### MATH 466/566 SPRING 2024

CHELSEA WALTON RICE U.

#### LAST TIME

· (SOM. AND EQUIV. OF MONOIDAL CATEGORIES LECTURE #14

- · MODULE CATEGORIES
- · BIMODULE CATEGORIES

## TOPICS:

I.	STRICTIFICATION	(§3.4.1)
II.	COHERENCE	(83.4.3)
III.	GRAPHICAL CALCULUS	(§3.5)

IV. RIGID CATEGORIES

(\$3.6)

= RECALL=

A MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR ⊗: & × & → & (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

= RECALL =

A MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR  $\otimes$ :  $\mathcal{C} \times \mathcal{C} \longrightarrow \mathcal{C}$ (MONOIDAL PRODUCT)
- (c) AN OBJECT 1 ∈ 6 (MONOIDAL UNIT)

(d,e,f) NATURAL L80MORPHISMS:

$$0 = \begin{cases} \alpha_{x,y,z} : (x \otimes y) \otimes \xi & \text{NATURAL IN } x_{,y,z} \in \mathcal{X} \\ \Rightarrow x \otimes (y \otimes \xi) & \\ x_{,y,z} \in \mathcal{X} \end{cases}$$
(ASSOCIATIVITY CONSTRAINT)

$$1 = \{l_X : 10 X \xrightarrow{\sim} X \}_{X \in \mathcal{C}} \text{ NATURAL IN } X$$
(LEFT UNITALITY CONSTRAINT)

SATISFYING THE COMPATIBILITY
(ONDITIONS BELOW:

= RECALL =

### A MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × & ---> & (MONOIDAL PRODUCT)
- (c) AN OBJECT 1 ∈ 6 (MONOIDAL UNIT)

(d,e,f) NATURAL (SOMORPHISMS:

$$0 = \begin{cases} \alpha_{x,y|z} : (x \otimes y) \otimes z \\ \longrightarrow x \otimes (y \otimes z) \end{cases}$$
NATURAL IN  $x,y,z \in \mathcal{X}$ 

$$(ASSOCIATIVITY CONSTRAINT)$$

 $1 = \{l_X : 1 \otimes X \xrightarrow{\sim} X \}_{X \in \mathcal{C}} \text{ NATURAL IN } X$ (LEFT UNITALITY CONSTRAINT)

right unitality constraint)

SATISFYING THE COMPATIBILITY
(ONDITIONS BELOW:

YWIXIYIZEC:

$$((W \otimes X) \otimes Y) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes$$

= RECALL =

### A MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & X & --- & (MONOIDAL PRODUCT)
- (c) AN OBJECT 1 ∈ 6 (MONOIDAL UNIT)

(d,e,f) NATURAL (SOMORPHISMS:

$$0 = \begin{cases} \alpha_{x,y,z} : (x \otimes y) \otimes \xi \\ \Rightarrow x \otimes (y \otimes \xi) \end{cases}$$

$$\text{NATURAL IN } x,y,\xi \in \mathcal{X}$$

$$(\text{ASSOCIATIVITY CONSTRAINT})$$

 $J = \{l_X : 1 \otimes X \xrightarrow{\sim} X \}_{X \in \mathcal{X}}$  NATURAL IN X (LEFT UNITALITY CONSTRAINT)

right unitality constraint)

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

¥W,X,Y, ₹ € €:

$$((W \otimes X) \otimes Y) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes X) \otimes Z)$$

$$(W \otimes X)$$

## A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR  $\otimes$ :  $& \times & \longrightarrow & \\ \text{(MONOIDAL PRODUCT)}$
- (c) AN OBJECT 1 ∈ 6 (MONOIDAL UNIT)
- (d,e,f) NATURAL (SOMORPHISMS:

$$\alpha = \left\{ \begin{array}{l} \alpha_{X,Y,1} \in (X \otimes Y) \otimes \xi \\ \Rightarrow X \otimes (Y \otimes \xi) \end{array} \right\}_{X,Y,1} \in \mathcal{X}$$

$$(ASSOCIATIVITY CONSTRAINT)$$

- $1 = \{l_X : 1 \otimes X \xrightarrow{\sim} X\}_{X \in \mathcal{X}} \text{ NATURAL IN } X$ (LEFT UNITALITY CONSTRAINT)
- right unitality constraint)

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

AMIXIAISEG:

$$((W \otimes X) \otimes Y) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes X) \otimes Z)$$

$$(W \otimes X) \otimes Z$$

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × & → & (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) NATURAL LEOMORPHISMS:

 $\alpha = \left\{ \begin{array}{l} \alpha_{X,Y,1} \neq : (X \otimes Y) \otimes \xi \\ & \stackrel{\sim}{\longrightarrow} X \otimes (Y \otimes \xi) \end{array} \right\}_{X,Y,1} \in \mathcal{X}$  (ASSOCIATIVITY CONSTRAINT)

 $1 = \{l_X : 1 \otimes X \xrightarrow{\sim} X\}_{X \in \mathcal{X}}$  NATURAL IN X (LEFT UNITALITY CONSTRAINT)

right unitality constraint)

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

4M1×1115€6:

$$((W \otimes X) \otimes Y) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes X) \otimes Z)$$

$$(W \otimes X) \otimes Z$$

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & x & → & (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) - NATURAL (SOMORPHISMS: idx,y,it) IDENTITIES $<math display="block">0 = \begin{cases} a_{x,y,t} : (x \otimes y) \otimes t \\ \Rightarrow x \otimes (y \otimes t) \end{cases}$ NATURAL (N X,Y,t) (ASSOCIATIVITY CONSTRAINT)

L= {Lx: 10 X ~ X Jxer NATURAL IN X

(LEFT UNITALITY CONSTRAINT)

= idx

= {Tx: X \omega 1 ~ X Jxer NATURAL IN X

(RIGHT UNITALITY CONSTRAINT)

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

AMIXINISEC:

$$(W \otimes X) \otimes Y \otimes Z$$

$$(W \otimes X) \otimes Y \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes X)$$

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × € → € (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) -NATURAL (80 MORPHISMS:

$$0 = \begin{cases} idx,y,z & IDENTITIES \\ ax,y,z : (x \otimes y) \otimes z & S \\ \Rightarrow x \otimes (y \otimes z) \end{cases}$$

$$(x \otimes y) \otimes z = x \otimes (y \otimes z)$$

$$(x \otimes y) \otimes z = x \otimes (y \otimes z)$$

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

: 33 51/1×1/WK

$$((W \otimes X) \otimes Y) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes$$

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × € → € (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) -NATURAL (SOMORPHISMS:

$$0 = \begin{cases} a_{X,Y|Z} & \text{IDENTITIES} \\ a_{X,Y|Z} : (X \otimes Y) \otimes Z & \text{NATURAL IN } X,Y,Z \\ \Rightarrow X \otimes (Y \otimes Z) \end{cases} \times_{X,Y,Z \in \mathcal{X}} \times_{X,Y,Z \in \mathcal{X}} \times_{X,Y,Z \in \mathcal{X}} \times_{X,Z} \times_{X,Z} \times_{X,Z} \times_{X,Z} \times_{X,Z} \times_{Z} \times_{X,Z} \times_{Z} \times_{Z$$

 $I = \{l_X : 1L \otimes X \xrightarrow{\sim} X\}_{X \in \mathcal{U}} \text{ NATURAL (N X X )}$   $(\text{LEFT UNITALITY CONSTRAINT}) \qquad \text{II}$   $= id_X \qquad \qquad \text{LOX}$   $\Gamma = \{\Gamma_X : X \otimes L \xrightarrow{\sim} X\}_{X \in \mathcal{U}} \text{ NATURAL (N X II)}$ 

(RIGHT UNITALITY CONSTRAINT) XOL

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

: 33 51/1×1/WK

$$((W \otimes X) \otimes Y) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y)) \otimes Z$$

$$(W \otimes (X \otimes Y) \otimes Z)$$

$$(W \otimes$$

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × € → € (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) -NATURAL (SOMORPHISMS:

 $0 = \begin{cases} idx,y,z & |DENTITIES \\ ax,y,z & (x \otimes y) \otimes z \end{cases}$   $x \otimes y \otimes z := (x \otimes y) \otimes z = x \otimes (y \otimes z)$   $x \otimes y \otimes z := (x \otimes y) \otimes z = x \otimes (y \otimes z)$ 

L= [Lx: L\omega X] xer NATURAL IN X X

(LEFT UNITALITY CONSTRAINT)

idx

[\omega \text{L} \times X] \times NATURAL IN X II

(RIGHT UNITALITY CONSTRAINT)

X\omega L

SATISFYING THE COMPATIBILITY
(ONDITIONS BELOW:

4M1×1115€6:

 $M \otimes (X \otimes \lambda) \otimes 5) \xrightarrow{\hspace{1cm}} M \otimes (X \otimes (\lambda \otimes 5))$   $Om^{1} \times^{2} \lambda^{1} \in I$   $Om^{1} \times^{2} \lambda^{1} \in I$   $Om^{1} \times^{2} \lambda^{1} \otimes f$   $Om^{1} \times^{2} \lambda^{1} \otimes f$ 

(PENTAGON AXIOM) VACUOUS

NO PARENTHESIS NEEDED

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × & → & (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) -NATURAL (80 MORPHISMS:

 $\alpha = \begin{cases} a_{X,Y,1}z & \text{id}_{X,Y,1}z \\ & \Rightarrow \chi_{0}(y\otimes z) \end{cases}$   $\text{NATURAL IN } \chi_{1,Y,1}z \\ \xrightarrow{\sim} \chi_{0}(y\otimes z) \end{cases} \times_{1} \chi_{1}z \in \mathcal{U}$ 

(504)0X = 20(10X) =: 2010X

1= [lx: 10x ~ X] Xer NATURAL IN X X (LEFT UNITALITY CONSTRAINT)

= idx

[ = {Tx: X \omega 1 \simplies \chi 3 \chi 2 (RIGHT UNITALITY CONSTRAINT) XOL SATISFYING THE COMPATIBILITY CONDITIONS RELOW:

AMIXINIEG:

 $((M \otimes X) \otimes A) \otimes 5$ (M & (X & Y)) & 5 (M & X) & (A & 5) UMIXWAIS II OMIX'ABS  $\mathsf{M}\otimes(\mathsf{X}\otimes\mathsf{Y})\otimes\mathsf{S})\longrightarrow \mathsf{M}\otimes(\mathsf{X}\otimes(\mathsf{Y}\otimes\mathsf{S}))$ 

(PENTAGON AXIOM) VACUOUS -> NO PARENTHESIS NEEDED

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR Ø: & × € → €
  (MONOIDAL PRODUCT)
- (c) AN OBJECT 11 € 6 (MONOIDAL UNIT)

(d,e,f) - NATURAL (80 MORPH(3 MS: idx,y,it)) = idx,y,it = idx,y,

SATISFYING THE COMPATIBILITY
CONDITIONS BELOW:

AMIXINIEG:

$$M \otimes (X \otimes \lambda) \otimes Y = A \otimes (X \otimes \lambda) \otimes A \otimes (X \otimes \lambda$$

PENTAGON AXIOM) VACUOUS

NO PARENTHESIS NEEDED

A , MONOIDAL CATEGORY CONSISTS OF:

- (a) A CATEGORY &
- (b) A BIFUNCTOR ⊗: & × & → & (MONOIDAL PRODUCT)
- (c) AN OBJECT 1 ∈ C (MONOIDAL UNIT)

COMPUTATIONS ARE SIMPLIFIED

SATISFYING THE COMPATIBILITY
(ONDITIONS BELOW:

AMIXIAISEG:

$$M \otimes (X \otimes \lambda) \otimes S) \stackrel{=}{=} M \otimes (X \otimes \lambda \otimes S)$$

$$Om^{1/2} \int_{\mathbb{R}^{3}} I \int_{\mathbb{R}^{3}} Om^{1/2} \int_{\mathbb{R}^{3}} Om^{1/$$

(PENTAGON AXIOM) VACUOUS

NO PARENTHESIS NEEDED

TAKE MONOIDAL CATEGORIES

& := (e, &, 11, 2, 12, 2) \$ 0 := (0, 8, 11, 2, 12, 2).

