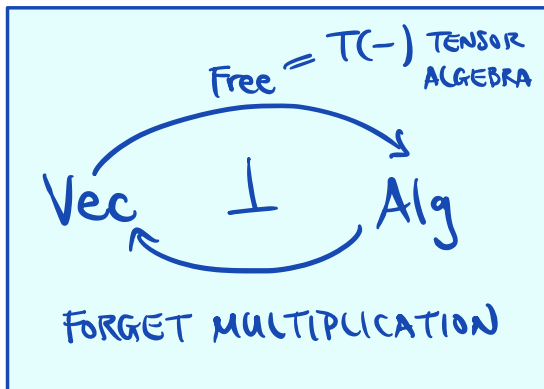
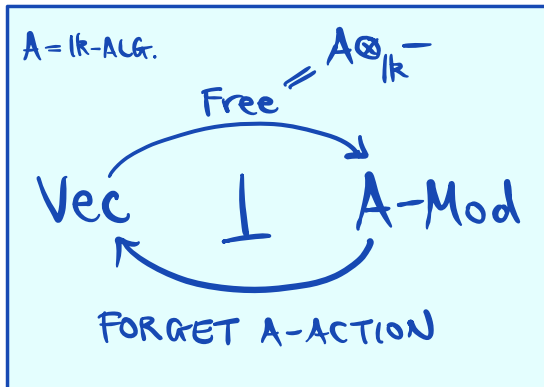
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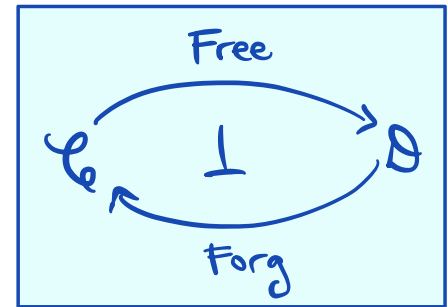
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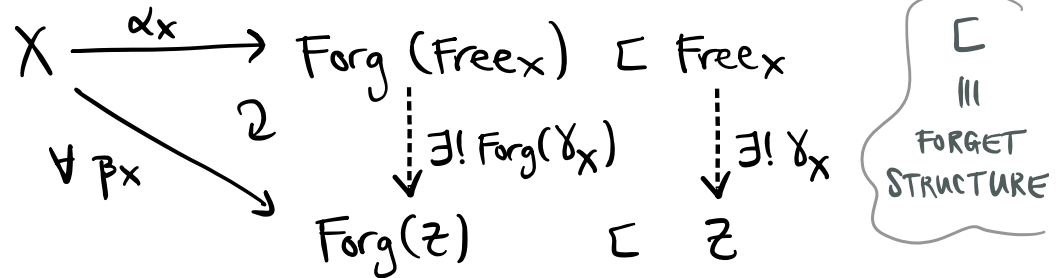


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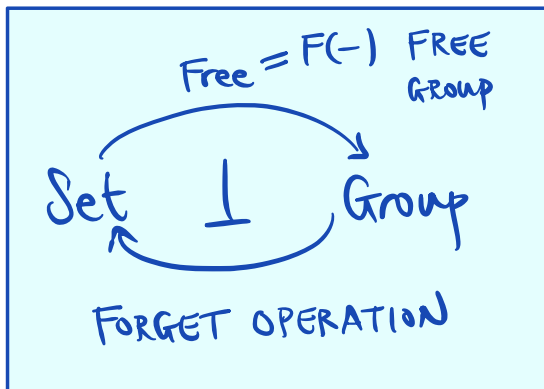
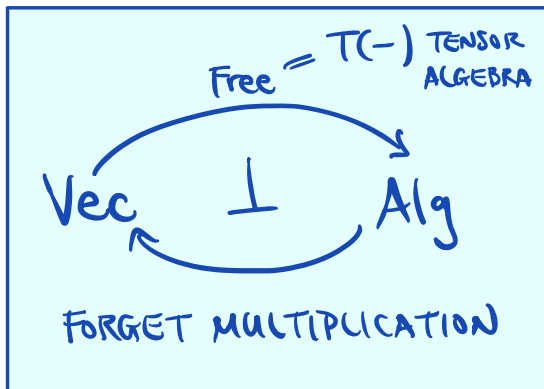
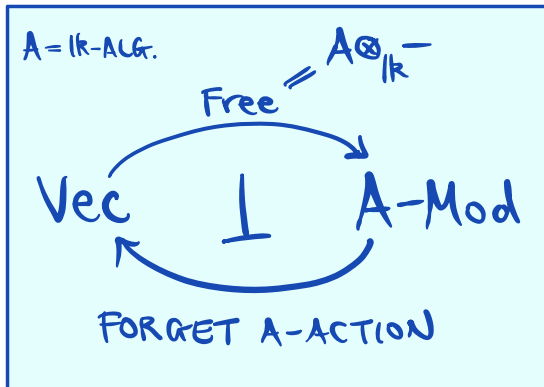
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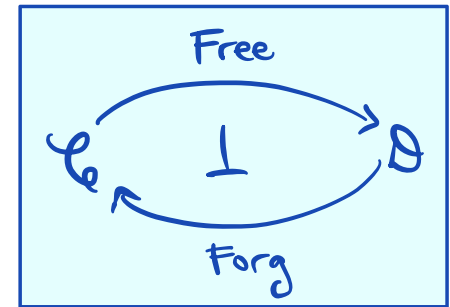
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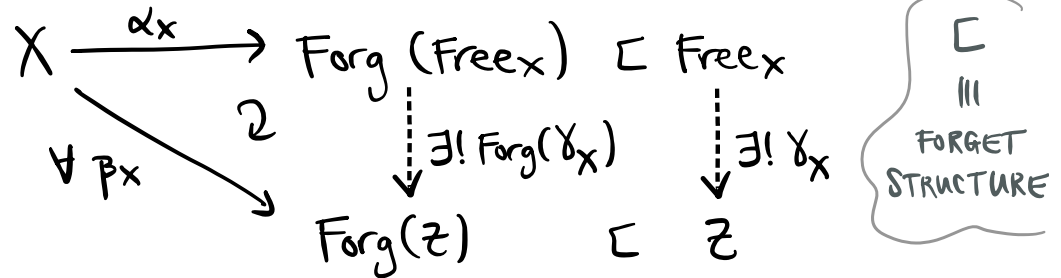
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LEFT ADJOINTS

F

PRESERVE

I. ADJUNCTION

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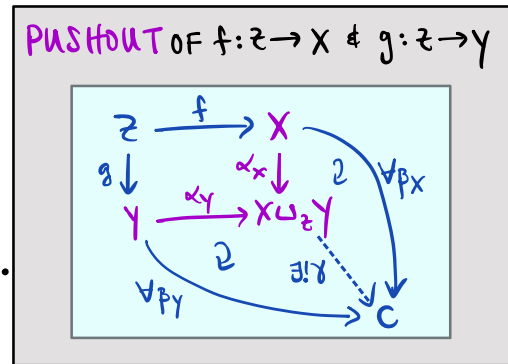
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LEFT ADJOINTS
 F
 PRESERVE

- PUSHOUTS ...



I. ADJUNCTION

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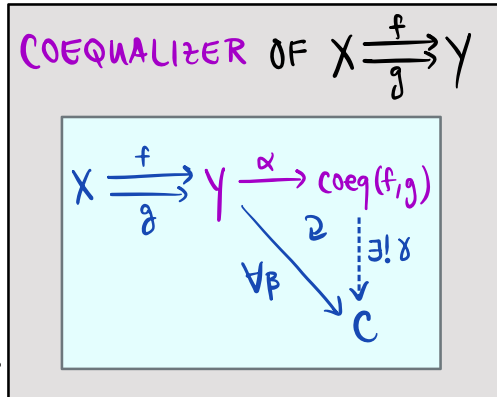
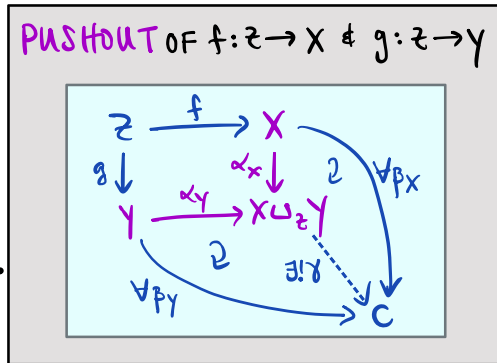
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LEFT ADJOINTS
 F
 PRESERVE

- PUSHOUTS
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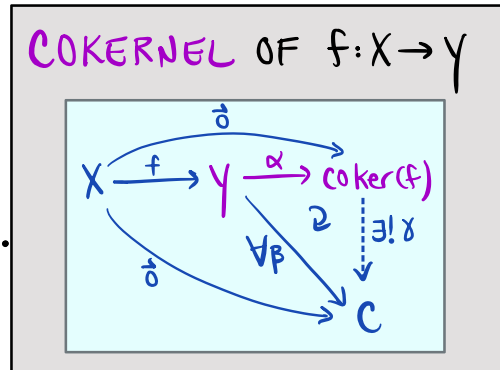
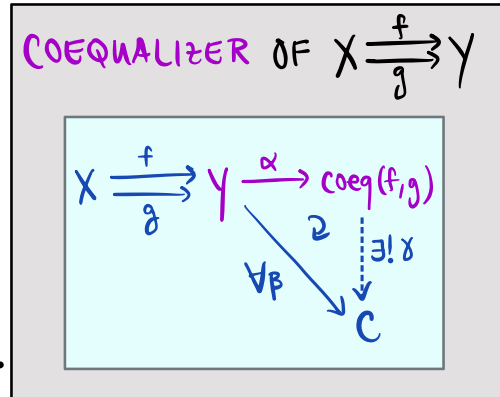
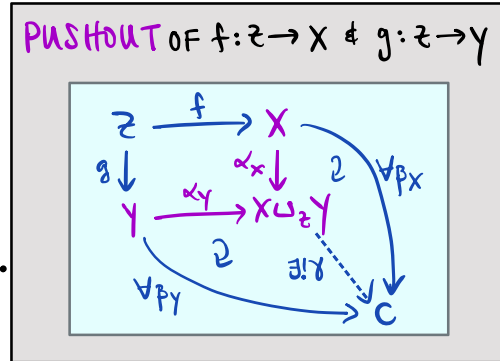
\exists BIJECTIONS
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COOL FACT: FOR $F \dashv G$



