

MATH 466/566  
SPRING 2024

CHELSEA WALTON  
RICE U.

LAST TIME

- CATEGORIES
- UNIV. CONSTRUCTIONS:  
I, T, O, U, M

LECTURE #7

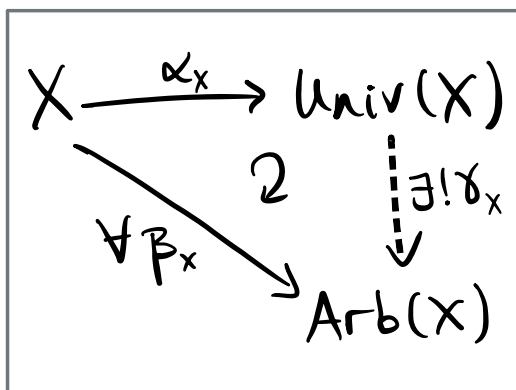
TOPICS:

I. UNIVERSAL CONSTRUCTIONS (§2.2.1)

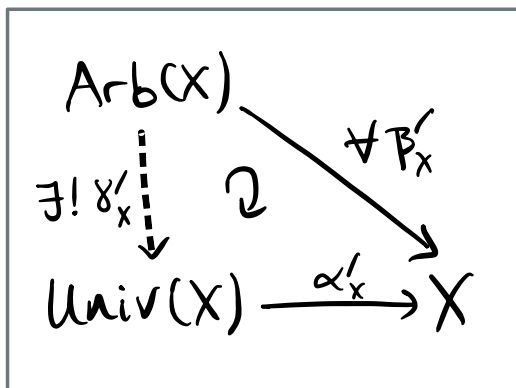
II. ABELIAN CATEGORIES (F2.2.2)

# I. UNIVERSAL CONSTRUCTIONS : REVIEW

## UNIVERSAL PROPERTY



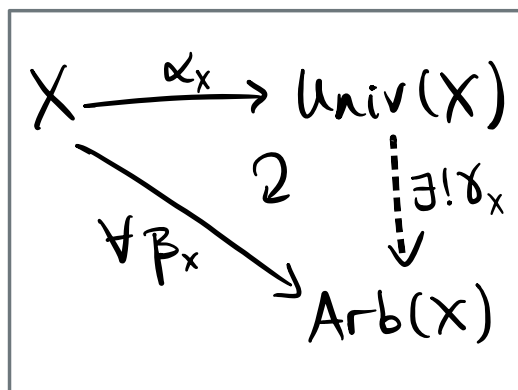
FORM I



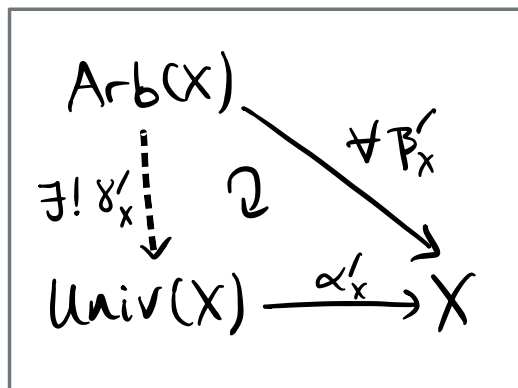
FORM II

# I. UNIVERSAL CONSTRUCTIONS : REVIEW

## UNIVERSAL PROPERTY



FORM I

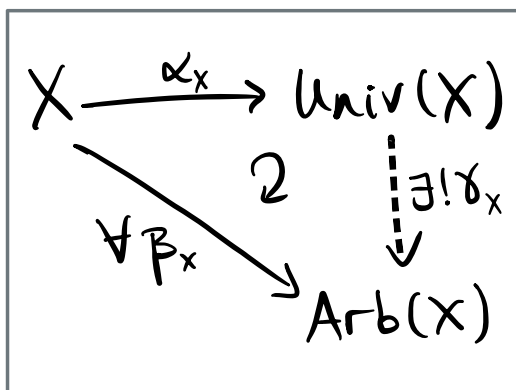


FORM II

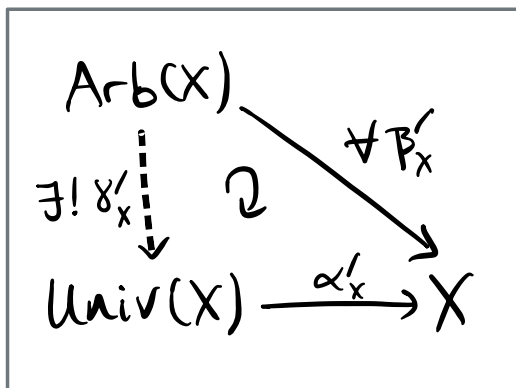
Univ(X)  
DOESN'T HAVE  
TO EXIST.  
IF EXISTS,  
THEN  
IT'S UNIQUE  
UP TO ISO.