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> 8.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \langle F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y)^{0} \times_{y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): 118 -> F(114) IN B.

SATISFYING:

COMPUTATIONS
ARE SIMPLIFIED

FOR INSTANCE RECALL

TAKE MONOIDAL CATEGORIES

A MONOIDAL FUNCTOR FROM & TO B CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> 8.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F^{(2)}_{X,Y} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y) \}_{X,Y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): 118 -> F(118) IN B.

SATISFYING: 
$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F(z) \otimes id} F(X \otimes^{e} Y) \otimes^{b} F(z) \xrightarrow{F(z) \otimes id} F((X \otimes^{e} Y) \otimes^{e} z)$$

$$(F(X) \otimes^{b} (F(Y)) \otimes^{b} F(z)) \xrightarrow{id \otimes F(z)} F(X \otimes^{e} Y) \otimes^{b} F(X \otimes^{e} Y) \otimes^{e} z)$$

$$F(X) \otimes^{b} (F(Y) \otimes^{b} F(z)) \xrightarrow{id \otimes F(z)} F(X \otimes^{e} Y) \otimes^{e} F(X \otimes^{e} Y) \otimes^{e} z)$$

$$F(x) \otimes^{b} 1^{b} \xrightarrow{\Gamma(x)} F(x)$$

$$id \otimes F^{(o)} \downarrow \qquad \qquad F(x)$$

$$F(x) \otimes^{b} \Gamma(x) \xrightarrow{F(x)} F(x)$$

COMPUTATIONS

TAKE MONOIDAL CATEGORIES ع := (ع , ه الر م الر

WHEN STRICT COMPUTATIONS

ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> 8.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{X,Y}^{(2)} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y) \}_{X,Y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(118) IN B.

SATISFYING: 
$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F(z)} F(X) \otimes^{b} F(y) \otimes^{b} F(z) \xrightarrow{F(z)} F(x) \otimes^{b} F(y) \otimes^{c} z)$$

$$(F(X) \otimes^{b} (F(Y)) \otimes^{b} F(z)) \xrightarrow{id \otimes F(z)} F(x) \otimes^{c} F(y) \otimes^{c} F($$

$$\begin{array}{ccc}
\mathbb{L}^{0} \otimes^{0} F(X) & \xrightarrow{\mathbb{L}^{0}} & F(X) \\
F^{(0)} \otimes id & & & & & \\
& & & & & \\
F(\mathbb{L}^{e}) \otimes^{0} F(X) & \xrightarrow{\mathbb{L}^{(2)}} & F(\mathbb{L} \otimes^{e} X)
\end{array}$$

WHEN STRICT TAKE MONOIDAL CATEGORIES & := (&, &, 1, a, 1, r) \$ 0 := (0, 8, 1, a, 1, r)

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: " -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F^{(2)}_{X,Y} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y) \}_{X,Y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): 118 -> F(118) IN B.

$$\begin{array}{ccc}
\mathbb{L}^{\delta} \otimes^{\delta} F(X) & \xrightarrow{\mathbb{L}^{\delta}} & F(X) \\
F^{(\bullet)} \otimes id & & & & & & \\
\mathbb{E}^{(\bullet)} \otimes id & & & & & & \\
\mathbb{E}^{(\bullet)} \otimes id & & & & & & \\
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\mathbb{E}^{(\bullet)} \otimes id & & \\
\mathbb{E}^{(\bullet)} \otimes$$

COMPUTATIONS

ARE SIMPLIFIED

WHEN STRICT TAKE MONOIDAL CATEGORIES C:= (e, &, 1, a, 1, r) \$ 0:= (0, 8, 1, a, 1, r)

COMPUTATIONS ARE SIMPLIFIE

BECOMES

A SQUARE

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F^{(2)}_{X,Y} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y)^{0} \times_{Y} \in \mathcal{C}.$

(c) A MORPHISM F(0): UB -> F(UC) IN B.

SATISFYING: 
$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F(z)} (X \otimes^{e} Y) \otimes^{b} F(z) \xrightarrow{F(x)} F((X \otimes^{e} Y) \otimes^{e} z)$$

$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F(z)} F(X \otimes^{e} Y) \otimes^{b} F(z) \xrightarrow{F(z)} F((X \otimes^{e} Y) \otimes^{e} z)$$

$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F(x)} F(X \otimes^{e} Y) \otimes^{e} F(z) \xrightarrow{F(x)} F(X \otimes^{e} Y) \otimes^{e} z)$$

$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{f(x)} F(X \otimes^{e} Y) \otimes^{e} F(z) \xrightarrow{F(x)} F(X \otimes^{e} Y) \otimes^{e} z)$$

$$(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{f(x)} F(X \otimes^{e} Y) \otimes^{e} F(z) \xrightarrow{F(x)} F(X \otimes^{e} Y) \otimes^{e} z)$$

$$(F(X) \otimes^{b} F(Y)) \otimes^{e} F(z) \xrightarrow{f(x)} F(X \otimes^{e} Y) \otimes^{e} F(z) \xrightarrow{F(x)} F(X \otimes^{e} Y) \otimes^{e} Z \xrightarrow{F(x)} Z \xrightarrow{F(x)} F(X \otimes^{e} Y) \otimes^{e} Z \xrightarrow{F(x)} Z \xrightarrow{F$$

WHEN STRICT TAKE MONOIDAL CATEGORIES & := (&, &, 1, a, 1, r) \$ 0 := (0, 8, 1, a, 1, r)

COMPUTATIONS ARE SIMPLIFIE

BECOMES

A SQUARE

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F^{(2)}_{X,Y} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y)^{0} \times_{Y} \in \mathcal{C}.$

(c) A MORPHISM F(0): UB -> F(UC) IN B.

SATISFYING: 
$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{e} Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{e} Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{e} Y) \otimes^{e} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{e} Y) \otimes^{e} F($$

WHEN STRICT TAKE MONOIDAL CATEGORIES C:= (e, &, 1, a, 1, r) \$ 0:= (0, 6, 1, a, 1, r). COMPUTATIONS ARE SIMPLIFIED

BECOMES

A SQUARE

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{X,Y}^{(2)} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y)\}_{X,Y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): 118 -> F(110) IN B.

SATISFY ING:

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$id \otimes F_{Y,z}^{(2)} \downarrow \qquad \qquad \downarrow F_{X\otimes Y,z}^{(2)}$$

$$F(X) \otimes^{b} F(Y\otimes^{e}z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X\otimes^{e}Y) \otimes^{b} F(z)$$

$$\begin{array}{ccc}
\mathbb{I}_{\emptyset} \otimes_{\mathbb{I}} E(X) & \xrightarrow{\mathbb{I}_{(S)}^{\emptyset}} & E(X) \\
\mathbb{I}_{(S)} \otimes_{\mathbb{I}} E(X) & \xrightarrow{\mathbb{I}_{(S)}^{(S)}} & E(X)
\end{array}$$

$$\begin{array}{ccc}
\mathbb{I}_{(S)} \otimes_{\mathbb{I}} E(X) & \xrightarrow{\mathbb{I}_{(S)}^{(S)}} & E(X)
\end{array}$$

WHEN STRICT TAKE MONOIDAL CATEGORIES C:= (e, &, 1, a, 1, r) \$ 0:= (0, 6, 1, a, 1, r). COMPUTATIONS ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{X,Y}^{(2)} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y)^{0} \times_{/Y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): UB -> F(UC) IN B.

SATISFYING:

$$F(X) \otimes^{0} F(Y) \otimes^{0} F(z) \xrightarrow{F_{X,Y}^{(2)}} \otimes^{id}$$

$$F(X) \otimes^{0} F(Y) \otimes^{0} F(z) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{0} Y) \otimes^{0} F(z)$$

$$F(X) \otimes^{0} F(Y \otimes^{0} z) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{0} Y \otimes^{0} z)$$

UKEWISE  $1^{0} \otimes F(x)$   $\xrightarrow{\int_{F(x)}^{0}} F(x)$   $F(x) \otimes 1^{0} \xrightarrow{\int_{F(x)}^{0}} F(x)$ THESE ARE  $F(x) \otimes F(x)$   $F(x) \otimes F(x)$ 

$$F(x) \otimes^{a} L^{0} \xrightarrow{F(x)} F(x)$$

$$id \otimes F^{(a)} \downarrow \qquad = \uparrow F(x)$$

$$F(x) \otimes^{a} F(L^{e}) \xrightarrow{F(x)} F(x \otimes^{e} L)$$

WHEN STRICT TAKE MONOIDAL CATEGORIES & := (&, &, 1, a, 1, r) \$ 0 := (0, 8, 1, a, 1, r)

COMPUTATIONS

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{X,Y}^{(2)} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y)\}_{X,Y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): UB -> F(UC) IN B.

SATISFYING:

$$F(X) \otimes^{0} F(Y) \otimes^{0} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{0} Y) \otimes^{0} F(z)$$

$$id \otimes F_{Y,z}^{(2)} \downarrow \qquad \qquad \downarrow F_{X \otimes Y,z}^{(2)} \downarrow \qquad \downarrow F_{X \otimes Y,z}^{(2)}$$

$$F(X) \otimes^{0} F(Y \otimes^{0} z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{0} Y) \otimes^{0} F(z)$$

UKEWISE  $F(x) \xrightarrow{\int_{F(x)}^{\delta}} F(x)$   $F(x) \xrightarrow{F(x)} F(x)$ THESE ARE  $F(x) \xrightarrow{F(x)} F(x)$   $F(x) \xrightarrow{F(x)} F(x)$ 

$$F(X) \overset{\text{id} \otimes F^{(o)}}{=} F(X)$$

$$F(X) \overset{\text{id} \otimes F^{(o)}}{=} F(X)$$

$$F(X) \overset{\text{form}}{=} F(X)$$

$$F(X) \overset{\text{form}}{=} F(X)$$

TAKE MONOIDAL CATEGORIES

(4, &, l', a', l', r') & D := (D, &, l', a', l', r').

COMPUTATIONS

ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{X,Y}^{(2)} : F(X) \otimes^{0} F(Y) \longrightarrow F(X \otimes^{0} Y)^{0} \times_{y} Y \in \mathcal{C}.$
- (c) A MORPHISM F(0): 118 -> F(118) IN B.

SATISFY ING:

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$id \otimes F_{Y,z}^{(2)} \downarrow \qquad \qquad \downarrow F_{X\otimes Y,z}^{(2)}$$

$$F(X) \otimes^{b} F(Y \otimes^{c} z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$F(X) \otimes^{b} F(Y \otimes^{c} z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

UKEWISE F(x) id

THESE ARE F(x) F(x) F(x) F(x) F(x) F(x) F(x)

$$F(x) \overset{\text{id}}{\otimes} F(x) \xrightarrow{\text{id}} F(x)$$

$$F(x) \overset{\text{id}}{\otimes} F(x) \xrightarrow{\text{id}} F(x)$$

TAKE MONOIDAL CATEGORIES

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ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{ F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y) \}_{x,y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(11c) IN B.