LEFT ADJOINTS
 F
 PRESERVE

- PUSHOUTS
- COEQUALIZERS
- COKERNELS



I. ADJUNCTION

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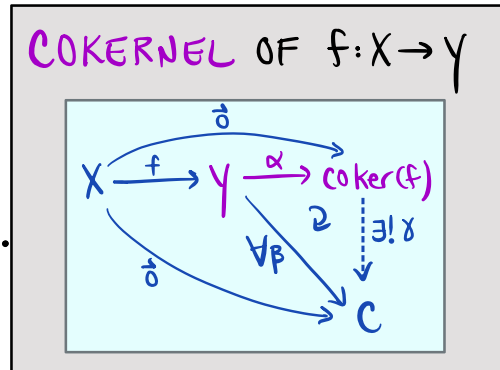
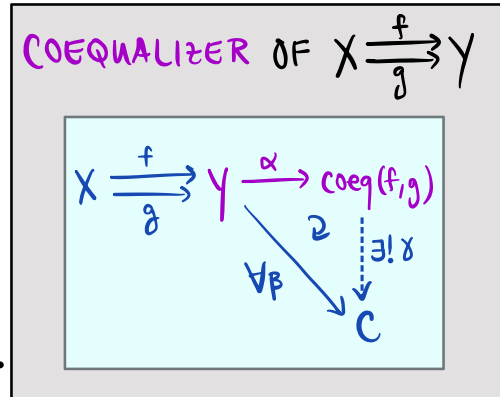
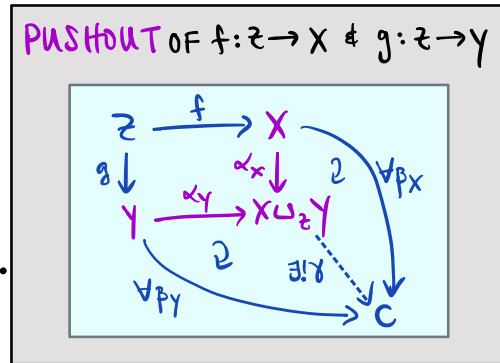
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~~**F**~~
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E.G. COKERNEL OF $F(f)$
 $= F(\alpha: Y \rightarrow \text{coker}(f))$

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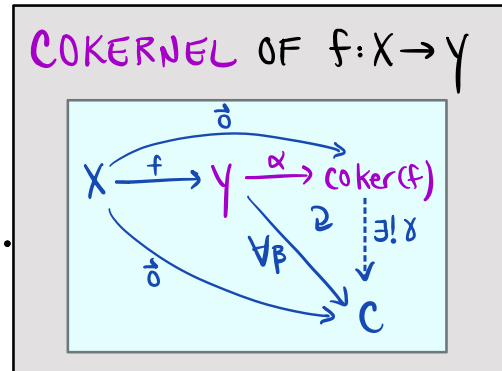
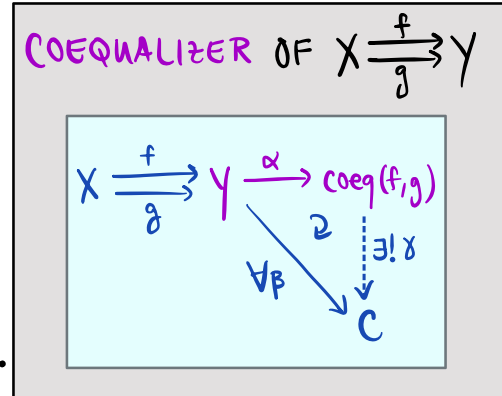
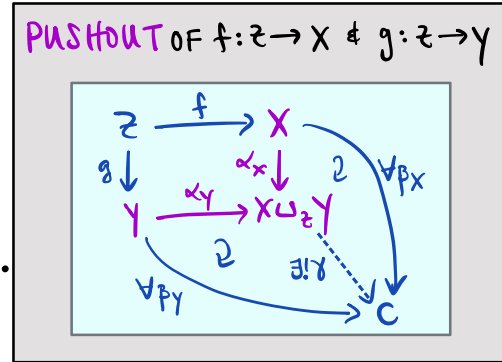
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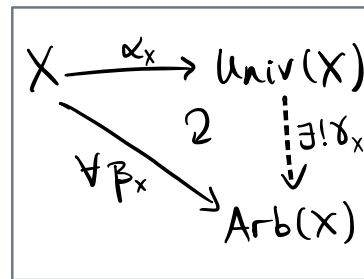
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LEFT ADJOINTS

F

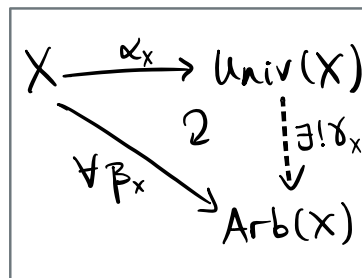
PRESERVE

RIGHT ADJOINTS

G

PRESERVE

- PUSHOUTS
- COEQUALIZERS
- COKERNELS
- ...



FORM I

I. ADJUNCTION

COOL FACT: FOR $F \dashv G$

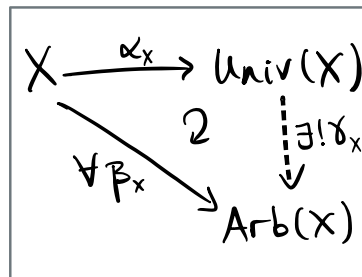
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LEFT ADJOINTS
F
 PRESERVE

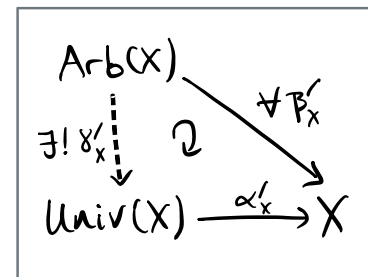
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- COEQUALIZERS
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- ...



FORM I

RIGHT ADJOINTS
G
 PRESERVE

- PULLBACKS
- EQUALIZERS
- KERNELS
- ...



FORM II

I. ADJUNCTION

COOL FACT: FOR $F \dashv G$

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LEFT ADJOINTS
 F
 PRESERVE

RIGHT ADJOINTS
 G
 PRESERVE

- PUSHOUTS
- COEQUALIZERS
- COKERNELS
- ...

- PULLBACKS
- EQUALIZERS
- KERNELS
- ...

"COLIMITS"

"LIMITS"

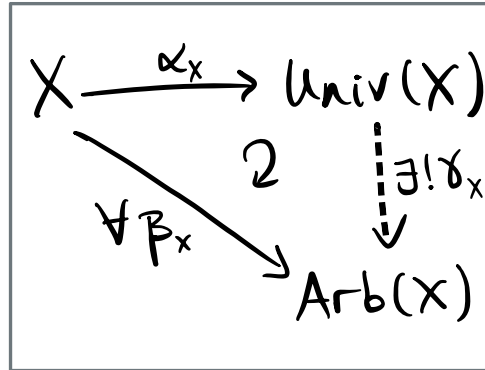
\equiv READ ABOUT THIS \equiv

II. UNIVERSALITY REVISITED

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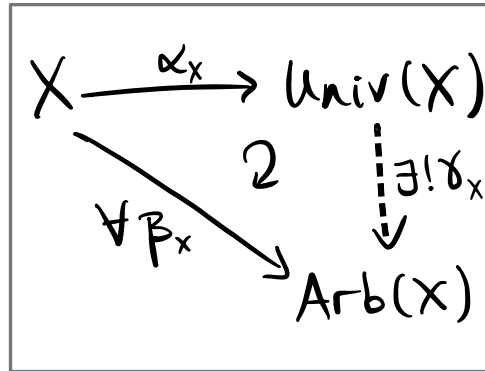
FORM I

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FORM I

GET:

$$\gamma_x \xleftrightarrow{|\cdot|} \beta_x$$

YIELDS:

$$Hom_{\text{Structure}}(\text{Univ}(X), \text{Arb}(X))$$

|||

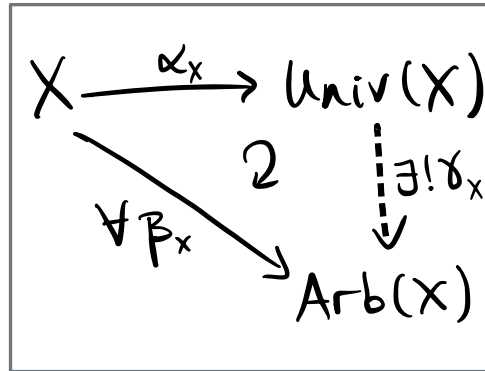
$$Hom_{\text{Gadget}}(X, \text{Arb}(X)_{\text{Gadget}})$$

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FORM I

GET:

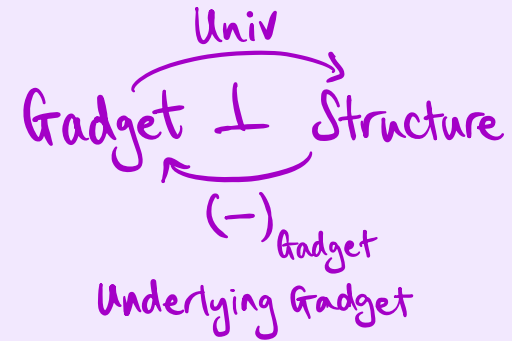
$$\delta_x \xleftarrow{|-|} \beta_x$$

YIELDS:

$$Hom_{Structure}(Univ(X), Arb(X))$$

|||

$$Hom_{Gadget}(X, Arb(X)_{Gadget})$$



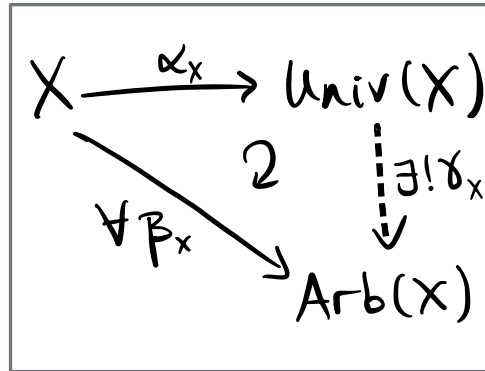
II. UNIVERSALITY REVISITED

(SIMILAR FOR FORM II...)