# I. UNIVERSAL CONSTRUCTIONS : REVIEW

## UNIVERSAL PROPERTY



FORM I



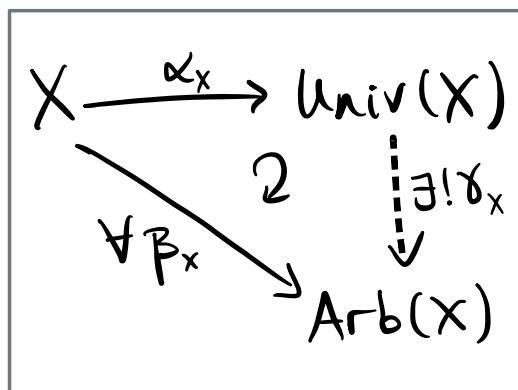
FORM II

Univ(X)  
DOESN'T HAVE  
TO EXIST.  
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THEN  
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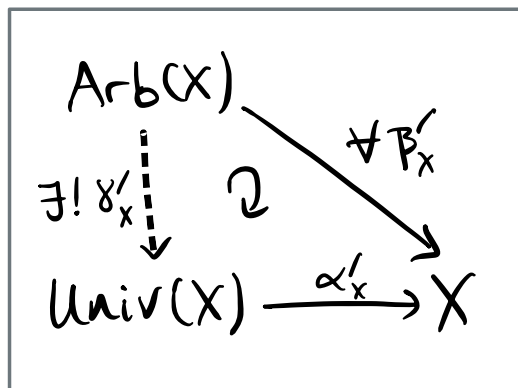
WE DON'T HAVE "ELEMENTS"  
TO WORK WITH IN GENERAL  
THIS IS THE MAIN WAY  
WE'LL DO COMPUTATIONS

# I. UNIVERSAL CONSTRUCTIONS : REVIEW

## UNIVERSAL PROPERTY



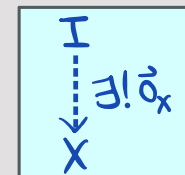
FORM I



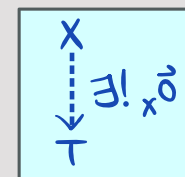
FORM II

GIVEN A CATEGORY  $\mathcal{C}$  :

AN OBJECT  $I \in \mathcal{C}$  IS INITIAL IF  
 $\forall X \in \mathcal{C} \exists!$  MORPHISM  $\vec{\delta}_x : I \rightarrow X$ .



AN OBJECT  $T \in \mathcal{C}$  IS TERMINAL IF  
 $\forall X \in \mathcal{C} \exists!$  MORPHISM  ${}_x\vec{\delta} : X \rightarrow T$ .



A ZERO OBJECT  $0$  IS AN INITIAL & TERMINAL OBJ.

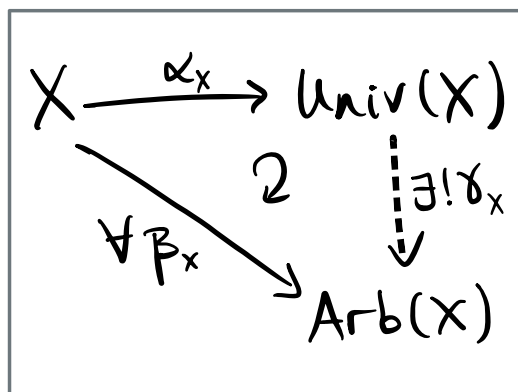
### EXAMPLES

	Set	Group	Ring	Vec
$I$	$\emptyset$	$\{e\}$	$\mathbb{Z}$	$0_{vs}$
$T$	$\{ \cdot \}$	$\{e\}$	$0_{RING}$	$0_{vs}$
$0$	N/A	$\{e\}$	N/A	$0_{vs}$

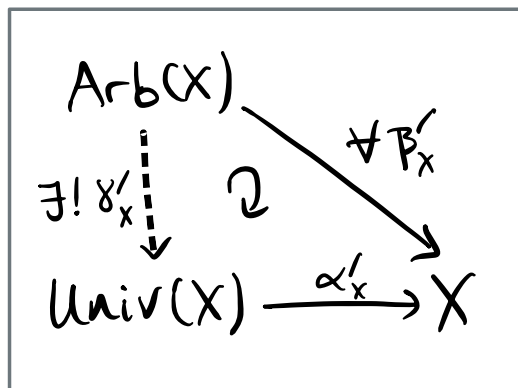
# I. UNIVERSAL CONSTRUCTIONS : REVIEW

## UNIVERSAL PROPERTY

GIVEN A CATEGORY  $\mathcal{C}$  & OBJECTS  $X, Y \in \mathcal{C}$  :



FORM I



FORM II

<u>EXAMPLES</u>	Set	Group	Ring	Vec
$\sqcup$	$(+)$ DISJOINT UNION	$\otimes$ FREE PRODUCT	$\otimes$ FREE PRODUCT	$\oplus$ DIRECT SUM
$\sqcap$	$\times$ CARTESIAN PRODUCT	$\times$ DIRECT PRODUCT	$\times$ DIRECT PRODUCT	$\times$ DIRECT PRODUCT

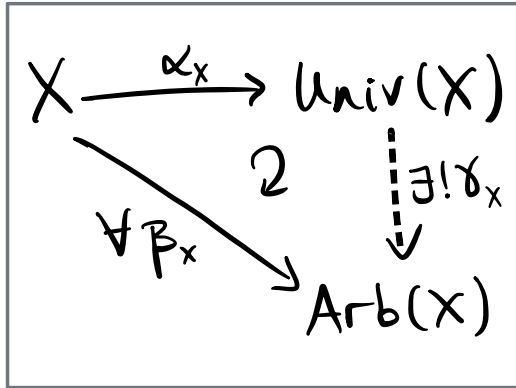
**COPRODUCT OF X AND Y**  
IS AN OBJECT  $X \sqcup Y \in \mathcal{C}$   
EQUIPPED WITH MORPHISMS  
 $\alpha_X: X \rightarrow X \sqcup Y$  AND  $\alpha_Y: Y \rightarrow X \sqcup Y$   
SATISFYING:

**PRODUCT OF X AND Y**  
IS AN OBJECT  $X \sqcap Y \in \mathcal{C}$   
EQUIPPED WITH MORPHISMS  
 $\alpha'_X: X \sqcap Y \rightarrow X$  AND  $\alpha'_Y: X \sqcap Y \rightarrow Y$   
SATISFYING:

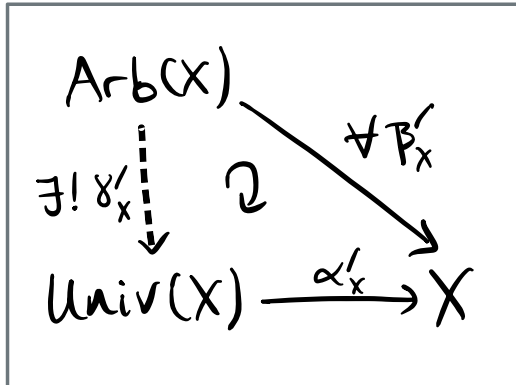
# I. UNIVERSAL CONSTRUCTIONS : REVIEW

## UNIVERSAL PROPERTY

GIVEN A CATEGORY  $\mathcal{C}$  & OBJECTS  $X, Y \in \mathcal{C}$  :

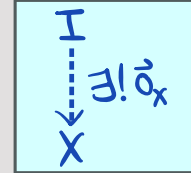


FORM I

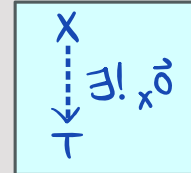


FORM II

AN OBJECT  $I \in \mathcal{C}$  IS INITIAL IF  $\forall X \in \mathcal{C} \exists!$  MORPHISM  $\delta_X: I \rightarrow X$ .



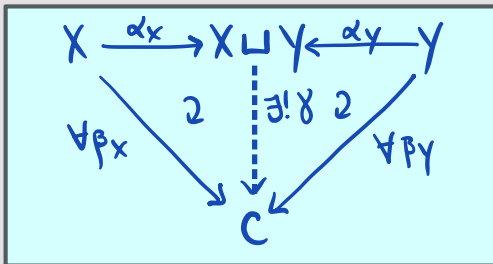
AN OBJECT  $T \in \mathcal{C}$  IS TERMINAL IF  $\forall X \in \mathcal{C} \exists!$  MORPHISM  $\delta_X: X \rightarrow T$ .



A ZERO OBJECT  $0$  IS AN INITIAL & TERMINAL OBJ.

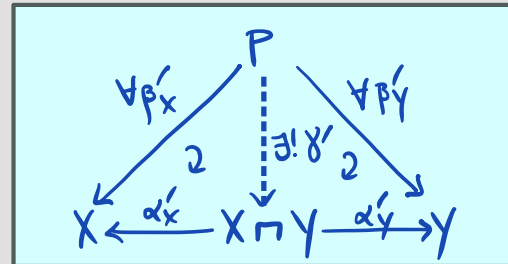
COPRODUCT OF  $X$  AND  $Y$  IS AN OBJECT  $X \sqcup Y \in \mathcal{C}$

EQUIPPED WITH MORPHISMS  $\alpha_X: X \rightarrow X \sqcup Y$  AND  $\alpha_Y: Y \rightarrow X \sqcup Y$  SATISFYING:



PRODUCT OF  $X$  AND  $Y$  IS AN OBJECT  $X \sqcap Y \in \mathcal{C}$

EQUIPPED WITH MORPHISMS  $\alpha'_X: X \sqcap Y \rightarrow X$  AND  $\alpha'_Y: X \sqcap Y \rightarrow Y$  SATISFYING:



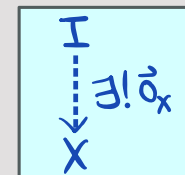
# I. UNIVERSAL CONSTRUCTIONS : REVIEW

GIVEN A CATEGORY  $\mathcal{C}$  & OBJECTS  $X, Y \in \mathcal{C}$  :

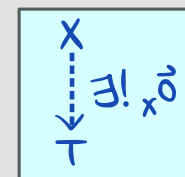
"STARTING" & "ENDING" OBJECTS



AN OBJECT  $I \in \mathcal{C}$  IS INITIAL IF  $\forall X \in \mathcal{C} \exists!$  MORPHISM  $\overset{I}{\delta}_X : I \rightarrow X$ .



AN OBJECT  $T \in \mathcal{C}$  IS TERMINAL IF  $\forall X \in \mathcal{C} \exists!$  MORPHISM  ${}_X\overset{\delta}{\delta} : X \rightarrow T$ .



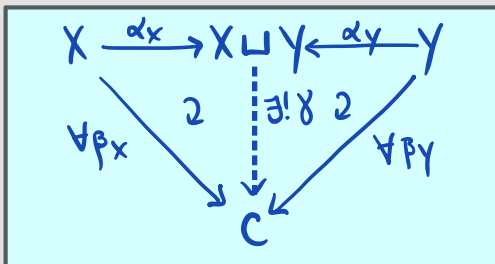
A ZERO OBJECT  $0$  IS AN INITIAL & TERMINAL OBJ.