SATISFYING:

GET:

$$F_{L_1X}^{(2)} \circ (F^{(0)} \otimes id_{F(X)}) = id_{F(X)}$$

 $F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$   $\downarrow^{id} \otimes^{e} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(x) \otimes^{e} F(z)$   $\downarrow^{id} \otimes^{e} F(X) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(x) \otimes^{e} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$   $\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$ 

$$F(\mathcal{L}^{e}) \otimes id \downarrow 2$$

$$F(\mathcal{L}^{e}) \otimes^{a} F(X) \xrightarrow{F_{\mathcal{L},X}^{(e)}} F(X)$$

$$F(x) \stackrel{\text{id}}{\approx} F(x) \xrightarrow{\text{id}} F(x)$$

$$F(x) \stackrel{\text{id}}{\approx} F(x) \xrightarrow{\text{krit.}} F(x)$$

TAKE MONOIDAL CATEGORIES

WHEN STRICT

WHEN STRICT

LIPENTITIES

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(OMPUTATIONS)
ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{ F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y) \}_{x,y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(11") IN B.

### SATISFYING:

GET:

$$F_{\mathbb{L}_{1}X}^{(2)} \circ (F^{(0)} \otimes id_{F(X)}) = id_{F(X)}$$

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$\downarrow^{id} \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$\downarrow^{id} \otimes^{b} F(X) \otimes^{b} F(Z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{b} F(Z)$$

$$\downarrow^{id} \otimes^{b} F(X) \otimes^{b} F(Z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(Z)$$

$$\downarrow^{id} \otimes^{b} F(X) \otimes^{b} F(Z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(Z)$$

$$\downarrow^{id} \otimes^{b} F(Z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(Z)$$

$$F(x) = id$$

$$F(x) = id$$

$$F(x) = F(x) \xrightarrow{F(x)} F(x)$$

$$F(x) \stackrel{\text{id} \otimes F^{(0)}}{=} 2$$

$$F(x) \stackrel{\text{e}}{\otimes} F(L^{e}) \stackrel{F^{(2)}}{\xrightarrow{x_{1}}} F(x)$$

TAKE MONOIDAL CATEGORIES

(4, &, l', a', l', r') & D := (D, &, l', a', l', r').

(OMPUTATIONS)
ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO B CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{ F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y) \}_{x,y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(11c) IN B.

### SATISFY ING:

GET:

$$F_{L_1X}^{(2)} \circ (F^{(0)} \otimes id_{F(X)}) = id_{F(X)}$$

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$\downarrow^{id} \otimes^{e} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(x) \otimes^{e} F(z)$$

$$\downarrow^{id} \otimes^{e} F(X) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(x) \otimes^{e} F(z)$$

$$\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(x) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$\downarrow^{id} \otimes^{e} F(x) \otimes^{e} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$F(\mathcal{L}^{(0)}) \stackrel{\text{id}}{=} F(x) \stackrel{\text{id}}{=} F(x) \stackrel{\text{id}}{=} F(x) \stackrel{\text{id}}{=} F(x)$$

$$F(\mathcal{L}^{(0)}) \stackrel{\text{id}}{=} F(x) \stackrel{\text{id}}{=} F(x) \stackrel{\text{id}}{=} F(x)$$

$$F(x) \stackrel{id}{\approx} F(x) \stackrel{id}{=} F(x)$$

$$F(x) \stackrel{a}{\approx} F(x) \stackrel{f(x)}{=} F(x)$$

WHEN STRICT

TAKE MONOIDAL CATEGORIES

(e.= (e, &, l', a, l', r') & D := (D, &, le, a, le, re).

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{ F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y) \}_{x,y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(11c) IN B.

### SATISFYING:

GET:

$$F_{\mathbb{L}_{1}X}^{(2)} \circ \left(F^{(0)} \otimes id_{F(X)}\right) = id_{F(X)}$$

$$\left(F^{(0)} \otimes id_{F(X)}\right) \circ F_{\mathbb{L}_{1}X}^{(2)} = id_{F(\mathbb{L}^{2})} \otimes^{a} F(X)$$

$$F(\mathcal{L}_{6}) \otimes_{g} F(x) \stackrel{id}{\leftarrow} F(\mathcal{L}_{6}) \otimes_{g} F(x) \xrightarrow{F_{(5)}} F(x)$$

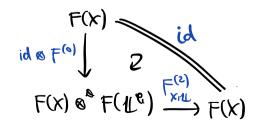
$$F(\mathcal{L}_{6}) \otimes_{g} F(x) \stackrel{id}{\leftarrow} F(x) \otimes_{g} F(x) \xrightarrow{F_{(5)}} F(x)$$

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$id \otimes F_{Y|z}^{(2)} \downarrow \qquad \qquad \downarrow F_{X\otimes Y|z}^{(2)}$$

$$F(X) \otimes^{b} F(Y \otimes^{c} z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$

$$F(X) \otimes^{b} F(Y \otimes^{c} z) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(z)$$



COMPUTATIONS

TAKE MONOIDAL CATEGORIES

(4, &, l', r') & D := (D, &, l', a, l', r').

COMPUTATIONS

ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{ F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y) \}_{x,y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(11c) IN B.

### SATISFYING:

GET:

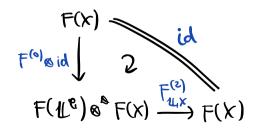
$$F_{\mathbb{L}_{1}X}^{(2)} \circ \left(F^{(0)} \otimes id_{F(X)}\right) = id_{F(X)}$$

$$\left(F^{(0)} \otimes id_{F(X)}\right) \circ F_{\mathbb{L}_{1}X}^{(2)} = id_{F(\mathbb{L}^{2})} \otimes^{a}_{F(X)}$$

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(\xi) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{e} Y) \otimes^{b} F(\xi)$$

$$id \otimes F_{Y|\xi}^{(2)} \downarrow \qquad \qquad \downarrow F_{X \otimes Y|\xi}^{(2)}$$

$$F(X) \otimes^{b} F(Y \otimes^{e} \xi) \xrightarrow{F_{X,Y}^{(2)} \otimes id} F(X \otimes^{e} Y) \otimes^{b} F(\xi)$$



$$F(x) = id$$

$$F(x) = F(x) = id$$

$$F(x) = F(x) = F(x)$$

$$F(x) = F(x)$$

$$F(x) = F(x)$$

$$F(x) = F(x)$$

TAKE MONOIDAL CATEGORIES

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{ F_{X,Y}^{(2)} : F(X) \otimes^{Q} F(Y) \longrightarrow F(X \otimes^{Q} Y) \}_{X,Y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): UB -> F(UC) IN B.

SATISFYING:  $(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)}} \otimes id$   $(F(X) \otimes^{b} F(Y)) \otimes^{b} F(z) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} Y) \otimes^{b} F(z) \xrightarrow{F_{X,X,Y}^{(2)}} F(X \otimes^{c} Y) \otimes^{c} z )$   $(F(X) \otimes^{b} (F(Y) \otimes^{b} F(z)) \xrightarrow{id \otimes F_{Y,z}^{(2)}} F(X) \otimes^{b} F(Y \otimes^{c} z) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} (Y \otimes^{c} z))$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X) \xrightarrow{F(X)} F(X \otimes^{b} I_{x}^{(2)} \xrightarrow{F_{X,Y}^{(2)}} F(X)$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X) \xrightarrow{F(X)} F(X \otimes^{c} I_{x}^{(2)} \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)})$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)})$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)})$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)})$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)})$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{F_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)})$   $(F(X) \otimes^{b} F(X)) \xrightarrow{f_{X,Y}^{(2)}} F(X \otimes^{c} I_{x}^{(2)}) \xrightarrow{$ 

TAKE MONOIDAL CATEGORIES

C:= (&, &, l', à, l', r') & D:= (D, &, le, a, le, a).

COMPUTATIONS

ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO & CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> B.
- (6) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y)^{0} \times_{y} \in \mathcal{C}.$
- (c) A MORPHISM F(0): 118 -> F(11c) IN B.

SATISFYING:

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$F(X) \otimes^{b} F(Y) \otimes^{c} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{c} F(z)$$

$$F(X) \otimes^{b} F(Y) \otimes^{c} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{c} F(z)$$

$$F(x) = \frac{1}{|x|} F(x)$$

$$F(x) = \frac{1}{|x|} F(x)$$

$$F(x) = \frac{1}{|x|} F(x)$$

$$F(x) = id$$

$$F(x) = F(x)$$

$$F(x) = F(x)$$

$$F(x) = F(x)$$

$$F(x) = F(x)$$

TAKE MONOIDAL CATEGORIES

(e, &, l', a', l', r') & D := (D, &, le, a, le, re).

COMPUTATIONS ARE SIMPLIFIED

A MONOIDAL FUNCTOR FROM & TO B CONSISTS OF:

- (a) A FUNCTOR BTW CATEGORIES F: & -> D.
- (b) A NATURAL TRANSFORMATION  $F^{(2)} = \{F_{x,y}^{(2)} : F(x) \otimes^{0} F(y) \longrightarrow F(x \otimes^{0} y)\}_{x,y \in \mathcal{C}}.$
- (c) A MORPHISM F(0): 118 -> F(114) IN B.

WOULD BENICE IF
WE WERE ALWAYS ABLE
TO REDUCE TO
THE STRICT CASE

SATISFYING:

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(X) \otimes^{b} F(z)$$

$$F(X) \otimes^{b} F(Y) \otimes^{b} F(z) \xrightarrow{F_{(2)}^{(2)} \otimes id} F(z)$$

$$F(x) = id$$

$$F(x) \otimes F(x) \xrightarrow{F(x)} F(x)$$

$$F(X) \otimes^{a} F(L^{e}) \xrightarrow{F(2)} F(X)$$

$$F(X) \otimes^{a} F(L^{e}) \xrightarrow{F(2)} F(X)$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

# STRATEGY

- RECALL (End(A), ⊗= 0, L=IdA) IS STRICT.
- DEFINE STRICT MONOIDAL CATEGORY & Str MODELLED ON Endmod-& (Greg).

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

# STRATEGY

- RECALL (End(A), ⊗= 0, L=IdA) IS STRICT.
- DEFINE STRICT MONOIDAL CATEGORY & str MODELLED ON Endmod-& (Greg).
- Build Fully FAITHFUL, STRONG MONOIDAL FUNCTOR:
   ρ: ← → ζ<sup>str</sup>

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

# STRATEGY

- RECALL (End(A), ⊗= 0, L=IdA) IS STRICT.
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- Build Fully FAITHFUL, STRONG MONOIDAL FUNCTOR:
   ρ: ← → ζ<sup>str</sup>
- GET "ESSENTIAL IMAGE OF P" WHICH IS STRICT. 111

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

• DEFINE STRICT MONOIDAL CATEGORY & str MODELLED ON Endmod-& (Greg).

OBJECTS: (FE END (C), 
$$u := (u_{M/X} : F(M) \otimes X \xrightarrow{\sim} F(M \otimes X)) \int_{M/X \in \mathcal{E}}$$
 $f(M \otimes X) \otimes Y$ 
 $f(M \otimes X) \otimes Y$ 

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

• DEFINE STRICT MONOIDAL CATEGORY & str MODELLED ON Endmod-& (Greg).