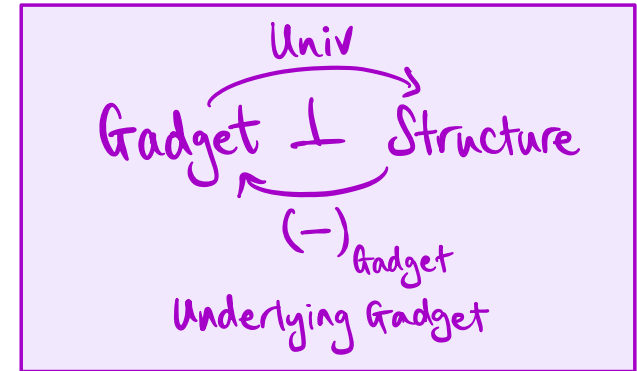
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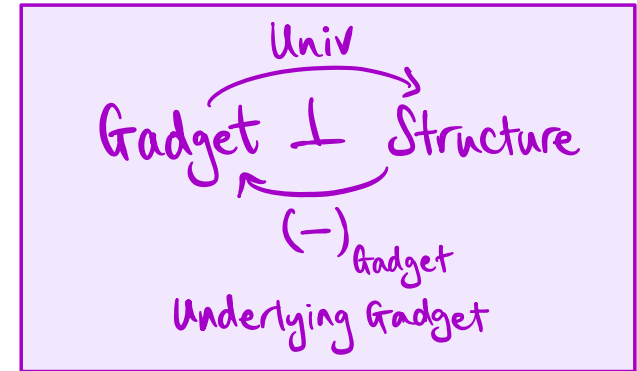
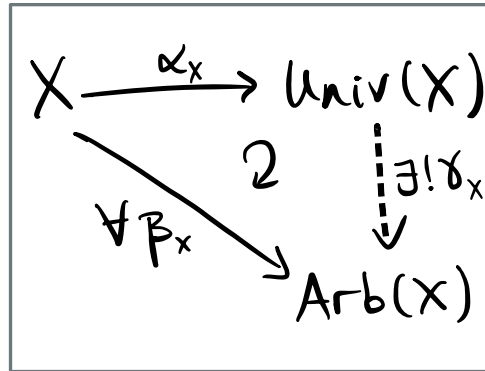
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EXAMPLES —

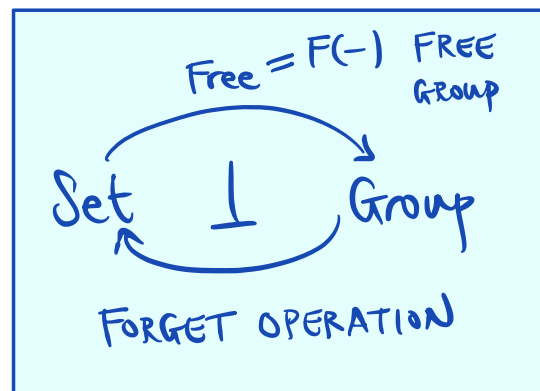
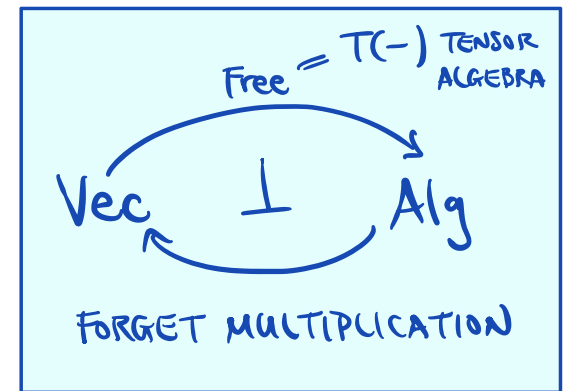
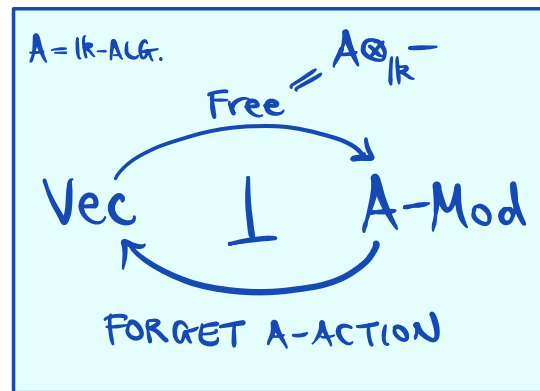
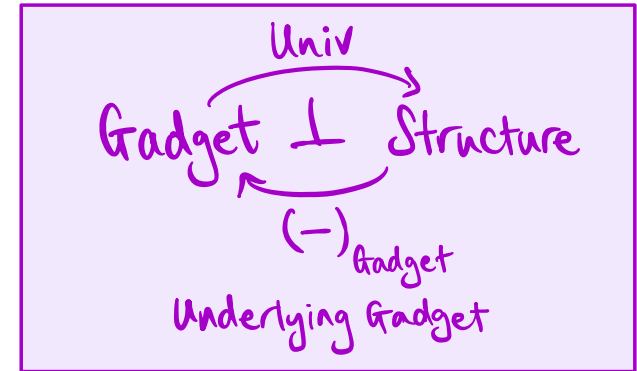
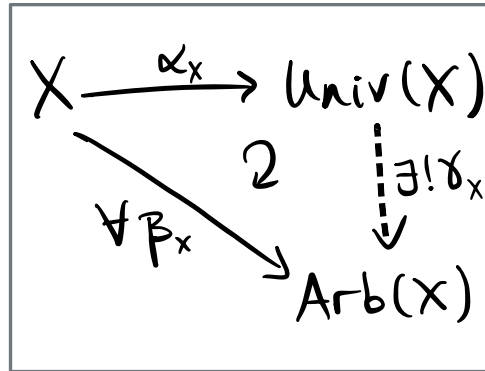
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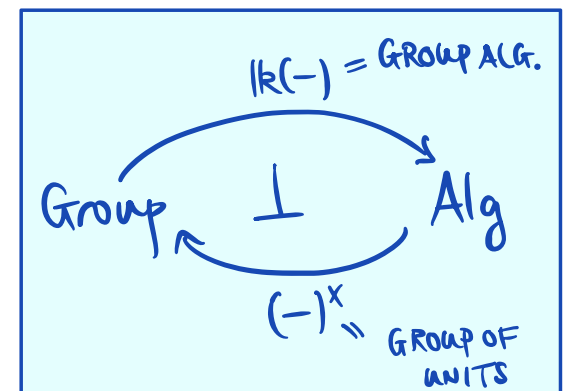
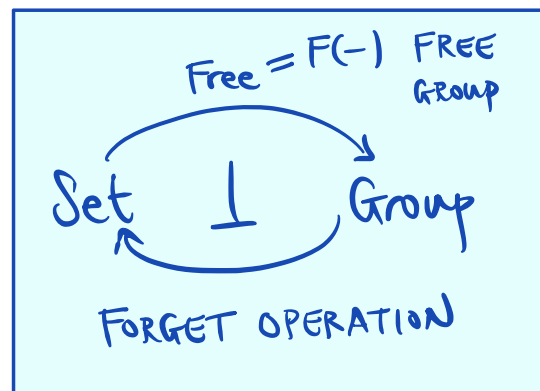
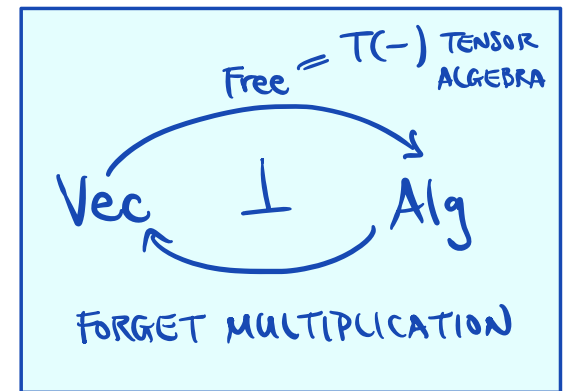
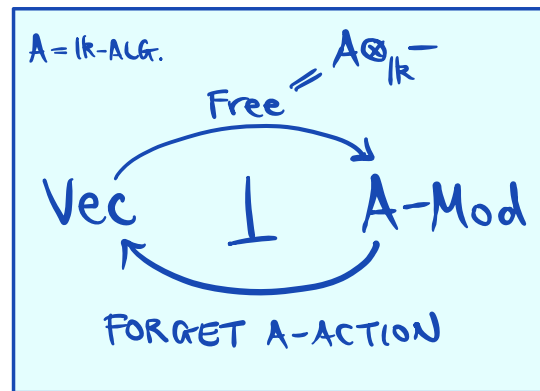
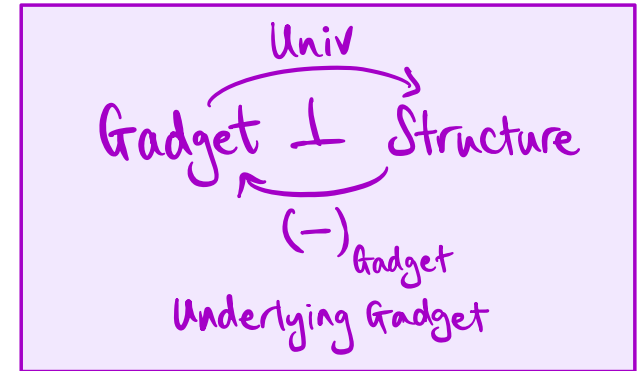
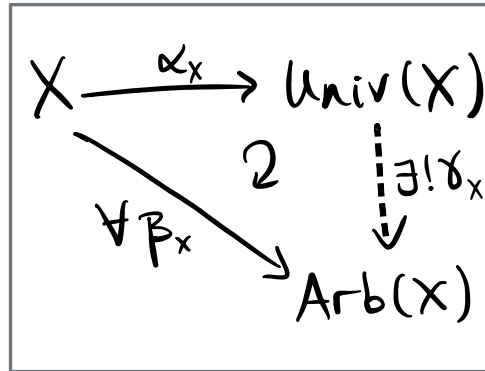
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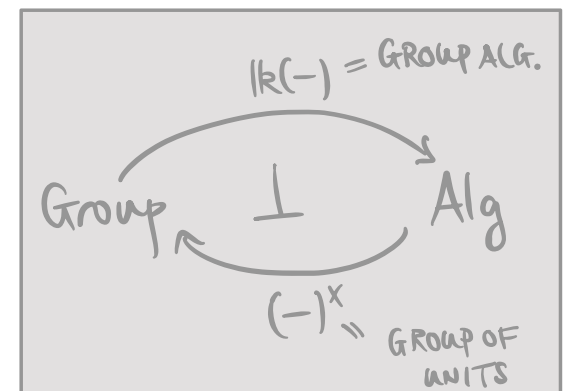
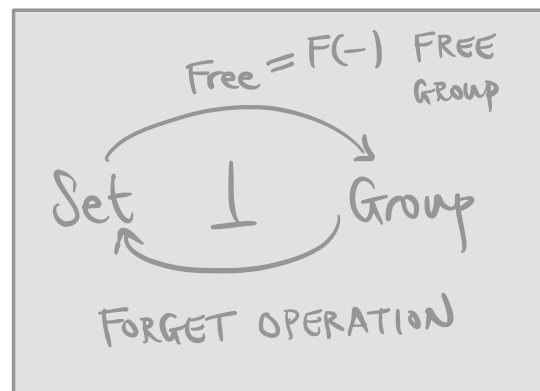
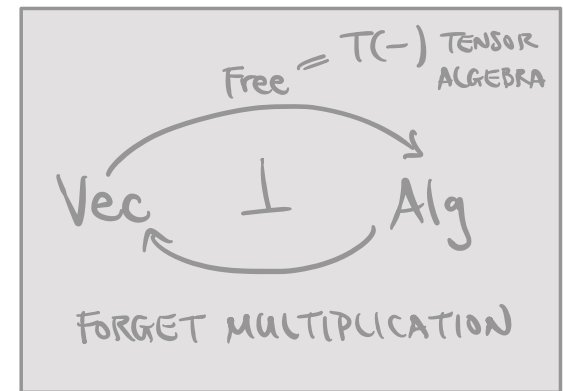
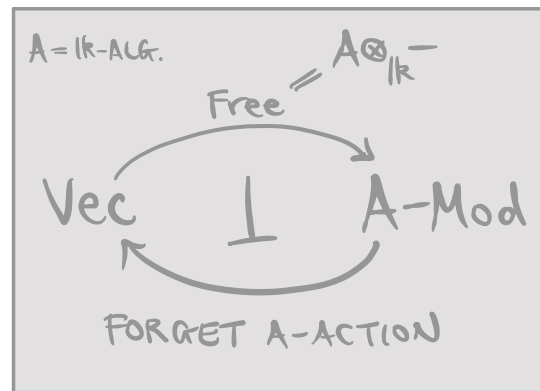
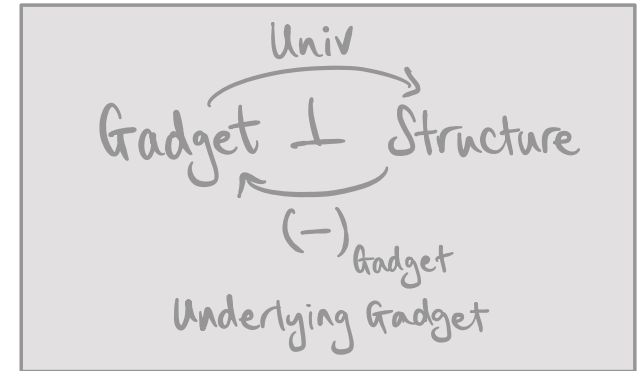
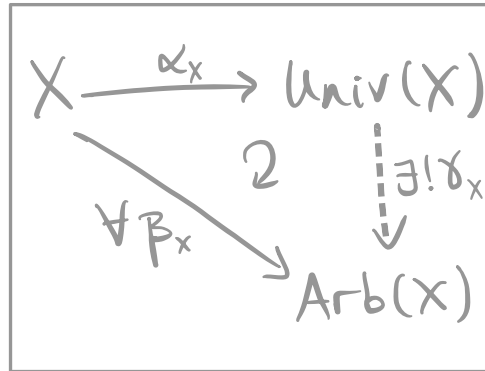
$H: \mathcal{D} \rightarrow \text{Set}$
COVARIANT