WAYS OF COMBINING OBJECTS



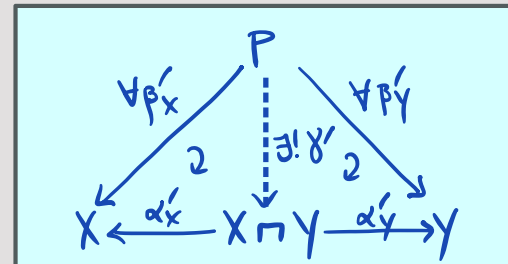
COPRODUCT OF  $X$  AND  $Y$  IS AN OBJECT  $X \sqcup Y \in \mathcal{C}$

EQUIPPED WITH MORPHISMS  $\alpha_x : X \rightarrow X \sqcup Y$  AND  $\alpha_y : Y \rightarrow X \sqcup Y$  SATISFYING:



PRODUCT OF  $X$  AND  $Y$  IS AN OBJECT  $X \sqcap Y \in \mathcal{C}$

EQUIPPED WITH MORPHISMS  $\alpha'_x : X \sqcap Y \rightarrow X$  AND  $\alpha'_y : X \sqcap Y \rightarrow Y$  SATISFYING:



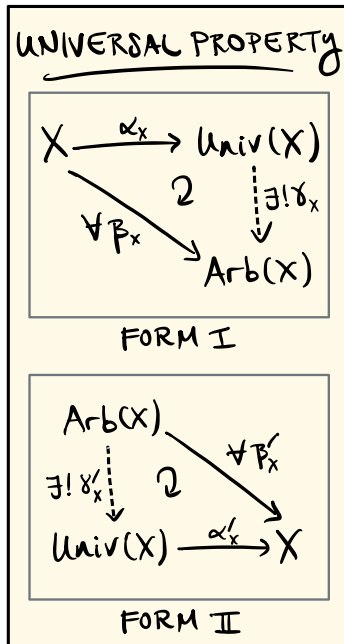
MORE UNIVERSAL CONSTRUCTIONS NEXT—



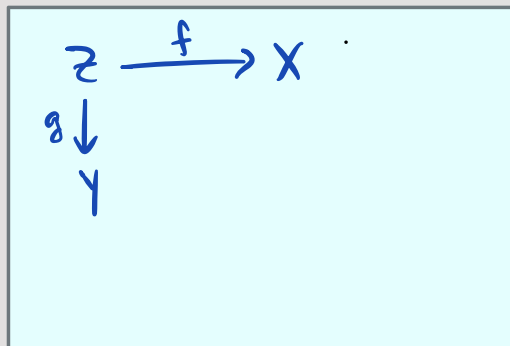
# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



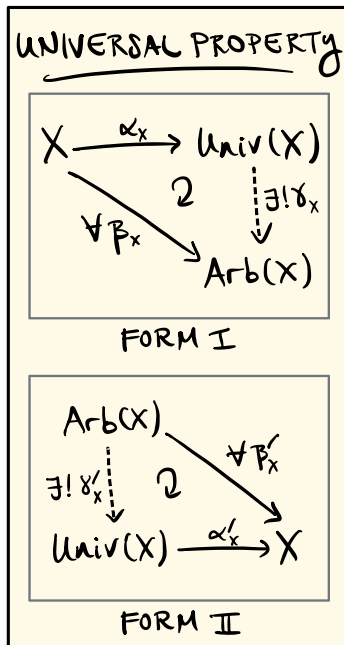
PUSHOUT OF  $f: z \rightarrow X$  &  $g: z \rightarrow Y$



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



PUSHOUT OF  $f: z \rightarrow X$  &  $g: z \rightarrow Y$   
IS AN OBJECT

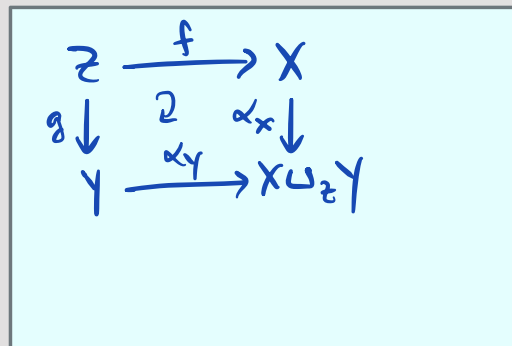
$$X \cup_z Y := X \cup_{z, f, g} Y \in \mathcal{C},$$

EQUIPPED WITH MORPHISMS

$$\alpha_x: X \rightarrow X \cup_z Y \text{ \& } \alpha_y: Y \rightarrow X \cup_z Y$$

$$\text{WHERE } \alpha_x f = \alpha_y g,$$

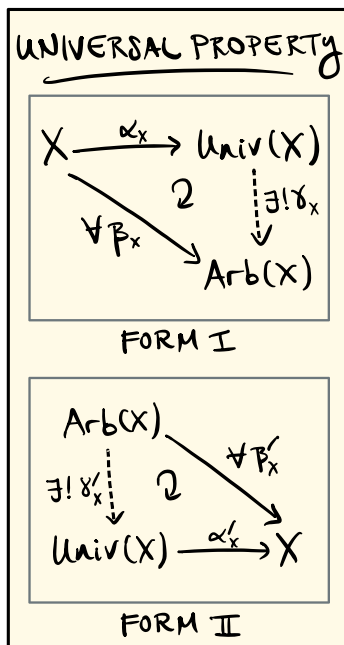
SATISFYING



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

↑ OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



**PUSHOUT** OF  $f: z \rightarrow X$  &  $g: z \rightarrow Y$   
IS AN OBJECT

$$X \cup_z Y := X \cup_{z, f, g} Y \in \mathcal{C},$$

EQUIPPED WITH MORPHISMS

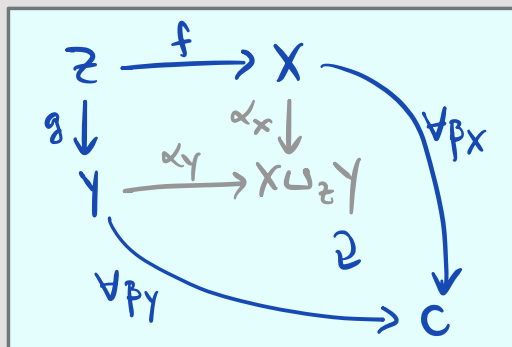
$$\alpha_x: X \rightarrow X \cup_z Y \quad \& \quad \alpha_y: Y \rightarrow X \cup_z Y$$