MORPHISMS:  $(F_{\mu}) \rightarrow (F'_{\mu}u') = NAT'LTRANSF \Theta: F \rightarrow F' \rightarrow$ 

$$F(M) \otimes X \xrightarrow{u_{M/X}} F(M \otimes X)$$

$$\Theta_{M} \otimes id_{X} \downarrow \qquad \qquad \downarrow \Theta_{M \otimes X}$$

$$F(M) \otimes X \xrightarrow{u'_{M/X}} F(M \otimes X)$$

YM,XEC.

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

• DEFINE STRICT MONOIDAL CATEGORY & str MODELLED ON Endmod-& (Greg).

MORPHISMS:  $(F,u) \rightarrow (F',u') = NATIL TRANSF : \theta:F \rightarrow F'$ + COMPATIBILITY AXIOM

$$\frac{MoNoIDAL PRODUCT: (F, u) \otimes^{Str} (F', u') := (FF', u'')}{U''_{M,X}: FF'(M) \otimes X \xrightarrow{U_{F'(M) \otimes X}} F(F'(M) \otimes X) \xrightarrow{F(U'_{M,X})} FF'(M \otimes X)}$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

• DEFINE STRICT MONOIDAL CATEGORY & str MODELLED ON Endmod-& (Greg).

OBJECTS: (FE End (L), U := (UMX: F(M) &X ~ F(M&X)) MX & + PENTAGON AXIOM

MORPHISMS:  $(F_{\mu}) \rightarrow (F'_{\mu}u') = NATIL TRANSF : \theta: F \rightarrow F'$ + COMPATIBILITY AXIOM

 $\frac{MoNOIDAL PRODUCT: (F,u) \otimes^{Str} (F,u') := (FF', u'')}{U''_{M,x}: FF'(M) \otimes X \xrightarrow{UF'(M) \otimes X} F(F'(M) \otimes X) \xrightarrow{F(U'_{M,x})} FF'(M \otimes X)}$ 

MONOIDAL UNIT: 11str := (Ide, lidnox Inxer)

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

# STRATEGY

PRECALL (End(A), ⊗= 0, L=IdA) IS STRICT.

DEFINE STRICT MONOIDAL CATEGORY & Str MODELLED ON Endmod-& (Greg).

- Build Fully FAITHFUL, STRONG MONOIDAL FUNCTOR:
   ρ: ← → ζ<sup>str</sup>
- GET & "ESSENTIAL IMAGE OF P" WHICH IS STRICT. 111

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

· BUILD FULLY FAITHFUL, STRONG MONOIDAL FUNCTOR:

$$(F_{\mu}) \rightarrow (F'_{\mu}u') \equiv \Theta: F \Rightarrow F'$$
  
+ COMPATIBILITY AXIOM

$$(F, u) \otimes^{Str} (F', u') := (FF', u'')$$

$$U''_{M,X} := F(U'_{M,X}) \circ U_{F'(M) \otimes X}$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

BUILD FULLY FAITHFUL, STRONG MONOIDAL FUNCTOR:

$$(F_{\mu}) \rightarrow (F'_{\mu}) = \theta : F \rightarrow F'$$
  
+ COMPATIBILITY AXIOM

$$(F, u) \otimes^{Str} (F', u') := (FF', u'')$$
 
$$u''_{M,X} := F(u'_{M,X}) \circ u_{F'(M) \otimes X}$$

$$\begin{array}{c}
\rho: \ell \longrightarrow \zeta^{str} \\
\downarrow U:= \{U_{M/X}: F(M) \otimes X \xrightarrow{\sim} F(M \otimes X)\}_{M/X} \in \ell \\
\uparrow PENTAGON AXION
\end{array}$$

$$\begin{array}{c}
DEFINE \rho(z) \\
\vdots = \left\{Z \otimes - : \ell \to \ell \\
U^{z}:= \left\{U_{M/X}: (Z \otimes M) \otimes X \xrightarrow{\sim} Z \otimes (H \otimes X)\right\}\right\} \\
\downarrow U^{z}:= \left\{U_{M/X}: (Z \otimes M) \otimes X \xrightarrow{\sim} Z \otimes (H \otimes X)\right\}$$

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\end{array}$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

BUILD FULLY FAITHFUL, STRONG MONOIDAL FUNCTOR:

$$(F_{\mu}) \rightarrow (F'_{\mu}u') \equiv \theta : F \Rightarrow F'$$
+ COMPATIBILITY AXIOM

$$(F, u) \otimes^{Str} (F', u') := (FF', u'')$$

$$U''_{M,X} := F(U'_{M,X}) \circ U_{F'(M) \otimes X}$$

$$\begin{array}{c}
\rho: \mathcal{C} \longrightarrow \mathcal{C}^{str} \\
\downarrow u := \{u_{M/X}: F(M) \otimes X \xrightarrow{\sim} F(M \otimes X)\}_{M/X} \in \mathcal{C} \\
\uparrow PENTAGON AXION
\end{array}$$

$$\begin{array}{c}
DEFINE \rho(z) \\
\vdots = \left( \frac{2 \otimes - : \mathcal{C}}{2 \otimes M} \otimes X \xrightarrow{\sim} 2 \otimes (M \otimes X) \right) \\
\downarrow u := \left( \frac{2 \otimes - : \mathcal{C}}{2 \otimes M} \otimes X \xrightarrow{\sim} 2 \otimes (M \otimes X) \right) \\
\downarrow u := \left( \frac{2 \otimes - : \mathcal{C}}{2 \otimes M} \otimes X \xrightarrow{\sim} 2 \otimes (M \otimes X) \right)$$

$$\begin{array}{c}
M_{X} \in \mathcal{C} \\
0 \neq 1 & M_{X} \neq 0
\end{array}$$

$$(F, u) \otimes^{Str} (F', u') := (FF', u'')$$

$$U''_{M,x} := F(U'_{M,x}) \circ U_{F'(M) \otimes X}$$

$$1 := (Ide, \{id_{M \otimes X}\}_{M,x \in E})$$

$$\downarrow p(2 + 2')$$

$$:= (\Theta : (2 \otimes -) \implies (2' \otimes -)$$

$$WITH \Theta_{Y} : 2 \otimes Y \longrightarrow 2 \otimes Y$$

$$f \otimes id_{Y}$$

$$F \in SURE$$

$$To CHECK \dots PROPERLY DEFINES p$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

$$\begin{cases}
F \in \text{End}(C), \\
u := \{u_{M/X} : F(M) \otimes X \xrightarrow{\sim} F(M \otimes X)\}_{M/X \in C} \\
+ PENTAGON AXION
\end{cases}$$

$$(F_{\mu}) \rightarrow (F'_{\mu}u') = \theta : F \Rightarrow F'$$
  
+ COMPATIBILITY AXIOM

$$(F, u) \otimes^{str} (F', u') := (FF', u'')$$
 
$$u''_{M,x} := F(u'_{M,x}) \circ u_{F'(M) \otimes x}$$

$$\sharp \rho(z \xrightarrow{f} z')$$

$$:= \left( \begin{array}{c} \theta : (z \otimes -) \\ \theta : (z \otimes -) \\ \theta : (z \otimes -) \end{array} \right) \xrightarrow{\text{def II}} z' \otimes \gamma \longrightarrow z' \otimes \gamma$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

BUILD FULLY FAITHFUL, STRONG MONOIDAL, FUNCTOR:
 p: 4 → C<sup>Str</sup>

DEFINE 
$$\rho(z) := \begin{cases} z \otimes - : & \Rightarrow z \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases} \\ \lambda^{z} := \begin{cases} \lambda^{z} : (z \otimes M) \otimes x \xrightarrow{\sim} z \otimes M \otimes x \end{cases}$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

BUILD FULLY FAITHFUL, STRONG MONOIDAL, FUNCTOR:
 P: 4 → 2<sup>str</sup>

DEFINE 
$$\rho(z) := \begin{cases} z \otimes - : & \varphi & \varphi \\ V^{z} := \begin{cases} V^{z} : & (z \otimes M) \otimes \chi & \xrightarrow{\sim} z \otimes (M \otimes \chi) \end{cases} \\ V^{z} := \begin{cases} V^{z} : & (z \otimes M) \otimes \chi & \xrightarrow{\sim} z \otimes (M \otimes \chi) \end{cases} \\ V^{z} := \begin{cases} V^{z} : & (z \otimes M) \otimes \chi & \xrightarrow{\sim} z \otimes (M \otimes \chi) \end{cases} \\ V^{z} := \begin{cases} V^{z} : & (z \otimes M) \otimes \chi & \xrightarrow{\sim} z \otimes (M \otimes \chi) \end{cases}$$

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$$\begin{vmatrix}
\alpha^{M_1 \xi} \cdot \\ -1 \\ 0 & 0
\end{vmatrix} = \beta(M) \circ \beta(\xi) - \beta(M \otimes \xi) \\
\beta^{M_1 \xi} \cdot \beta(M) \otimes^{SHL} \beta(\xi) = \beta(M) \circ \beta(\xi) - \beta(M \otimes \xi)$$

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\end{vmatrix} = \beta(M) \circ \beta(K) \circ \beta(K) \circ \beta(K)$$

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STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

BUILD FULLY FAITHFUL, STRONG MONOIDAL, FUNCTOR:
 p: 4 → C<sup>Str</sup>
 ≡ READ DETAILS =

DEFINE 
$$\rho(z) := \begin{cases} z \otimes - : & \varphi & \varphi \\ U^{z} := \begin{cases} U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \end{cases} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\} \\ 0 : = \left\{ U_{M|X} : & (z \otimes M) \otimes X \xrightarrow{\sim} z \otimes (M \otimes X) \right\}$$

$$\begin{vmatrix}
\alpha_{M,\varsigma} : \\ \beta_{M,\varsigma} : \\ \beta$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

# STRATEGY

- PRECALL (End(A), ⊗= 0, L=IdA) IS STRICT.
- DEFINE STRICT MONOIDAL CATEGORY & Str MODELLED ON Endmod-& (Greg).
- BUILD FULLY FAITHFUL, STRONG MONOIDAL FUNCTOR:

  p: 4 -> Cstr
  - GET & "ESSENTIAL IMAGE OF P" WHICH IS STRICT. 111

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

# STRATEGY

RECALL (End(A),  $\otimes = \circ$ ,  $L = Id_A$ ) IS STRICT.

DEFINE STRICT MONOIDAL CATEGORY & STr MODELLED ON Endmod-r. (Greg).

BUILD FULLY FAITHFUL, STRONG MONOIDAL FUNCTOR:

P: L -> CSTr

THE FULL SUBCATEGORY OF COTT

ON OBJECTS = TO p(X) FOR SOME X = C

WHICH IS STRICT. //

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

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BUILD FULLY FAITHFUL, STRONG MONOIDAL FUNCTOR:

P: C -> CST EXER3.16. MONOIDAL

THE FULL SUBCATEGORY OF CST

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WHICH IS STRICT. //

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE N OBJECTS:  $X_1, X_2, ..., X_n \in \mathcal{C}$ .

AN EXPRESSION  $w(X_1, X_2, ..., X_n)$  is an object in  $\mathcal{C}$ .