IS REPRESENTABLE IF
 $H(-) \cong \text{Hom}_{\mathcal{D}}(U, -)$

FOR SOME $U \in \mathcal{D}$
↑

UNIVERSAL OBJECT
THAT REPRESENTS H

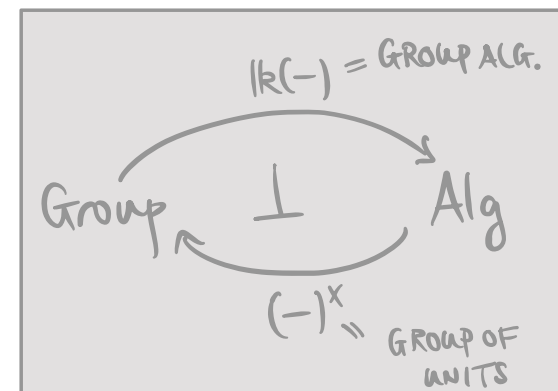
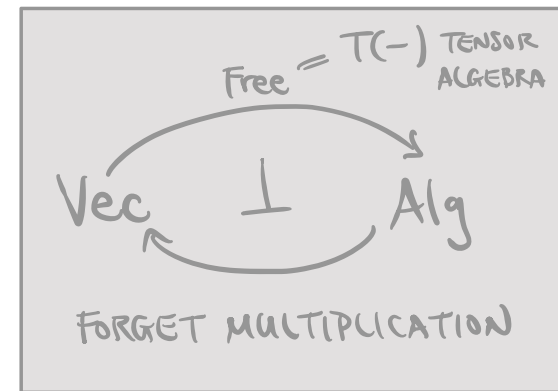
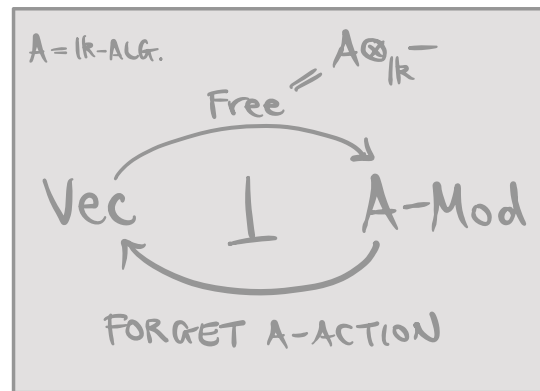
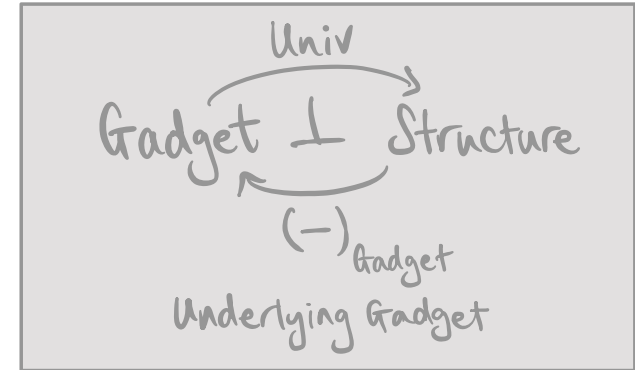
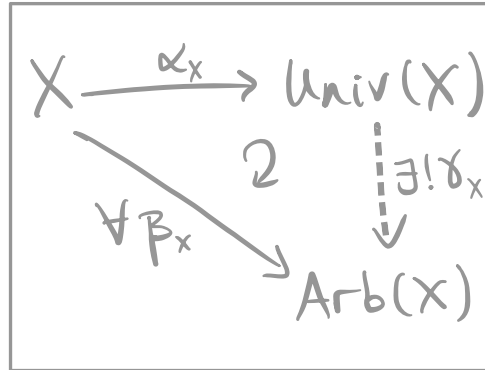
FORM I



II. UNIVERSALITY REVISITED

$H: \mathcal{D} \rightarrow \text{Set}$
 COVARIANT
 (RESP., CONTRAVARIANT)
 IS REPRESENTABLE IF
 $H(-) \cong \text{Hom}_{\mathcal{D}}(U, -)$
 (RESP.,
 $H(-) \cong \text{Hom}_{\mathcal{D}}(-, U)$)
 FOR SOME $U \in \mathcal{D}$
 \uparrow
 UNIVERSAL OBJECT
 THAT REPRESENTS H