$$\text{WHERE } \alpha_x f = \alpha_y g,$$

SATISFYING

$$\forall \begin{cases} \beta_x: X \rightarrow C \\ \beta_y: Y \rightarrow C \end{cases} \text{ WHERE } \beta_x f = \beta_y g,$$

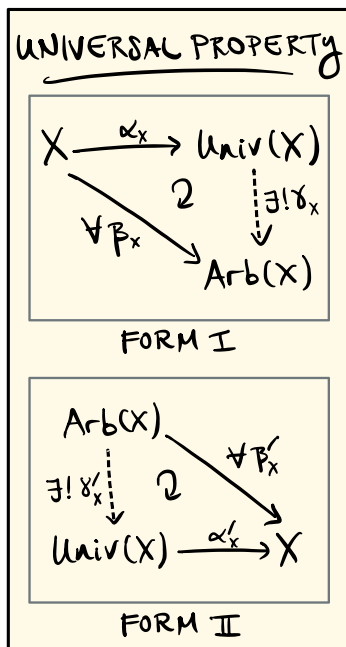
WE GET:



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

↑ OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



PUSHOUT OF  $f: z \rightarrow X$  &  $g: z \rightarrow Y$   
IS AN OBJECT

$$X \cup_z Y := X \cup_{z, f, g} Y \in \mathcal{C},$$

EQUIPPED WITH MORPHISMS

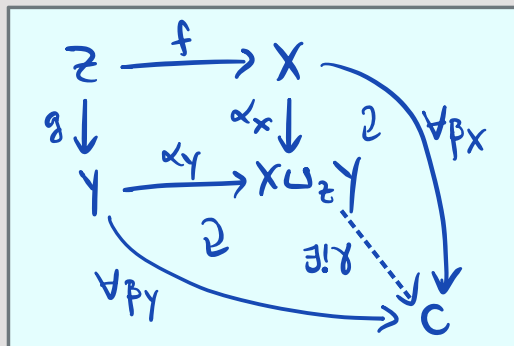
$$\alpha_x: X \rightarrow X \cup_z Y \quad \& \quad \alpha_y: Y \rightarrow X \cup_z Y$$

$$\text{WHERE } \alpha_x f = \alpha_y g,$$

SATISFYING

$$\forall \begin{cases} \beta_x: X \rightarrow C \\ \beta_y: Y \rightarrow C \end{cases} \text{ WHERE } \beta_x f = \beta_y g,$$

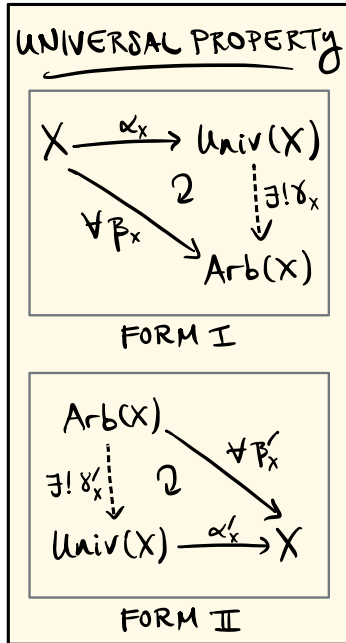
WE GET:



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

↑ OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



**PUSHOUT** OF  $f: z \rightarrow X$  &  $g: z \rightarrow Y$   
IS AN OBJECT

$$X \cup_z Y := X \cup_{z, f, g} Y \in \mathcal{C},$$

EQUIPPED WITH MORPHISMS

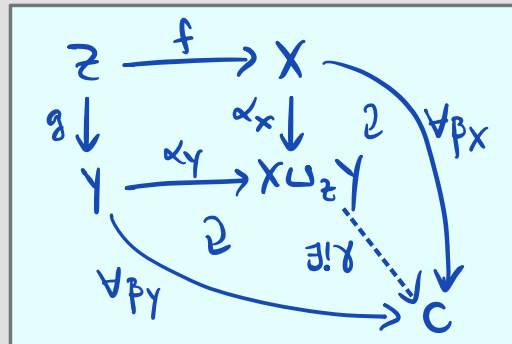
$$\alpha_x: X \rightarrow X \cup_z Y \text{ \& \ } \alpha_y: Y \rightarrow X \cup_z Y$$

$$\text{WHERE } \alpha_x f = \alpha_y g,$$

SATISFYING

$$\forall \left\{ \begin{array}{l} \beta_x: X \rightarrow C \\ \beta_y: Y \rightarrow C \end{array} \right. \text{ WHERE } \beta_x f = \beta_y g,$$

WE GET:



**PULLBACK** OF  $f: X \rightarrow z$  &  $g: Y \rightarrow z$   
IS AN OBJECT

$$X \cap_z Y := X \cap_{z, f, g} Y \in \mathcal{C},$$

EQUIPPED WITH MORPHISMS

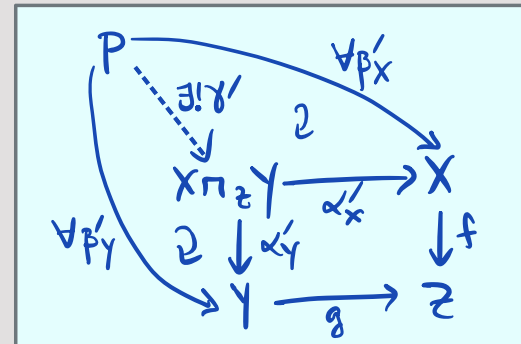
$$\alpha'_x: X \cap_z Y \rightarrow X \text{ \& \ } \alpha'_y: X \cap_z Y \rightarrow Y$$

$$\text{WHERE } f \alpha'_x = g \alpha'_y,$$

SATISFYING

$$\forall \left\{ \begin{array}{l} \beta'_x: P \rightarrow X \\ \beta'_y: P \rightarrow Y \end{array} \right. \text{ WHERE } f \beta'_x = g \beta'_y,$$

WE GET:



"PULLBACKS"  
IN  $\mathcal{C}$

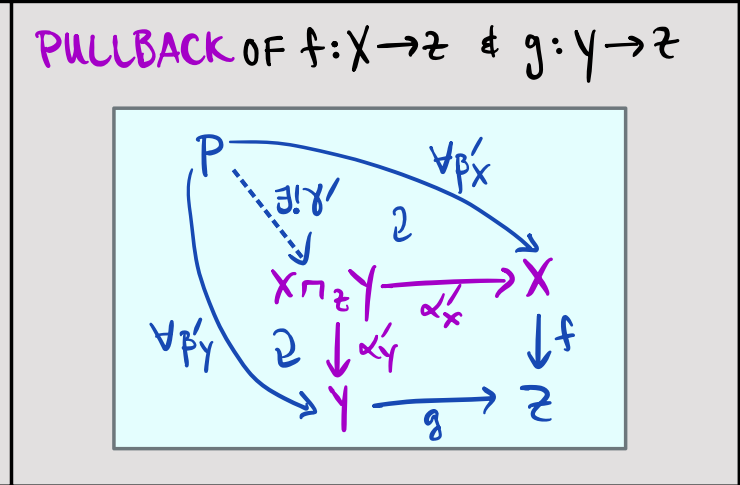
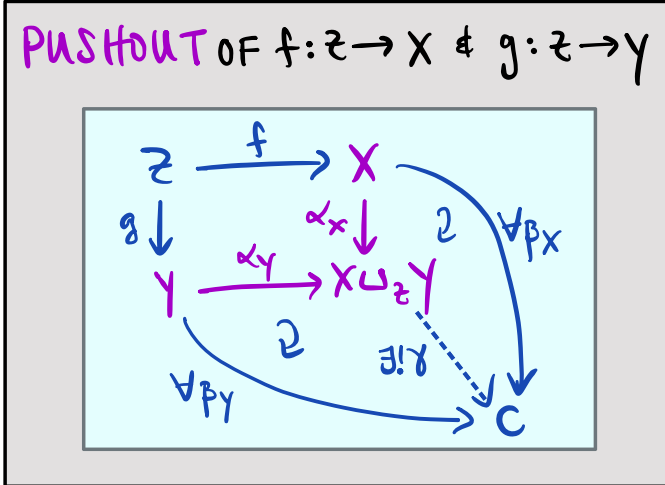
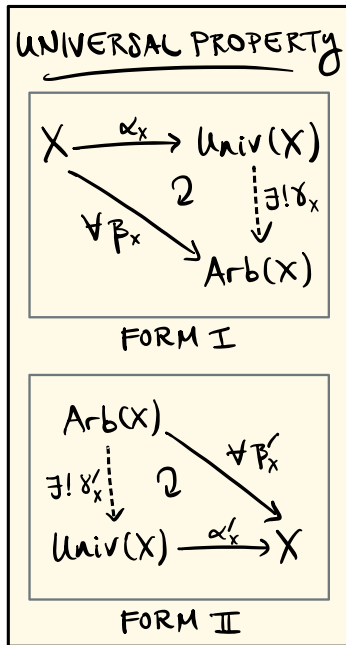
III

PUSHOUTS  
IN  $\mathcal{C}^{op}$

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



↑  
OF "FORM I"

↑  
OF "FORM II"

"PULLBACKS"  
IN  $\mathcal{C}$

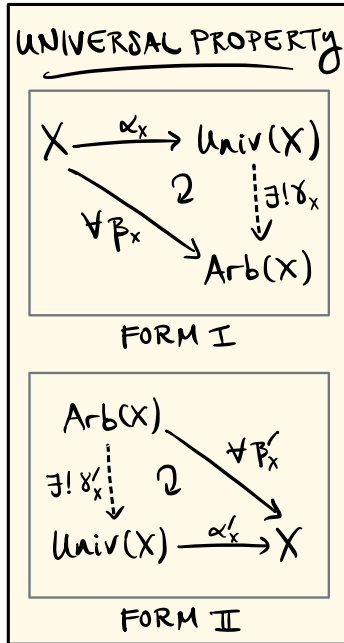
|||

PUSHOUTS  
IN  $\mathcal{C}^{op}$

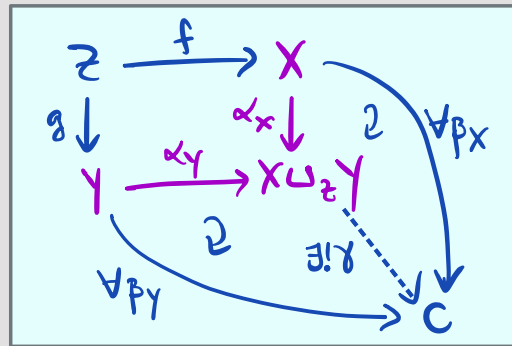
# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