FORMED BY INSERTING PARENTHESES AND LS

IN  $X_1 \otimes X_2 \otimes ... \otimes X_n$ 

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E.G. EXPRESSIONS IN XOYOZ:

$$(X \otimes Y) \otimes (1 \otimes z)$$
  $1 \otimes (X \otimes ((Y \otimes L) \otimes z))$ 

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I way of moving from one expression in 3 OBJECTS TO ANOTHER

\* VIA ASSOCIATIVITY & WITALITY MORPHISMS

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$$\mathbb{T}\otimes\left(\times\otimes\left((\lambda\otimes\pi)\otimes\varsigma\right)\right)$$

\* VIA ASSOCIATIVITY & WITALITY MORPHISMS

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FORMED BY INSERTING PARENTHESES AND LS

IN  $X_1 \otimes X_2 \otimes ... \otimes X_n$ 

I way of moving from one expression in 3 OBJECTS TO ANOTHER

$$(X \otimes Y) \otimes (1 \otimes 2)$$

$$(X \otimes Y) \otimes (1 \otimes 2)$$

$$* VIA ASSOCIATIVITY & WITALITY MORPHISMS
$$(X \otimes ((Y \otimes L) \otimes 2))$$$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE N OBJECTS:  $X_1, X_2, ..., X_n \in \mathcal{C}$ .

AN EXPRESSION  $w(X_1, X_2, ..., X_n)$  is an object in  $\mathcal{C}$ .

FORMED BY INSERTING PARENTHESES AND LS

IN  $X_1 \otimes X_2 \otimes ... \otimes X_n$ 

I way of moving from one expression in 3 OBJECTS TO ANOTHER

$$(X \otimes Y) \otimes (1 \otimes Z)$$

$$(X$$

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TAKE N OBJECTS:  $X_1, X_2, ..., X_n \in \mathcal{C}$ .

AN EXPRESSION  $w(X_1, X_2, ..., X_n)$  is an object in  $\mathcal{C}$ .

FORMED BY INSERTING PARENTHESES AND LS

IN  $X_1 \otimes X_2 \otimes ... \otimes X_n$ 

I WAY OF MOVING FROM ONE EXPRESSION IN 3 OBJECTS TO ANOTHER

$$(X \otimes Y) \otimes (X \otimes Y) \otimes (X \otimes (Y \otimes U)) \otimes (X \otimes U)$$

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE N OBJECTS: X1, X2, ..., Xn & C. AN EXPRESSION W(X1, X2, ..., Xn) IS AN OBJECT IN & FORMED BY INSERTING PARENTHESES AND LS IN XI & X2 & ... & Xn

I WAY OF MOVING FROM ONE EXPRESSION IN 3 OBJECTS TO ANOTHER

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE N OBJECTS:  $X_1, X_2, ..., X_n \in \mathcal{C}$ .

AN EXPRESSION  $w(X_1, X_2, ..., X_n)$  is an object in  $\mathcal{C}$ .

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THE PENTAGON & TRIANGLE AXIOMS IMPLY THAT

J! WAY OF MOVING FROM ONE EXPRESSION IN 3 OBJECTS TO ANOTHER

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THE PENTAGON & TRIANGLE AXIOMS IMPLY THAT

J! WAY OF MOVING FROM ONE EXPRESSION IN 4 OBJECTS TO ANOTHER

\* VIA ASSOCIATIVITY & WITALITY MORPHISMS

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

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THE PENTAGON & TRIANGLE AXIOMS IMPLY THAT

J! WAY OF MOVING FROM ONE EXPRESSION IN 4 OBJECTS TO ANOTHER

?? WHAT ABOUT >5 OBJECTS ??

\* VIA ASSOCIATIVITY & WITALITY MORPHISMS

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE N OBJECTS:  $X_1, X_2, ..., X_n \in \mathcal{C}$ .

AN EXPRESSION  $w(X_1, X_2, ..., X_n)$  is an object in  $\mathcal{C}$ .

FORMED BY INSERTING PARENTHESES AND LS

IN  $X_1 \otimes X_2 \otimes ... \otimes X_n$ 

THE PENTAGON & TRIANGLE AXIOMS IMPLY THAT

J! WAY OF MOVING FROM ONE EXPRESSION IN 4 OBJECTS TO ANOTHER

OF STILL HOLDS ... ?? WHAT ABOUT > 5 OBJECTS ??

COHERENCE THEOREM: LET  $f,g:w(X_1,...,X_n)\longrightarrow w'(X_1,...,X_n)\in \mathcal{C}$  BE ISOS COMPRISED OF ASSOC. & UNITALITY MORPHISMS\*. THEN, f=g.

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

PROOF IDEA:

COHERENCE THEOREM: LET  $f,g:w(X_1,...,X_n)\longrightarrow w'(X_1,...,X_n)\in \mathcal{C}$ BE ISOS COMPRISED OF ASSOC. & UNITALITY MORPHISMS\*. THEN, f=g.

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

PROOF IDEA: VIA THE STRICTIFICATION THEOREM,

HAVE FAITHFUL, STRONG MONOIDAL FUNCTOR

P: C -> cstr

COHERENCE THEOREM: LET  $f,g:w(X_1,...,X_n)\longrightarrow w'(X_1,...,X_n)\in \mathcal{C}$ BE ISOS COMPRISED OF ASSOC. & UNITALITY MORPHISMS\*. THEN, f=g.

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

PROOF IDEA: VIA THE STRICTIFICATION THEOREM,

HAVE FAITHFUL, STRONG MANOIDAL FUNCTOR  $\rho\colon \mathcal{C}\longrightarrow \mathcal{C}^{str}$ ARGUE THAT  $\rho(f)=\rho(g)$  USING STRONGNESS.

COHERENCE THEOREM: LET  $f,g:w(X_1,...,X_n)\longrightarrow w'(X_1,...,X_n)\in \mathcal{C}$ BE ISOS COMPRISED OF ASSOC. & UNITALITY MORPHISMS\*. THEN, f=g.

FAITHFULNESS => f=g.//

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE MONOIDAL CATEGORY (E, Ø, L, a, l, r)
WLOG VIA STRICTIFICATION THEOREM
CAN ASSUME STRICT (E, Ø, L)

STRICTIFICATION THEOREM: ANY MONOIDAL CATEGORY & IS MONOIDALLY EQUIVALENT TO A STRICT MONOIDAL CATEGORY.

TAKE MONOIDAL CATEGORY (6,0, 1, a, l, r) WLOG VIA STRICTIFICATION THEOREM CAN ASSUME STRICT (6, 0, 11) CAN ILLUSTRATE OBJECTS & MORPHISMS IN STRICT MONOIDAL CATEGORIES (6,0, L) HERE WITH NICE PICTURES

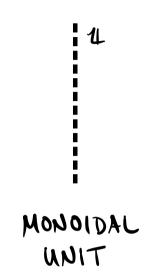
CAN ILLUSTRATE

OBJECTS & MORPHISMS IN

STRICT MONOIDAL CATEGORIES

(%, 0, 1L)

HERE WITH NICE PICTURES



CAN ILLUSTRATE

OBJECTS & MORPHISMS IN

STRICT MONOIDAL CATEGORIES

(%, 0, 1L)

HERE WITH NICE PICTURES

OR SOMETIMES
THE EMPTY STRING
IS USED

MONOIDAL

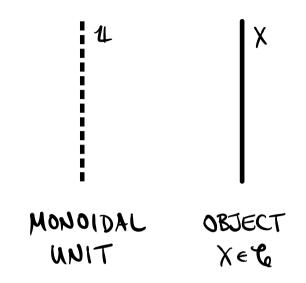
CAN ILLUSTRATE

OBJECTS & MORPHISMS IN

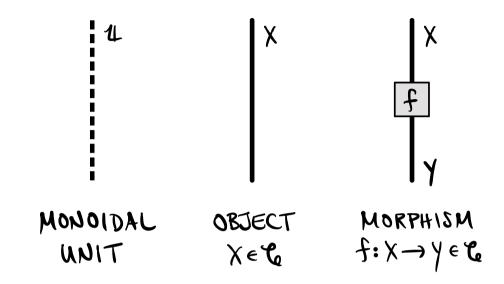
STRICT MONOIDAL CATEGORIES

(°C, ®, 1L)

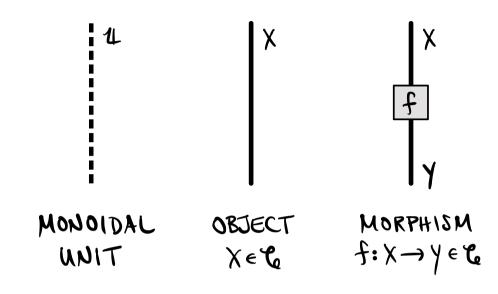
HERE WITH NICE PICTURES



CAN ILLUSTRATE
OBJECTS & MORPHISMS IN
STRICT MONOIDAL CATEGORIES
(%, 0, 1L)
HERE WITH NICE PICTURES



CAN ILLUSTRATE
OBJECTS & MORPHISMS IN
STRICT MONOIDAL CATEGORIES
(%, 0, 1L)
HERE WITH NICE PICTURES



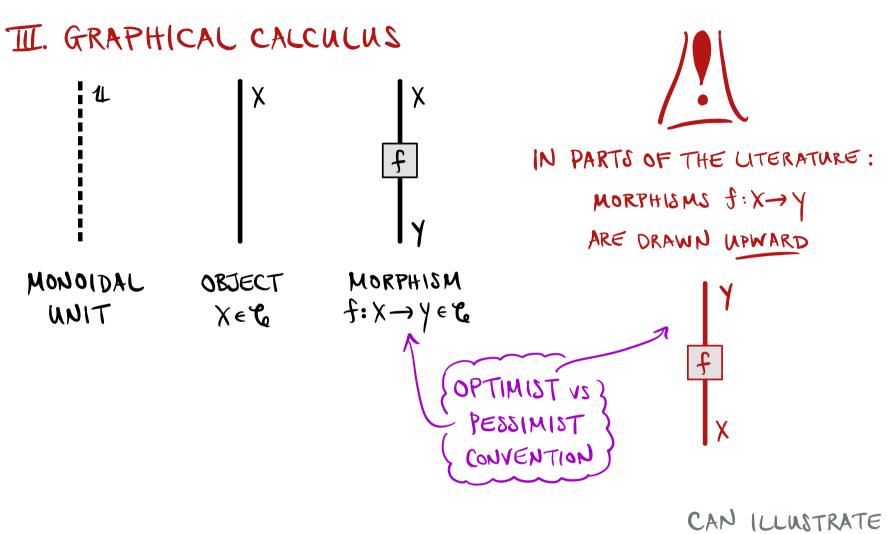


IN PARTS OF THE LITERATURE:

MORPHISMS f:X-Y ARE DRAWN UPWARD



CAN ILLUSTRATE
OBJECTS & MORPHISMS IN
STRICT MONOIDAL CATEGORIES
(%, 0), 1L)
HERE WITH NICE PICTURES



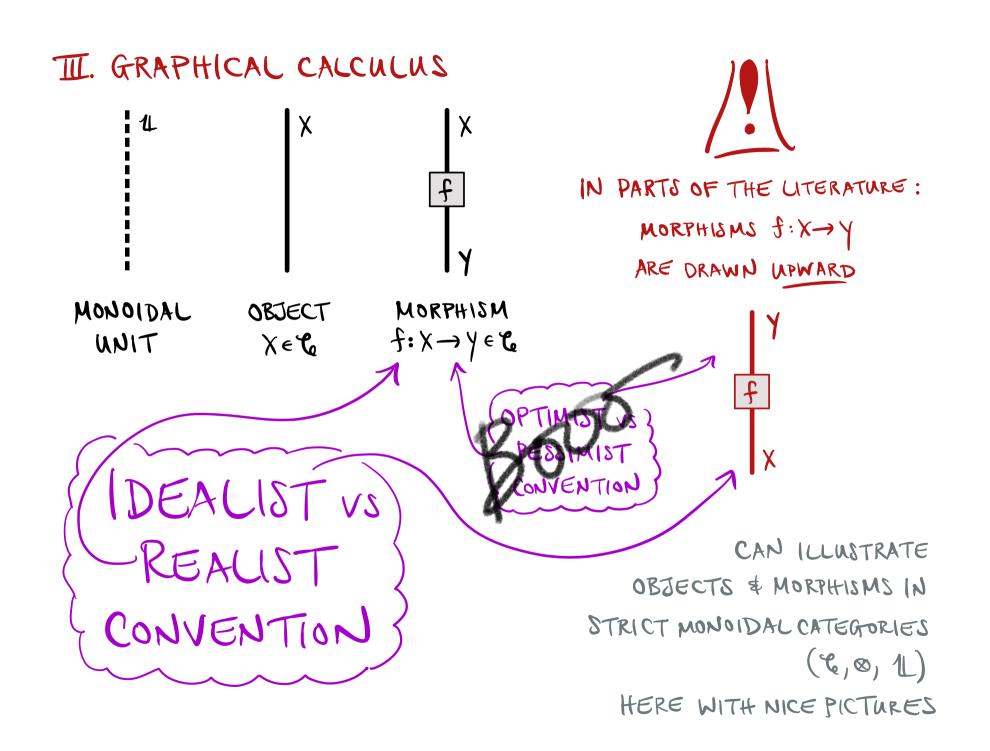
CAN ILLUSTRATE

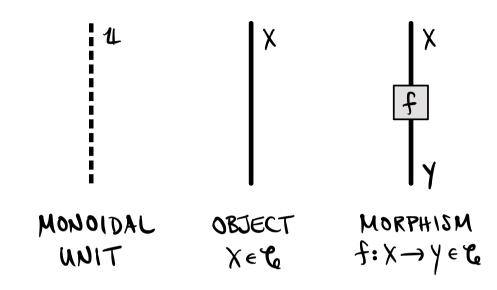
OBJECTS & MORPHISMS IN

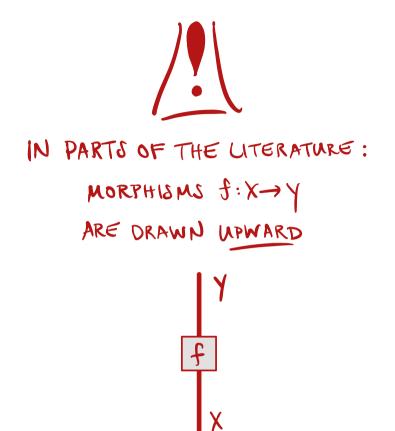
STRICT MONOIDAL CATEGORIES

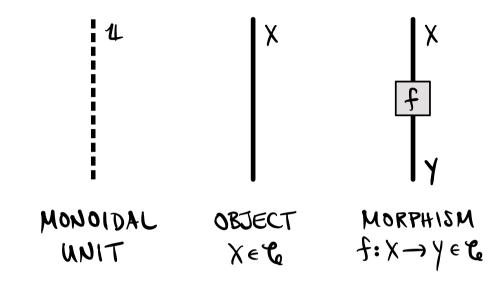
(°C, ®, 1L)

HERE WITH NICE PICTURES









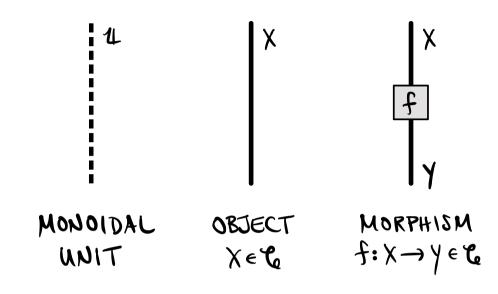


IN PARTS OF THE LITERATURE:

MORPHISMS f:X-Y ARE DRAWN UPWARD



COMPOSITION OF MORPHISMS f: X → Y, g: Y → z ∈ C



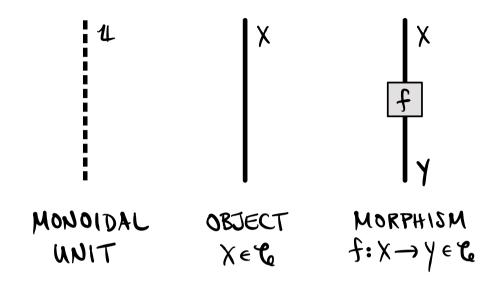


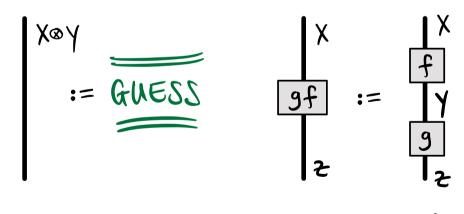
IN PARTS OF THE LITERATURE:

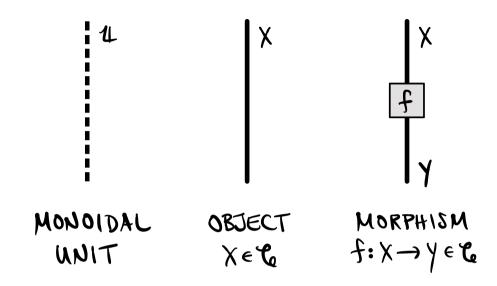
MORPHISMS f:X-Y ARE DRAWN UPWARD

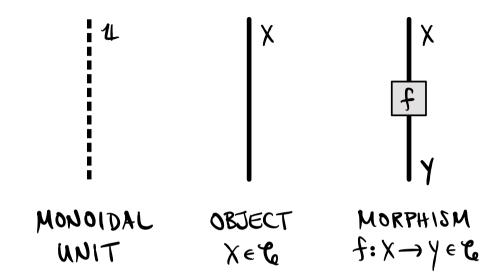


COMPOSITION OF MORPHISMS f: X → Y, g: Y → z ∈ C

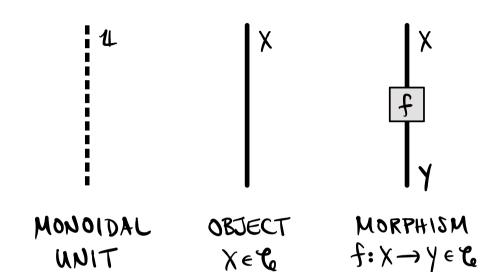


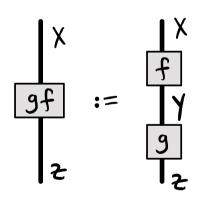




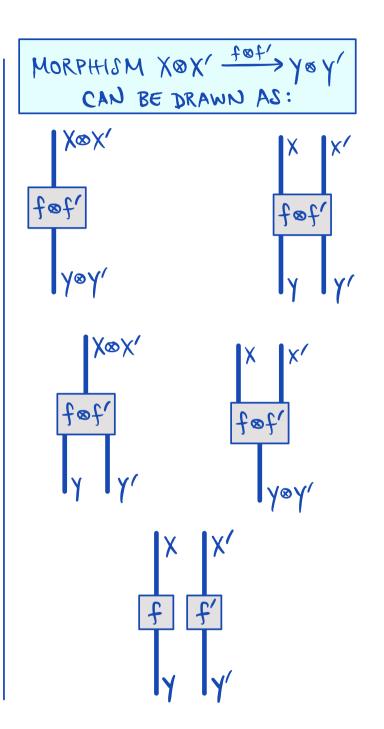


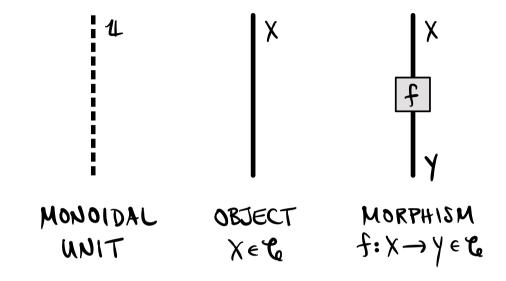
COMPOSITION OF MORPHISMS f:X-y, g:Y-zec MORPHISM XXX fot you



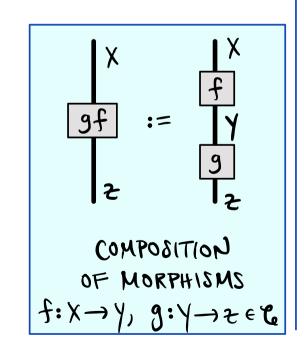


COMPOSITION OF MORPHISMS f: X→Y, g:Y→z ∈ C



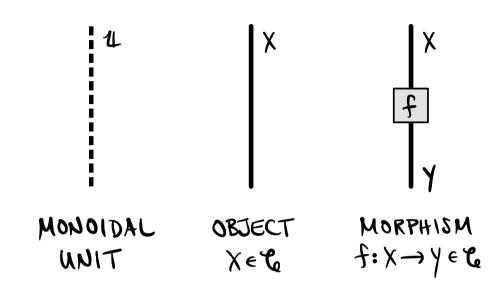


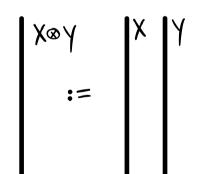
⊗ OF OBJECTS X,y ∈ C



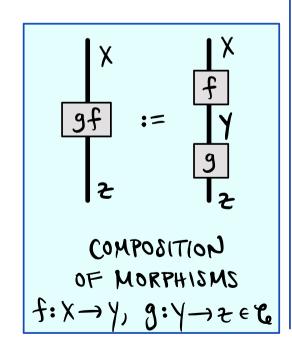
ASSOCIATIVITY:

UNITAUTY:

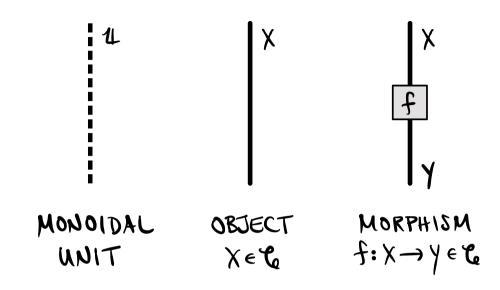


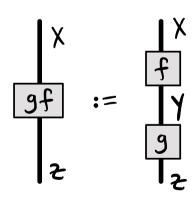


⊗ OF OBJECTS X, y ∈ C



$$\begin{array}{c} \text{UNITAUTy}: \\ \downarrow X \\ \downarrow f \\ \downarrow X \\ \downarrow f \\ \downarrow Y \\ \downarrow$$

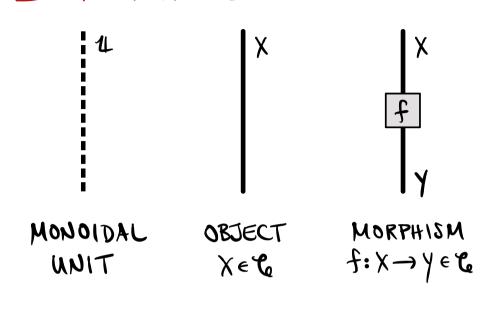


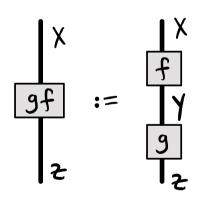


COMPOSITION OF MORPHISMS f:X→Y, g:Y→z ∈ C

### COMMUTATIVE DIAGRAM:

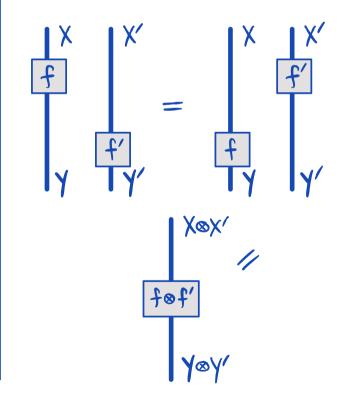
$$\begin{array}{c|c}
X \otimes X' \xrightarrow{f \otimes id_{X'}} Y \otimes X' \\
id_X \otimes f' \downarrow & 2 & \downarrow id_Y \otimes f' \\
X \otimes Y' \xrightarrow{f \otimes id_{Y'}} Y \otimes Y'
\end{array}$$