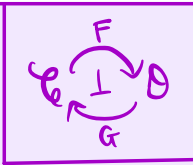
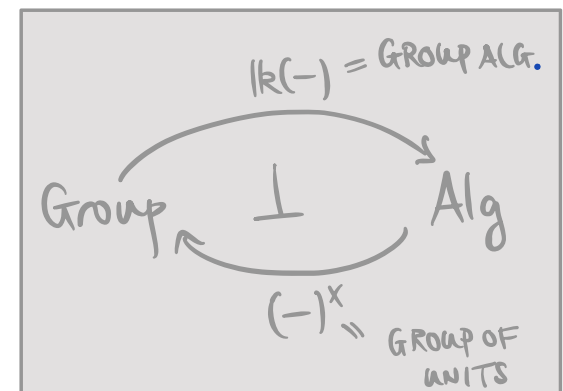
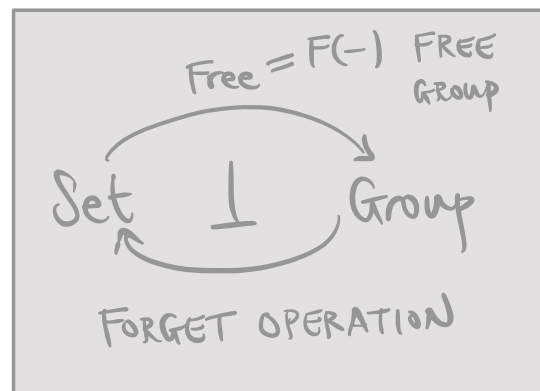
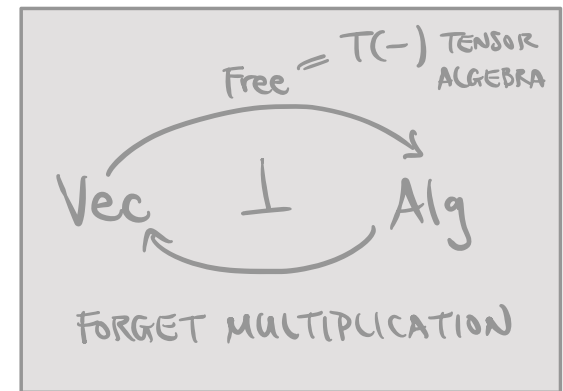
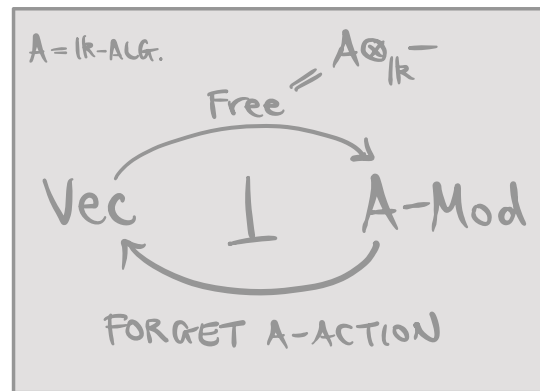
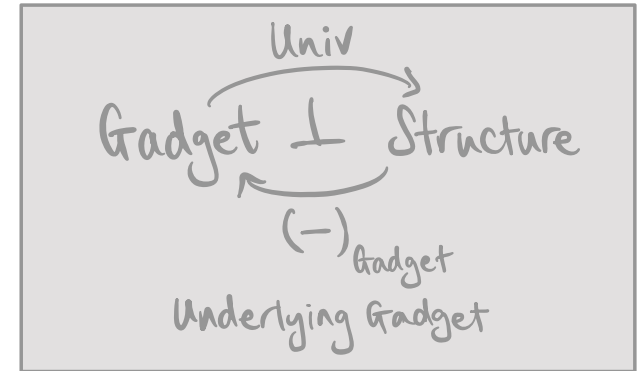
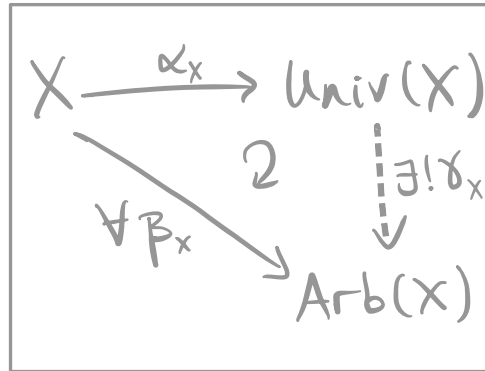
FORM I



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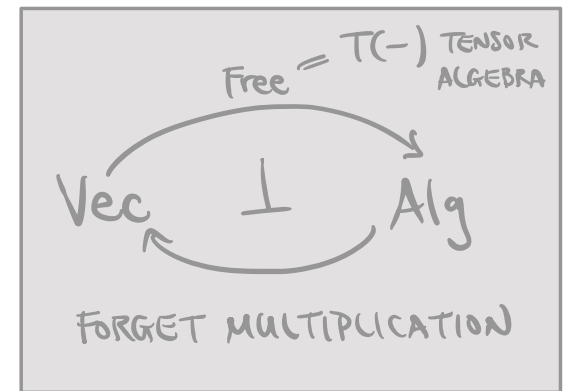
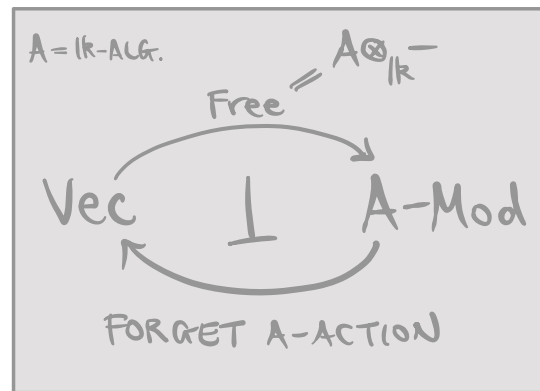
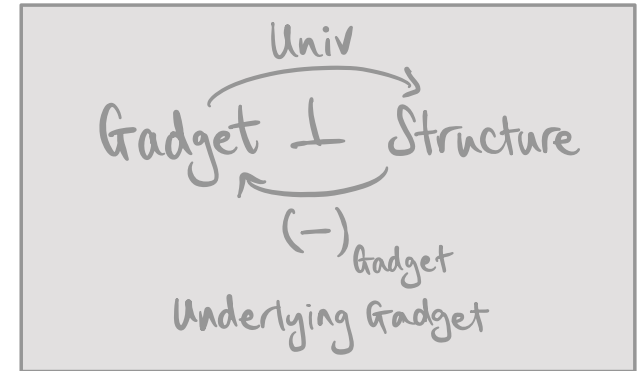
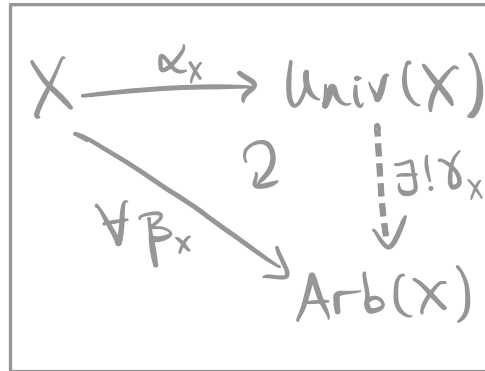
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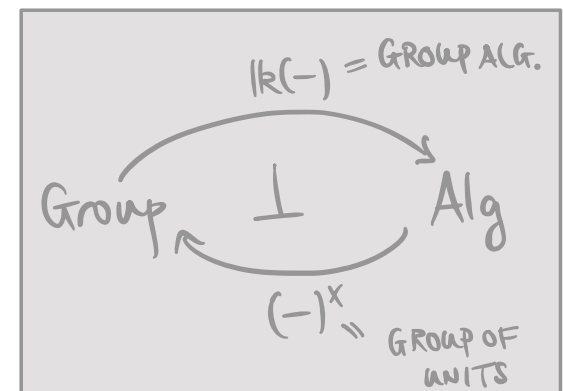
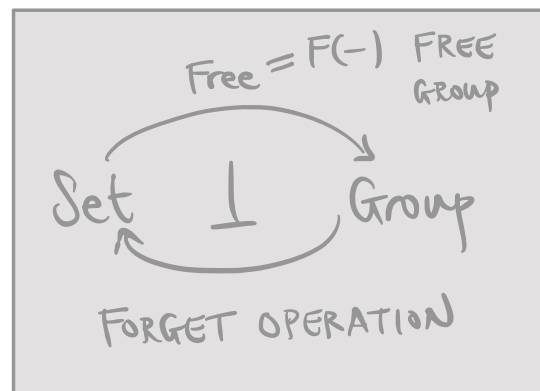
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FORM I