↑ OPERATION ON MORPHISMS

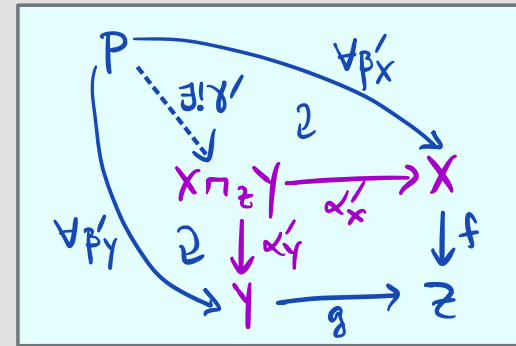
GIVEN A CATEGORY  $\mathcal{C}$ :



**PUSHOUT** OF  $f: z \rightarrow X$  &  $g: z \rightarrow Y$



**PULLBACK** OF  $f: X \rightarrow z$  &  $g: Y \rightarrow z$



EXERCISE 2.9 FOR  $X, Y, z \in \text{Set}$ , WE GET

WITH FUNCTIONS

$$f: z \rightarrow X, g: z \rightarrow Y :$$

$$X \sqcup_z Y \cong (X \sqcup Y) / \sim$$

QUOTIENT SET OF  $X \sqcup Y$

HERE,  $f(z) \sim g(z)$  IN  $X \sqcup Y \quad \forall z \in z$

WITH FUNCTIONS

$$f: X \rightarrow z, g: Y \rightarrow z :$$

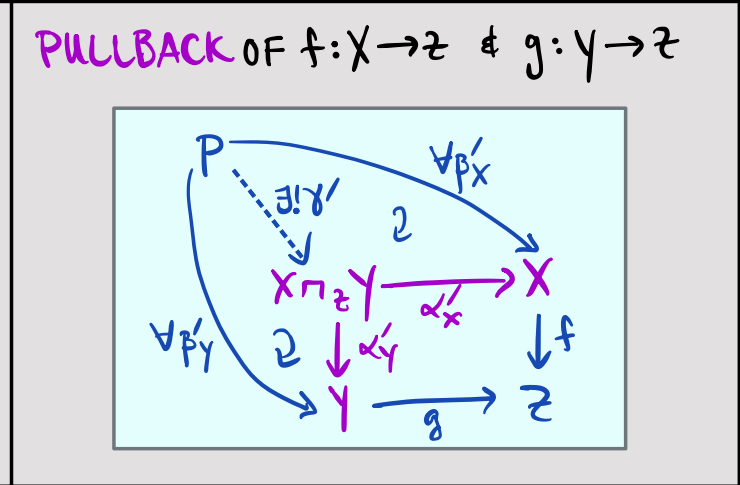
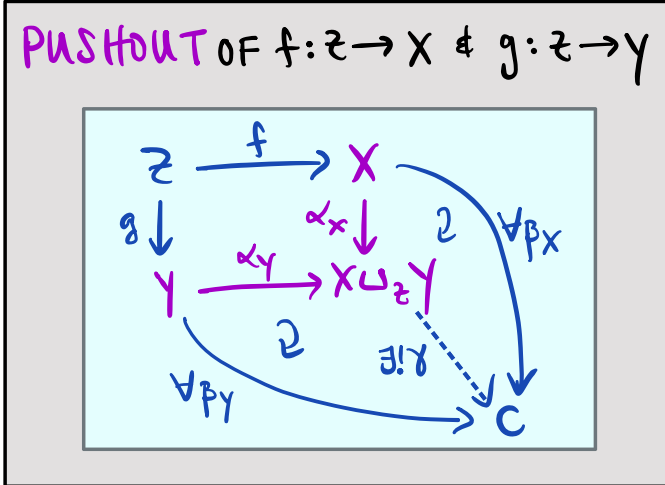
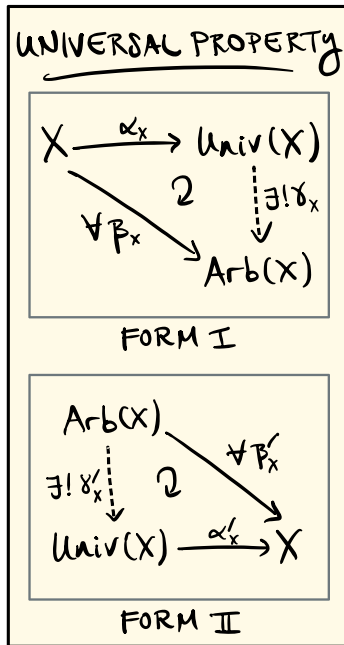
$$X \pi_z Y = \{ (x, y) \in X \times Y \mid f(x) = g(y) \text{ IN } z \}$$

SUBSET OF  $X \times Y$

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



**EXERCISE 2.9** FOR  $X, Y, z \in \text{Set}$ , WE GET

<p>WITH FUNCTIONS  <math>f: z \rightarrow X, g: z \rightarrow Y</math>:</p> $X \sqcup_z Y \cong (X \uplus Y) / \sim$ <p>QUOTIENT SET OF <math>X \uplus Y</math></p> <p>HERE, <math>f(z) \sim g(z)</math> IN <math>X \uplus Y \quad \forall z \in z</math></p>	<p>WITH FUNCTIONS  <math>f: X \rightarrow z, g: Y \rightarrow z</math>:</p> $X \sqcap_z Y = \{(x, y) \in X \times Y \mid f(x) = g(y) \text{ IN } z\}$ <p>SUBSET OF <math>X \times Y</math></p>
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IN Set