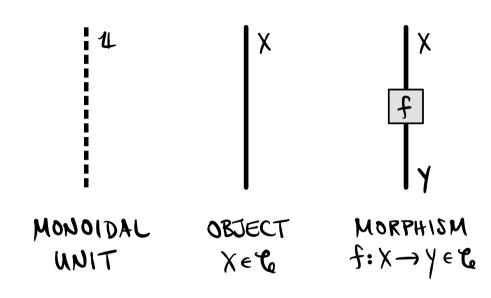


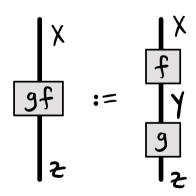


COMPOSITION OF MORPHISMS f:X→Y, g:Y→zeC

### COMMUTATIVE DIAGRAM:

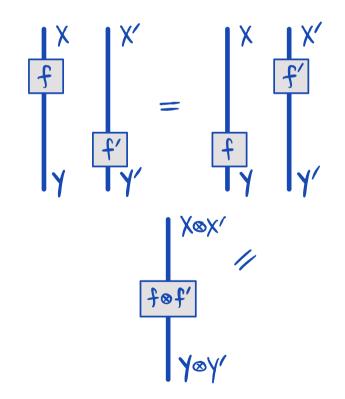


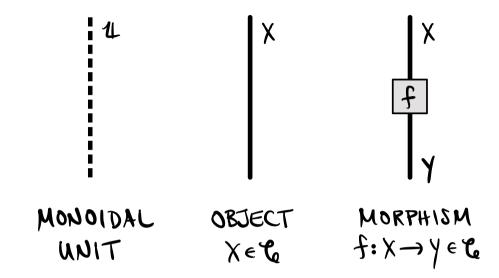




COMPOSITION OF MORPHISMS f: X→Y, g:Y→z ∈ C

### LEVEL EXCHANGE



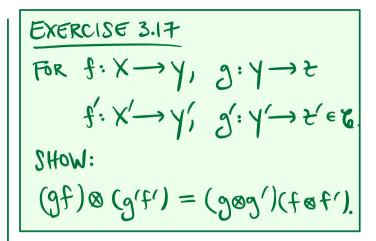


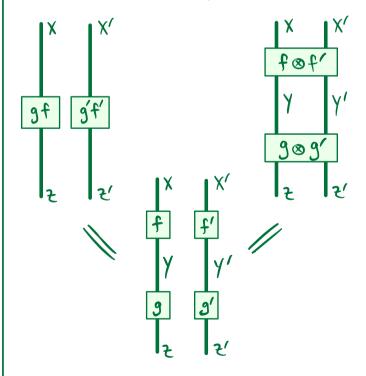
EXERCISE 3.17

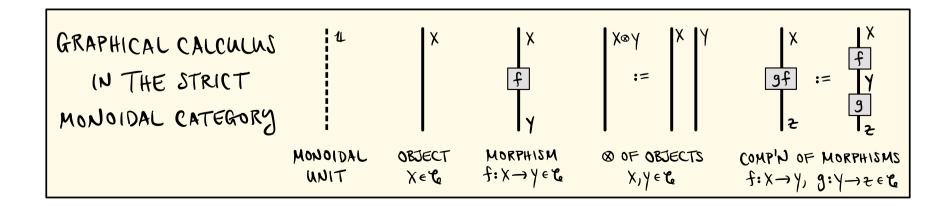
FOR 
$$f: X \rightarrow Y$$
,  $g: Y \rightarrow \xi$ 
 $f': X' \rightarrow Y'$ ,  $g': Y' \rightarrow \xi' \in \mathcal{E}$ 

SHOW:

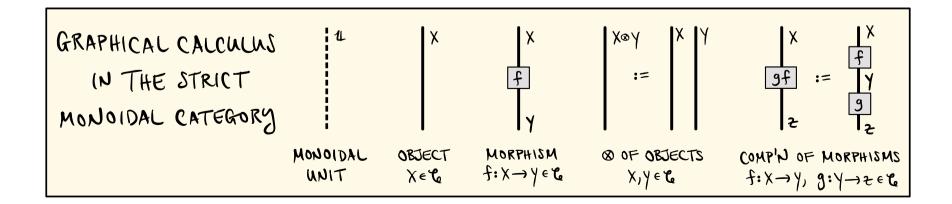
 $(9f) \otimes (g'f') = (g \otimes g')(f \otimes f')$ .



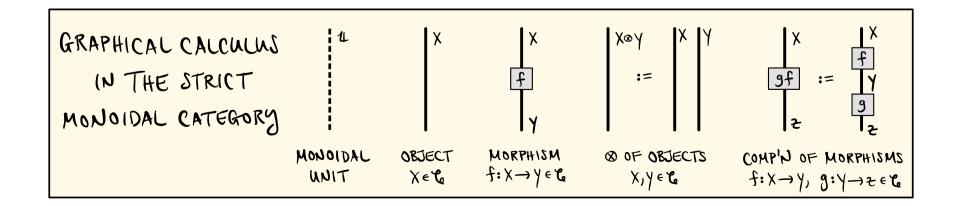




THESE ARE MONOIDAL CATEGORIES THAT CONTAIN DUAL STRUCTURES

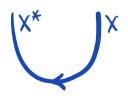


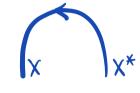
# IV. RIGID CATEGORIES: IN STRICT CASE A LEFT DUAL OF AN OBJECT X & US AN OBJECT X\* & C EQUIPPED WITH MORPHISMS



A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS





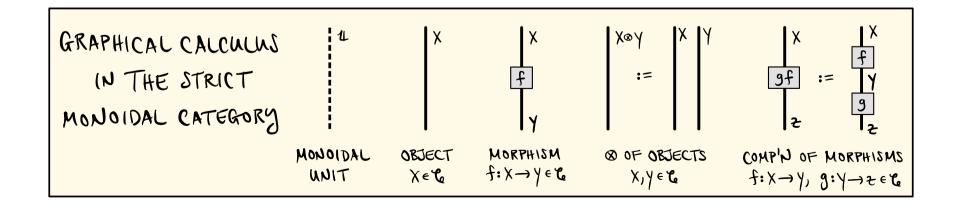
.<del>)</del>.

 $ev_{\chi}^{\chi}: \chi^{*} \otimes \chi \longrightarrow U$ 

LEFT EVALUATION

 $\operatorname{coev}_{\mathsf{X}}^{\mathsf{X}}: \mathbb{L} \to \mathsf{X} \otimes \mathsf{X}^{\mathsf{X}}$ 

LEFT COEVALUATION

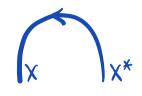


A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

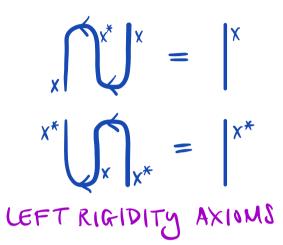


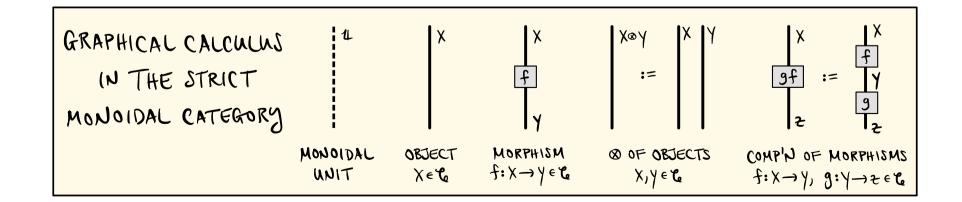


evX: X\*⊗X → L LEFT EVALUATION



COEVX: (L→X®X\*



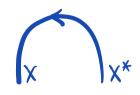


A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

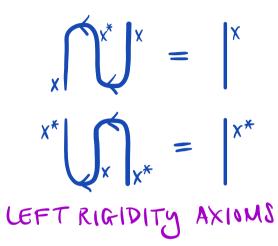
EQUIPPED WITH MORPHISMS



evX: X\*⊗X→1L LEFT EVALUATION

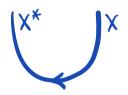


COEVX: (L→X&X\*



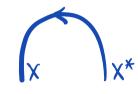
A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & C

EQUIPPED WITH MORPHISMS



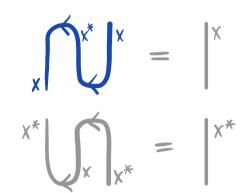
 $ev_{x}^{L}: \chi^{*} \otimes \chi \longrightarrow L$ 

LEFT EVALUATION

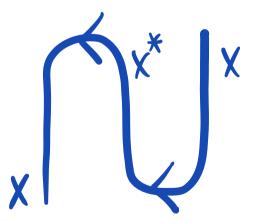


 $\operatorname{coev}_{\mathsf{X}}^{\mathsf{L}}: \mathbb{L} \to \mathsf{X} \otimes \mathsf{X}^{\mathsf{X}}$ 

LEFT COEVACUATION

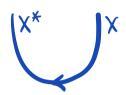


LEFT RIGIDITY AXIOMS

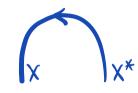


A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & C

EQUIPPED WITH MORPHISMS

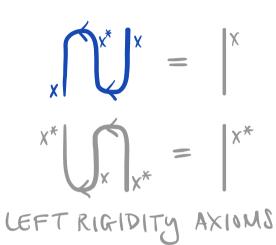


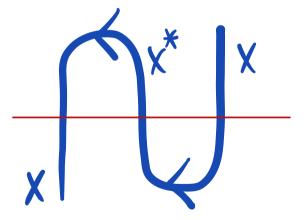
LEFT EVALUATION



 $ev_{x}^{\prime}: \chi^{*} \otimes \chi \rightarrow \mathbb{L}$   $coev_{x}^{\prime}: \mathbb{L} \rightarrow \chi \otimes \chi^{*}$ 

LEFT COEVALUATION



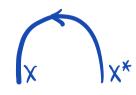


A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS



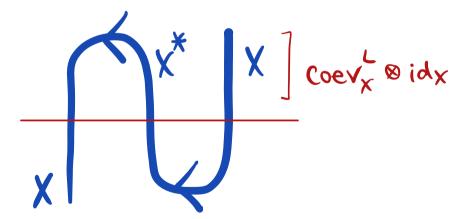
LEFT EVALUATION



 $ev_{x}^{L}: \chi^{*} \otimes \chi \longrightarrow U$   $coev_{x}^{L}: U \longrightarrow \chi \otimes \chi^{*}$ 

LEFT COEVALUATION

$$\begin{array}{ccc}
x^* & & \\
x^* & & \\
x^* & & \\
x^* & & \\
\end{array}$$
LEFT RIGIDITY AXIOMS



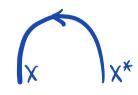
A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS



 $ev_{x}: X^{*} \otimes X \longrightarrow L$ 

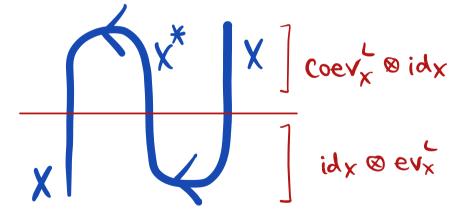
LEFT EVALUATION



 $\operatorname{coev}_{\mathsf{X}}: \mathbb{L} \to \mathsf{X} \otimes \mathsf{X}^*$ 

LEFT COEVALUATION

$$\begin{array}{ccc}
x^* & & & \\
x^* & &$$



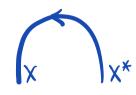
$$X = 1 \times X \xrightarrow{\text{Coev}_{\times}^{\times} \otimes \text{id}_{\times}} X \otimes X \xrightarrow{\text{id}_{\times} \otimes \text{ev}_{\times}^{\times}} X \otimes 1 = X$$

A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS

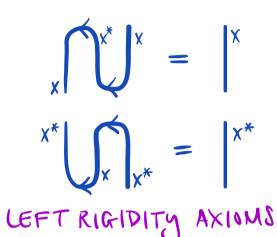


LEFT EVALUATION



 $ev_{x}^{L}: \chi^{*} \otimes \chi \longrightarrow L$   $coev_{x}^{L}: L \longrightarrow \chi \otimes \chi^{*}$ 

LEFT COEVALUATION



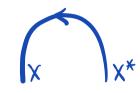
WHEN THESE EXIST, & IS CALLED LEFT RIGID.