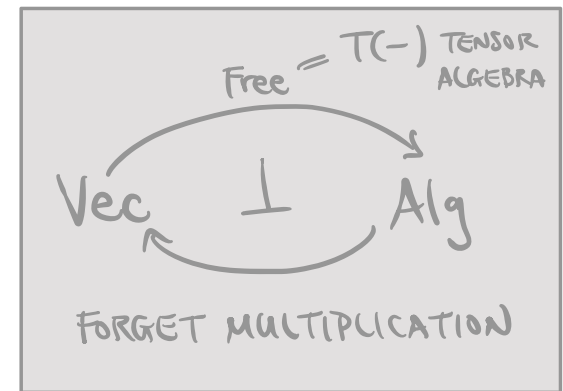
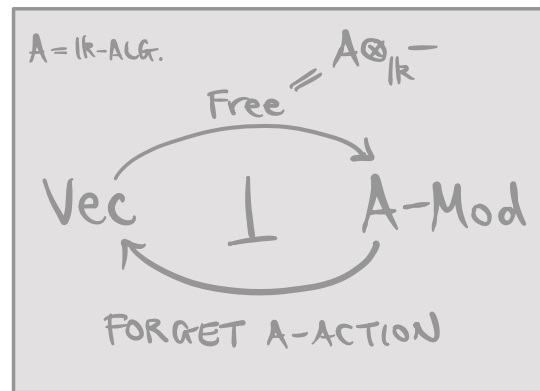
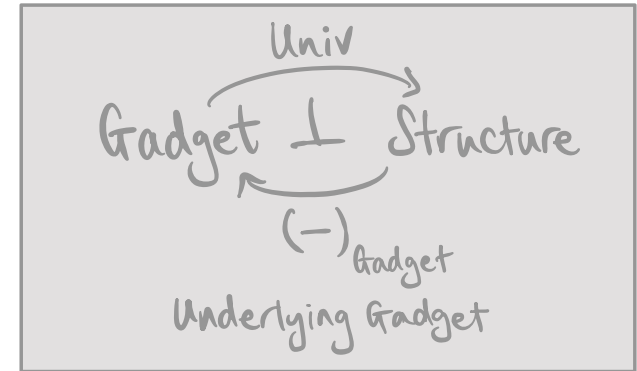
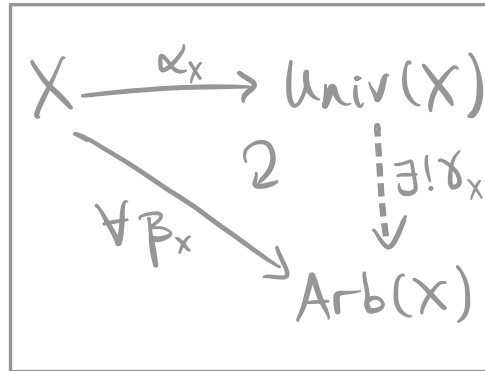
$\text{Hom}_{\mathcal{D}}(F(X), -)$
 s.t.
 $\text{Hom}_{\mathcal{D}}(X, G(-))$



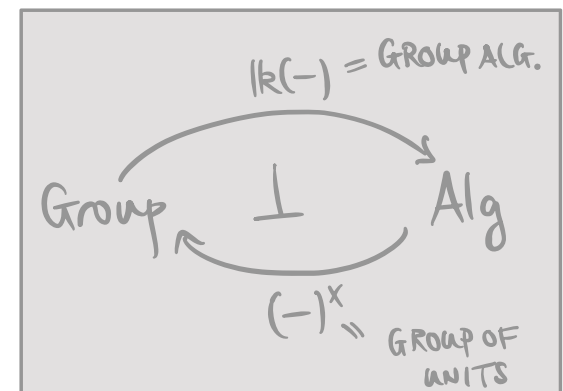
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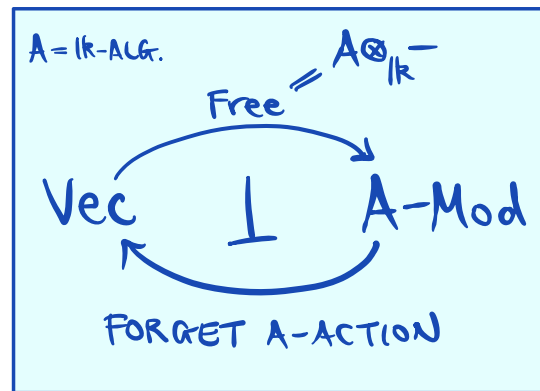
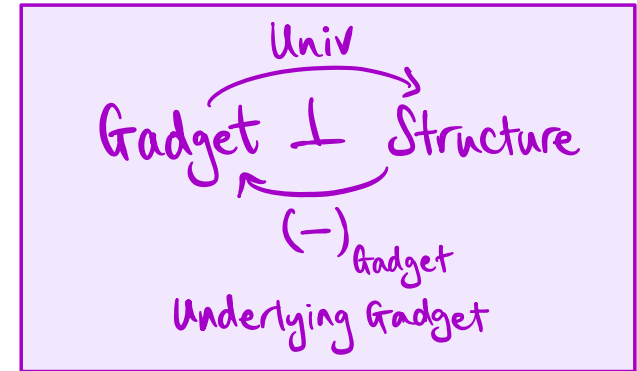
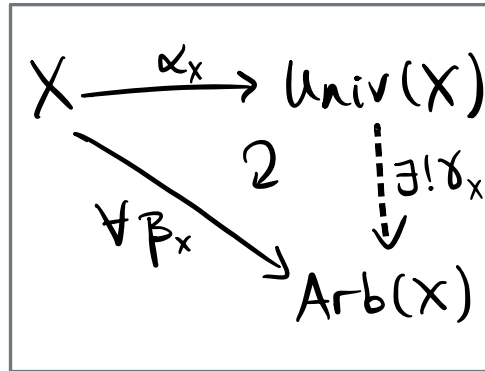
$\text{Hom}_{\mathcal{D}}(F(X), -)$
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 $\text{Hom}_{\mathcal{D}}(X, G(-)) : \mathcal{D} \rightarrow \text{Set}$
 $H_X \cong \text{Hom}_{\mathcal{D}}(X, G(-))$ REPRESENTED BY $F(X) \in \mathcal{D}$



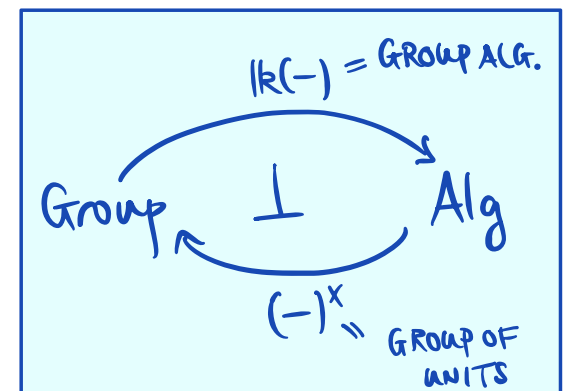
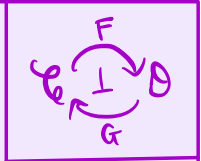
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FORM I

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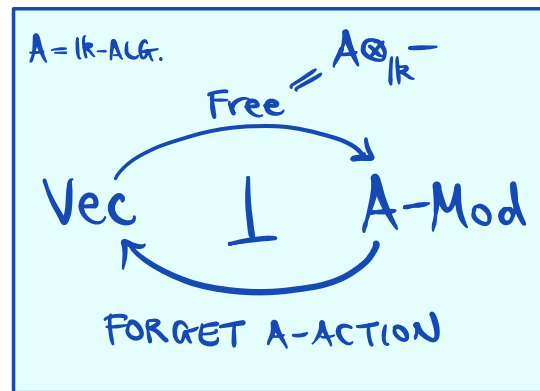
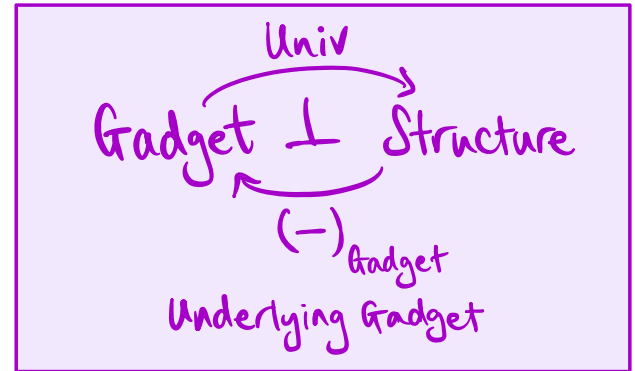
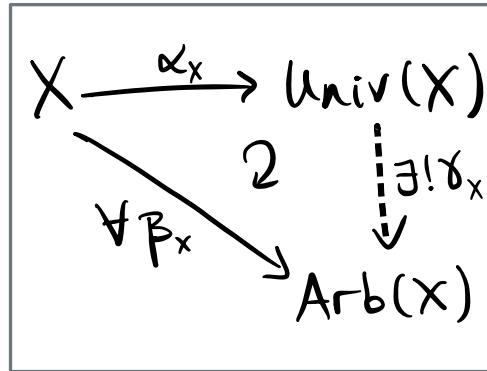
$\text{Hom}_{\mathcal{D}}(F(X), -)$
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II. UNIVERSALITY REVISITED

FORM I

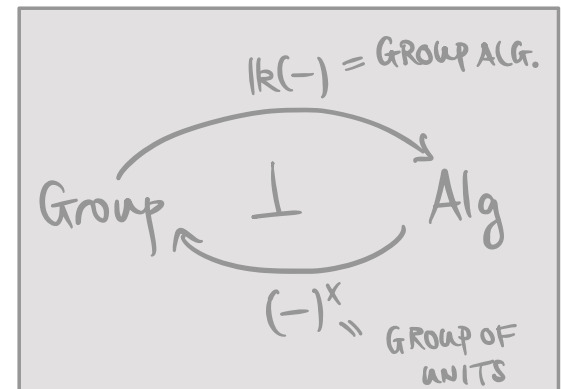
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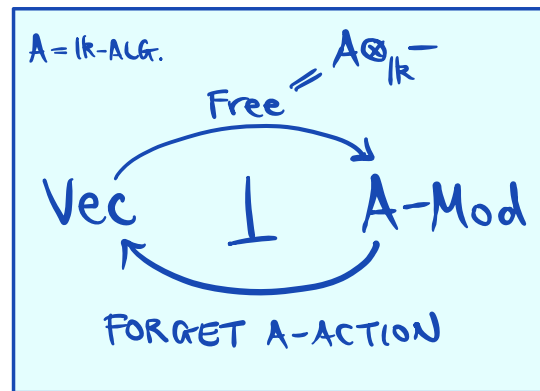
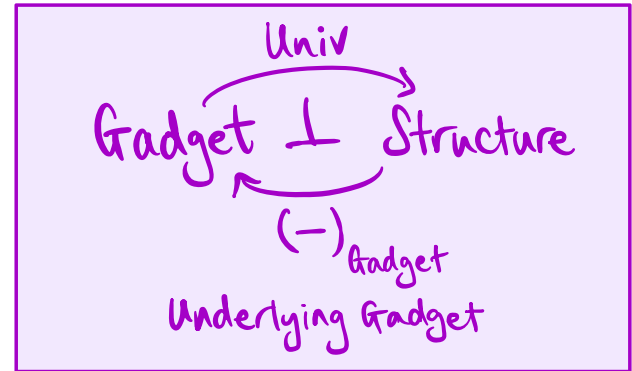
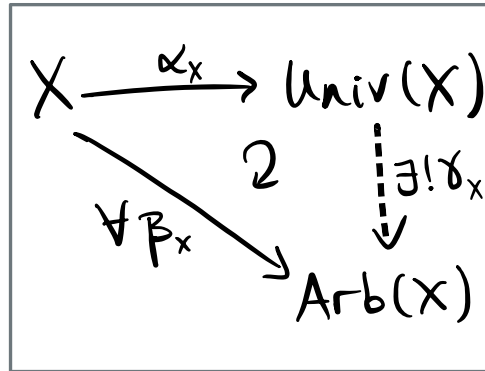
\downarrow
 $\text{Hom}_{A\text{-MOD}}(A \otimes_{\mathbb{K}} V, A W)$
 \cong
 $\text{Hom}_{\text{Vec}}(V, W)$



II. UNIVERSALITY REVISITED

FORM I

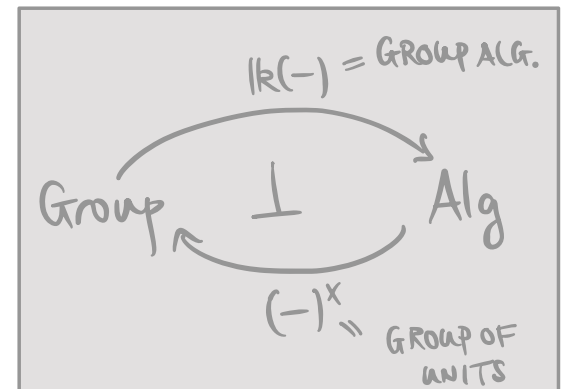
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 $\text{Hom}_{\mathcal{G}}(X, G(-)) : \mathcal{D} \rightarrow \text{Set}$
 H_X REPRESENTED BY $F(X) \in \mathcal{D}$



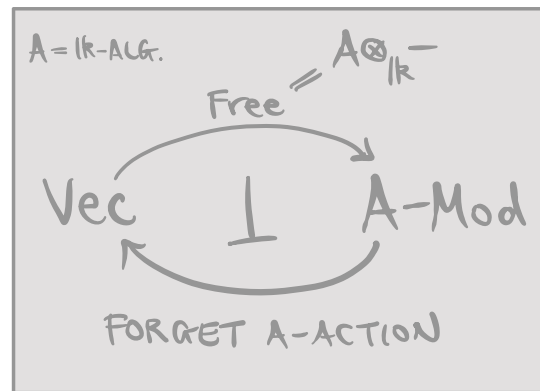
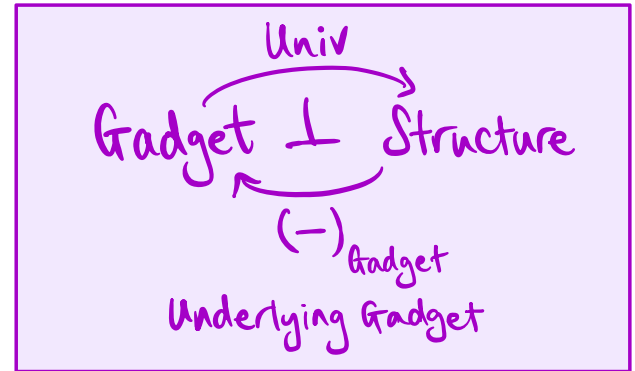
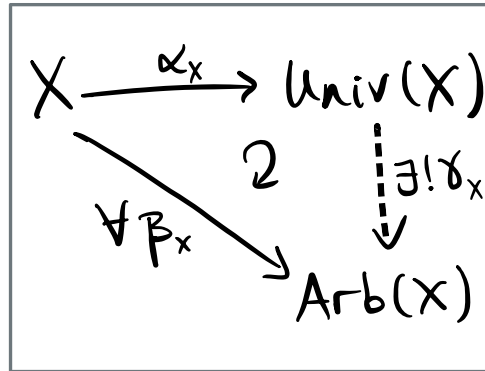
\downarrow
 $\text{Hom}_{A\text{-MOD}}(A \otimes_k V, A W)$
 \cong
 $\text{Hom}_{\text{Vec}}(V, W)$
 $H_V := \text{Hom}_{\text{Vec}}(V, \text{Forg}(-))$
 IS REPRESENTED BY $A \otimes_k V$



II. UNIVERSALITY REVISITED

FORM I

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 ↑
 UNIVERSAL OBJECT
 THAT REPRESENTS H



$$\text{Hom}_{\text{Alg}}(kG, A)$$

$$\cong \text{Hom}_{\text{Group}}(G, A^{\times})$$

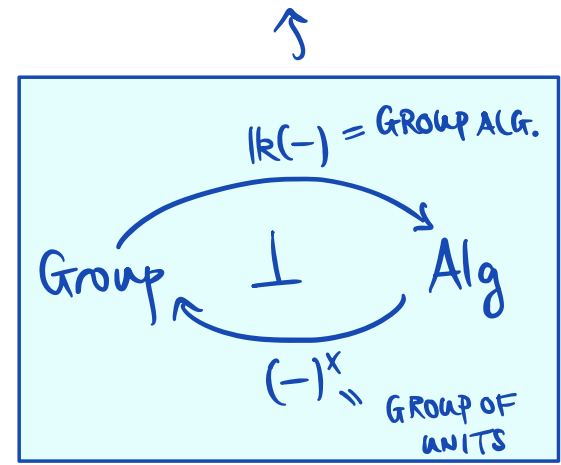
$\text{Hom}_{\mathcal{D}}(F(X), -)$
 $\cong \text{Hom}_{\mathcal{D}}(X, G(-)) : \mathcal{D} \rightarrow \text{Set}$
 H_X REPRESENTED BY $F(X) \in \mathcal{D}$



$$\text{Hom}_{A\text{-mod}}(A \otimes_k V, A W)$$

$$\cong \text{Hom}_{\text{Vec}}(V, W)$$

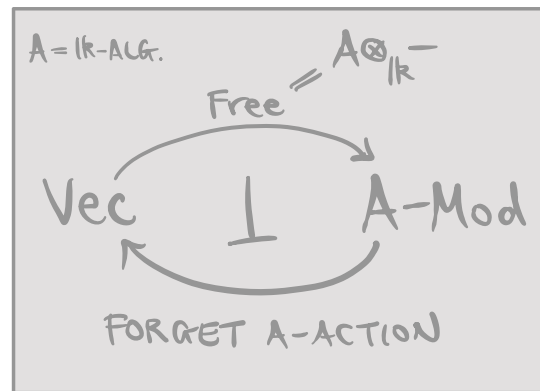
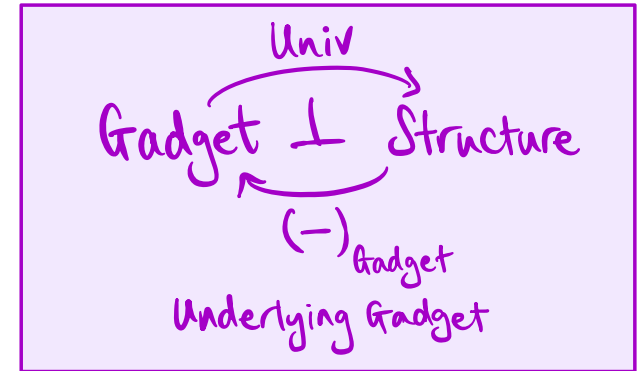
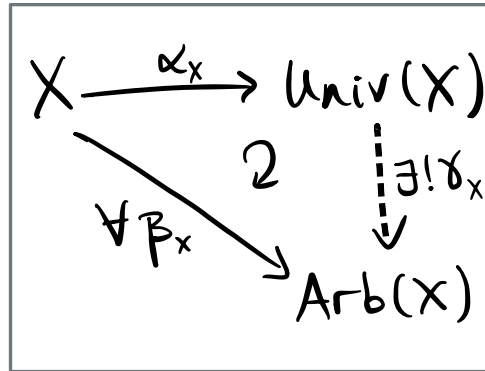
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FORM I

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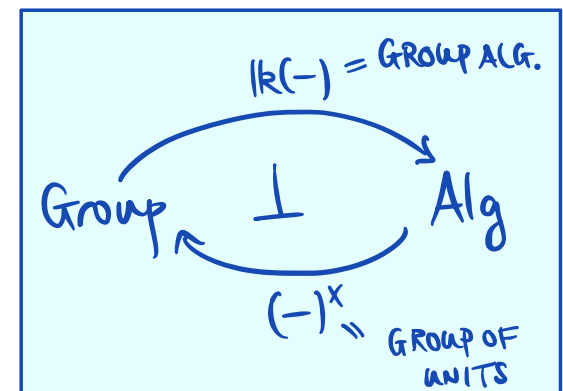