A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS

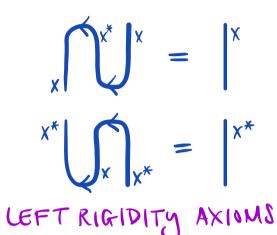


LEFT EVALUATION



 $ev_{x}^{L}: \chi^{*} \otimes \chi \longrightarrow U$   $coev_{x}^{L}: U \longrightarrow \chi \otimes \chi^{*}$ 

LEFT COEVALUATION



WHEN THESE EXIST, & IS CALLED LEFT RIGID.

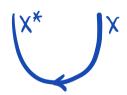
GET LEFT DUALITY FUNCTOR

$$(-)^*: \mathcal{C} \longrightarrow \mathcal{C}$$

$$\times \longmapsto \times^*$$

A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS

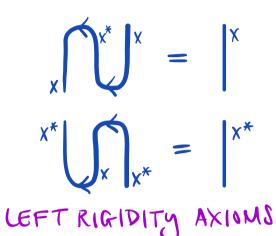


LEFT EVALUATION



 $ev_{X}^{L}: X^{*} \otimes X \longrightarrow \mathbb{L}$   $coev_{X}^{L}: \mathbb{L} \longrightarrow X \otimes X^{*}$ 

LEFT COEVALUATION



WHEN THESE EXIST, & IS CALLED LEFT RIGID.

GET LEFT DUALITY FUNCTOR

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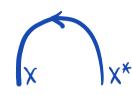
A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS



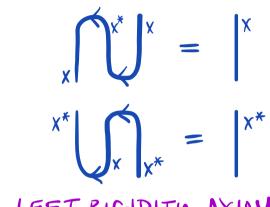
 $ev_{X}^{L}: X^{*} \otimes X \longrightarrow \mathbb{L}$ 

LEFT EVALUATION



 $\operatorname{coev}_{\mathsf{X}}^{\mathsf{L}}: \mathbb{L} \to \mathsf{X} \otimes \mathsf{X}^{\mathsf{X}}$ 

LEFT COEVACUATION



LEFT RIGIDITY AXIOMS

WHEN THESE EXIST, & IS CALLED LEFT RIGID.

GET LEFT DUALITY FUNCTOR

$$(-)^*: \mathcal{C} \longrightarrow \mathcal{C}$$

$$\times \longmapsto \times^*$$

EXERCISE 3.23

FOR  $f: X \rightarrow Y$ ,  $g: Y \rightarrow Z \in C$ 

VERIFY:

$$(gf)^* = f^*g^*$$

AS MORPHISMS 2\* -> X\*

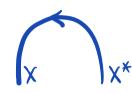
A LEFT DUAL OF AN OBJECT X & & IS AN OBJECT X\* & &

EQUIPPED WITH MORPHISMS



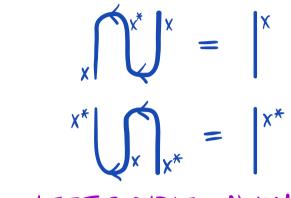
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VERIFY:

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AS MORPHISMS 2\* X\*

$$\begin{array}{ccc}
x^* & & & & \\
x^* & & \\
x^* & & & \\
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x^$$

GET CEFT DUALITY FUNCTOR

$$(-)^*: C \longrightarrow C$$

$$\times \longmapsto X^*$$

$$\downarrow^{X} \longmapsto^{F^*} := \downarrow^{Y^*} \bigvee^{X}$$

$$\downarrow^{X} \bigvee^{X} \bigvee^{X$$

EXERCISE 3.23

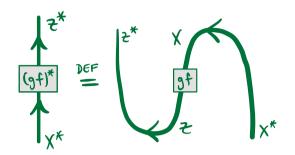
FOR 
$$f: X \rightarrow Y$$
,  $g: Y \rightarrow Z \in C$ 

VERIFY:

 $(gf)^* = f^*g^*$ 

AS MORPHISMS  $Z^* \longrightarrow X^*$ 

= LET(S DO IT =



$$\begin{array}{ccc}
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x^* & \times & \times & \times \\
x^*$$

GET LEFT DUALITY FUNCTOR

$$(-)^*: \mathcal{C} \longrightarrow \mathcal{C}$$

$$X \longleftrightarrow X^*$$

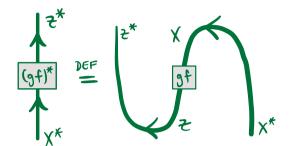
EXERCISE 3.23

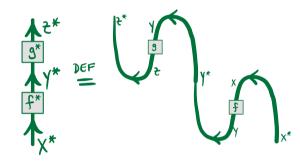
FOR  $f: X \rightarrow Y$ ,  $g: Y \rightarrow Z \in \mathcal{C}$ 

VERIFY:

$$(gf)^* = f^*g^*$$

AS MORPHISMS 2\* -> X\*





$$\begin{array}{ccc}
x^* & \times & \times & \times & \times & \times \\
x^* & \times & \times & \times & \times & \times \\
x^* & \times & \times & \times & \times & \times \\
\text{LEFT RIGIDITY AXIOMS} & & & & & & & & & \\
\end{array}$$

GET LEFT DUALITY FUNCTOR

$$(-)^*: \mathcal{C} \longrightarrow \mathcal{C}$$

$$X \longleftrightarrow X^*$$

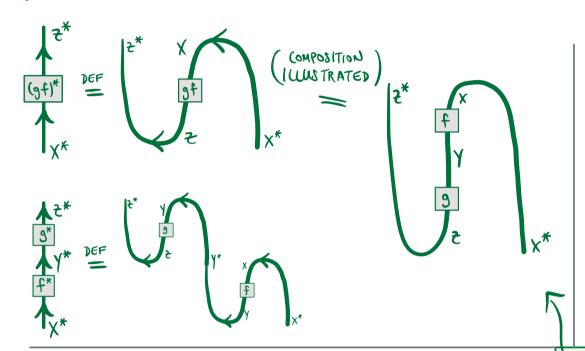
EXERCISE 3.23

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VERIFY:

$$(gf)^* = f^*g^*$$

AS MORPHISMS 2\* X\*



$$\begin{array}{ccc}
x^* & & & \\
x^* & &$$

GET LEFT DUALITY FUNCTOR

$$(-)^*: \mathcal{C} \longrightarrow \mathcal{C}$$

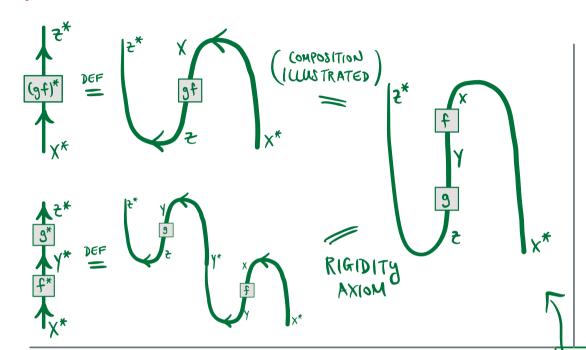
EXERCISE 3.23

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AS MORPHISMS 2\* -> X\*



$$\begin{array}{ccc}
x^* & & & \\
x^* & & \\
x^* & & & \\
x^* & & \\
x^*$$

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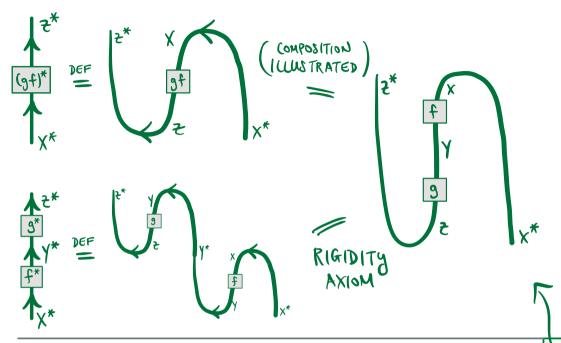
EXERCISE 3.23

FOR  $f: X \rightarrow Y$ ,  $g: Y \rightarrow Z \in C$ 

VERIFY:

$$(gf)^* = f^*g^*$$

AS MORPHISMS 2\* -> X\*



$$\begin{array}{ccc}
x^* & \times & \times & \times & \times & \times \\
x^* & \times & \times & \times & \times & \times \\
x^* & \times & \times & \times & \times & \times \\
\text{LEFT RIGIDITY AXIOMS}$$

GET LEFT DUALITY FUNCTOR

$$(-)^*: \mathcal{C} \longrightarrow \mathcal{C}$$

$$\times \longmapsto \times^*$$

EXERCISE 3.23

FOR  $f: X \rightarrow Y$ ,  $g: Y \rightarrow Z \in \mathcal{C}$ VERIFY:  $(gf)^* = f^*g^*$ AS MORPHISMS  $Z^* \longrightarrow X^*$ 

MATH 466/566 SPRING 2024

MORE)

CHELSEA WALTON RICE U.

LECTURE #14

# TOPICS:

I. STRICTIFICATION (§3.4.1)

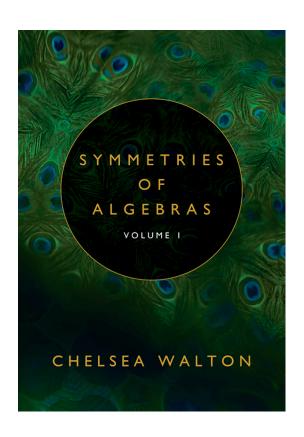
I. COHERENCE (83.4.3)

III. GRAPHICAL CALCULUS (§3.5)

→ IV. RIGID CATEGORIES (53.6)

# Enjoy this lecture? You'll enjoy the textbook!

### C. Walton's "Symmetries of Algebras, Volume 1" (2024)



**Available for purchase at:** 

619 Wreath (at a discount)

https://www.619wreath.com/

Also on Amazon & Google Play

<u>Lecture #14 keywords</u>: Coherence Theorem, dual of an object, graphical calculus, rigidity axioms, rigid category, rigid object, Strictification Theorem