$\text{Hom}_{\text{Alg}}(\mathbb{K}G, A)$
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 $H_G := \text{Hom}_{\text{Group}}(G, (-)^{\times})$
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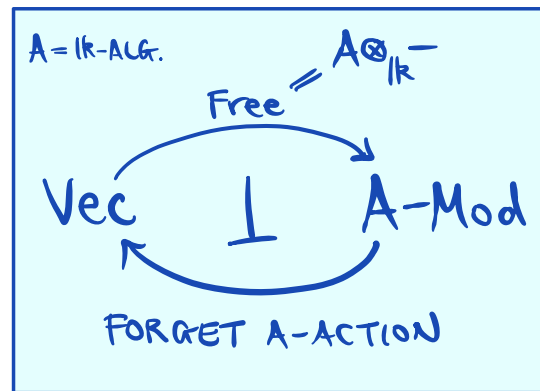
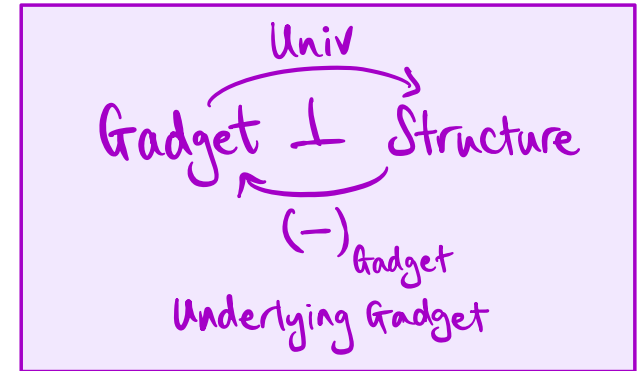
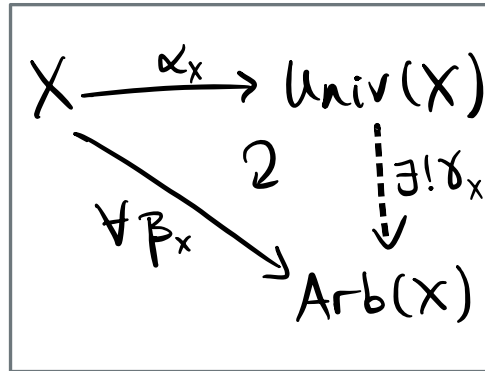
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 $\text{Hom}_{A\text{-mod}}(A \otimes_{\mathbb{K}} V, A W)$
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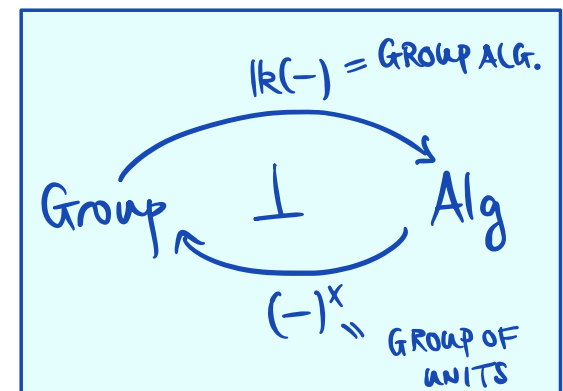
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III. YONEDA'S LEMMA



SOME CALL THIS THE IMPORTANT RESULT
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COVARIANT

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YONEDA'S LEMMA:

TAKE A FUNCTOR $F: \mathcal{D} \rightarrow \text{Set}$

THEN $\forall u \in \mathcal{D}$, \exists BIJECTION:

$$\Phi_{F,u}: \text{Nat}_{\mathcal{D}, \text{Set}}(\text{Hom}_{\mathcal{D}}(u, -), F) \xrightarrow{\sim} F(u).$$

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PF/ DEFINE $\mathcal{D} \begin{array}{c} \xrightarrow{\text{Hom}_{\mathcal{D}}(u, -)} \\ \Downarrow \varphi \\ \xrightarrow{F} \end{array} \text{Set} \xrightarrow{\Phi_{F,u}} \text{Set} \ni \bullet \in F(u).$

III. YONEDA'S LEMMA



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$$[\varphi_u: \text{Hom}_{\mathcal{D}}(u, u) \rightarrow F(u)]$$

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$[\varphi_u: \text{Hom}_{\mathcal{D}}(u, u) \rightarrow F(u)]$

DEFINE $\Psi_{F,u}: F(u) \longrightarrow \text{Nat}_{\mathcal{D}, \text{Set}}(\text{Hom}_{\mathcal{D}}(u, -), F)$

ω
 $x \longmapsto \mathcal{D} \begin{array}{c} \xrightarrow{\text{Hom}_{\mathcal{D}}(u, -)} \\ \Downarrow \Psi_{F,u}(x) \\ \text{Set} \\ \uparrow F \end{array}$

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 FOR SOME $u \in \mathcal{D}$
 \uparrow
 UNIVERSAL OBJECT
 THAT REPRESENTS H

YONEDA'S LEMMA:

TAKE A FUNCTOR $F: \mathcal{D} \rightarrow \text{Set}$
 THEN $\forall u \in \mathcal{D}$, \exists BIJECTION:

$$\Phi_{F,u}: \text{Nat}_{\mathcal{D}, \text{Set}}(\text{Hom}_{\mathcal{D}}(u, -), F) \xrightarrow{\sim} F(u).$$

PF/ DEFINE $\mathcal{D} \begin{array}{c} \xrightarrow{\text{Hom}_{\mathcal{D}}(u, -)} \\ \Downarrow \phi \\ \text{Set} \\ \uparrow F \end{array} \xrightarrow{\Phi_{F,u}} \phi_u(\text{id}_u) \in F(u).$

$[\phi_u: \text{Hom}_{\mathcal{D}}(u, u) \rightarrow F(u)]$

DEFINE $\Psi_{F,u}: F(u) \longrightarrow \text{Nat}_{\mathcal{D}, \text{Set}}(\text{Hom}_{\mathcal{D}}(u, -), F)$

ω
 $x \longmapsto \mathcal{D} \begin{array}{c} \xrightarrow{\text{Hom}_{\mathcal{D}}(u, -)} \\ \Downarrow \Psi_{F,u}(x) \\ \text{Set} \\ \uparrow F \end{array}$
 WHERE
 $(\Psi_{F,u}(x))_z: \text{Hom}_{\mathcal{D}}(u, z) \rightarrow F(z)$
 $f \longmapsto F(f)(x)$

III. YONEDA'S LEMMA



SOME CALL THIS THE IMPORTANT RESULT
IN CATEGORY THEORY

$H: \mathcal{D} \rightarrow \text{Set}$
 COVARIANT
 (RESP., CONTRAVARIANT)
 IS REPRESENTABLE IF
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$[\varphi_u: \text{Hom}_{\mathcal{D}}(u, u) \rightarrow F(u)]$

$\left. \begin{array}{l} \swarrow \\ \searrow \end{array} \right\} \text{CHECK THESE ARE MUTUALLY INVERSE} //$

DEFINE $\Psi_{F,u}: F(u) \longrightarrow \text{Nat}_{\mathcal{D}, \text{Set}}(\text{Hom}_{\mathcal{D}}(u, -), F)$

$x \mapsto \mathcal{D} \begin{array}{c} \xrightarrow{\text{Hom}_{\mathcal{D}}(u, -)} \\ \Downarrow \Psi_{F,u}(x) \\ \text{Set} \\ \uparrow F \end{array}$
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III. YONEDA'S LEMMA

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MAIN CONSEQUENCE:

REPRESENTABLE FUNCTORS ARE
DETERMINED BY THEIR UNIV. OBJECTS

III. YONEDA'S LEMMA

$H: \mathcal{D} \rightarrow \text{Set}$
COVARIANT

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REPRESENTABLE FUNCTORS ARE
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COROLLARY:

IN SET

$$\text{Hom}_{\mathcal{D}}(u, -) \cong \text{Hom}_{\mathcal{D}}(u', -) \Leftrightarrow u \cong u'.$$

$$\text{Hom}_{\mathcal{D}}(-, u) \cong \text{Hom}_{\mathcal{D}}(-, u') \Leftrightarrow u \cong u'.$$

IN \mathcal{D}

III. YONEDA'S LEMMA



SOME CALL THIS THE IMPORTANT RESULT
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$H: \mathcal{D} \rightarrow \text{Set}$
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MAIN CONSEQUENCE:

REPRESENTABLE FUNCTORS ARE
DETERMINED BY THEIR UNIV. OBJECTS

PROOF \equiv
EXERCISE 2.48

COROLLARY:

IN SET IN \mathcal{D}
 $\text{Hom}_{\mathcal{D}}(u, -) \cong \text{Hom}_{\mathcal{D}}(u', -) \iff u \cong u'.$
 $\text{Hom}_{\mathcal{D}}(-, u) \cong \text{Hom}_{\mathcal{D}}(-, u') \iff u \cong u'.$

MATH 466/566
SPRING 2024

CHELSEA WALTON
RICE U.

LECTURE #10

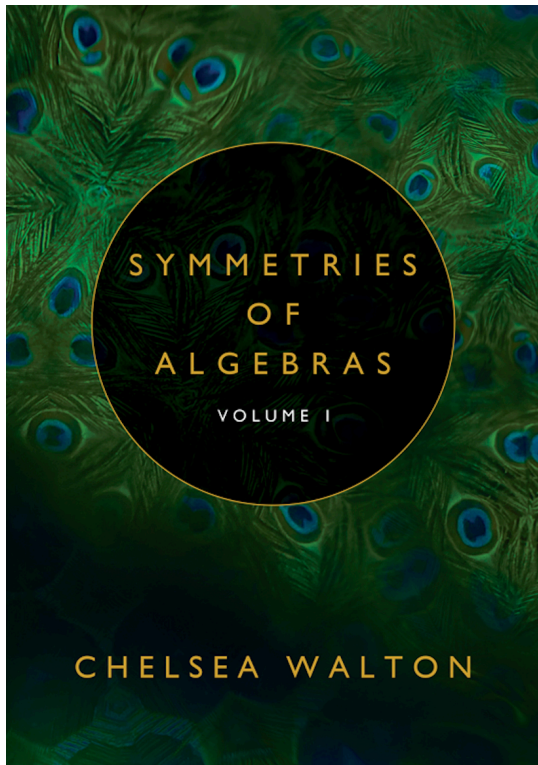
TOPICS:

- I. ADJUNCTION (§2.5)
- II. UNIVERSALITY REVISITED (§2.6.1)
- III. YONEDA'S LEMMA (§2.6.2)

NEXT TIME: BITS OF HOMOLOGICAL ALGEBRA

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You'll enjoy the textbook!**

C. Walton's "Symmetries of Algebras, Volume 1" (2024)



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Lecture #10 keywords: adjoint functors, adjunction, Free-Forgetful adjunction, representable functor, Tensor-Hom adjunction, universal object, Yoneda's Lemma