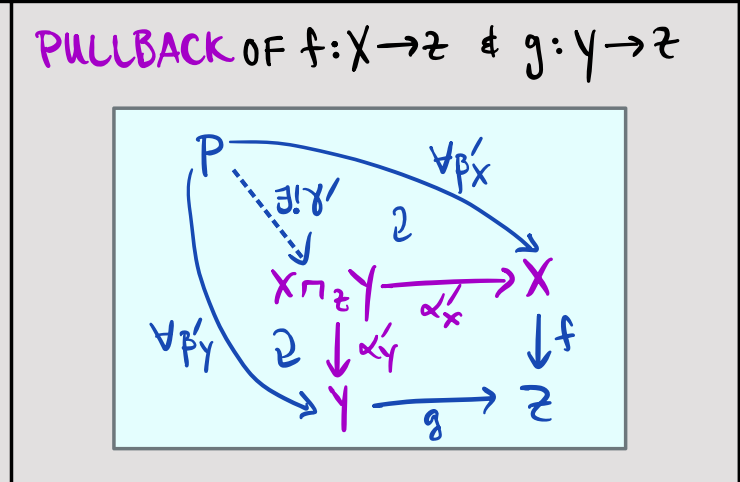
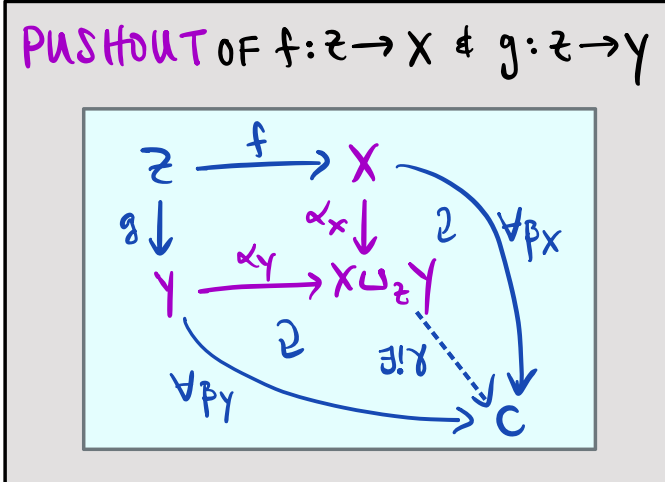
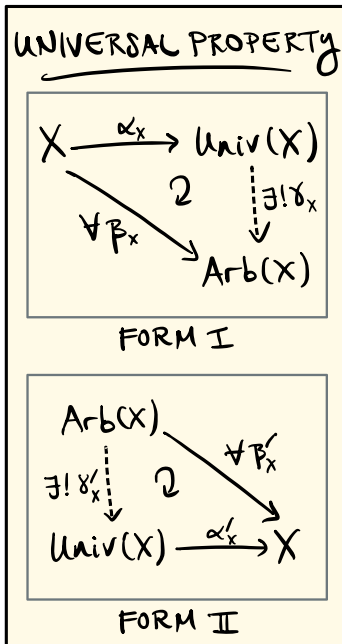
$\emptyset$	$\emptyset$
$\{ \cdot \}$	$\{ \cdot \}$
$\uplus$	$\uplus$ DISJOINT UNION
$\times$	$\times$ CARTESIAN PRODUCT



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

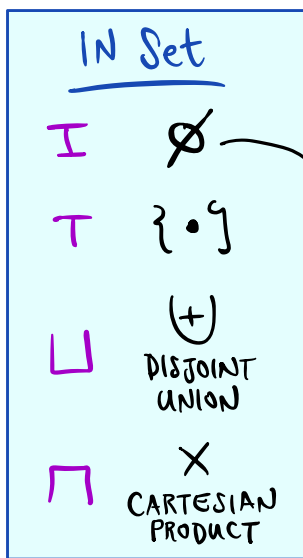
OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



**EXERCISE 2.9** FOR  $X, Y, z \in \text{Set}$ , WE GET

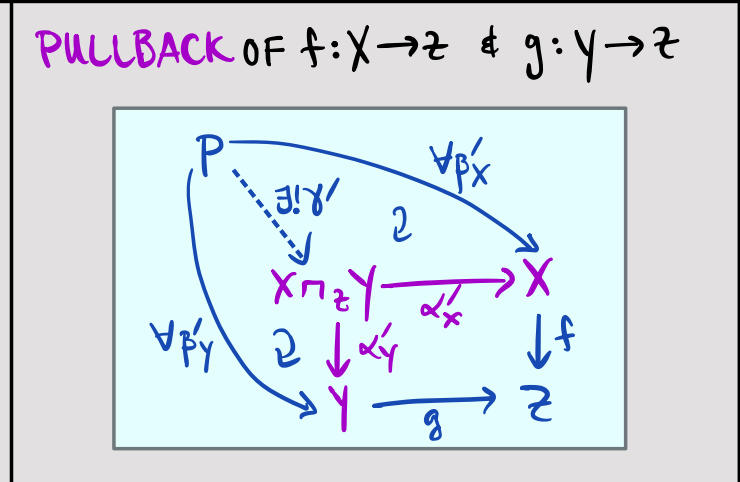
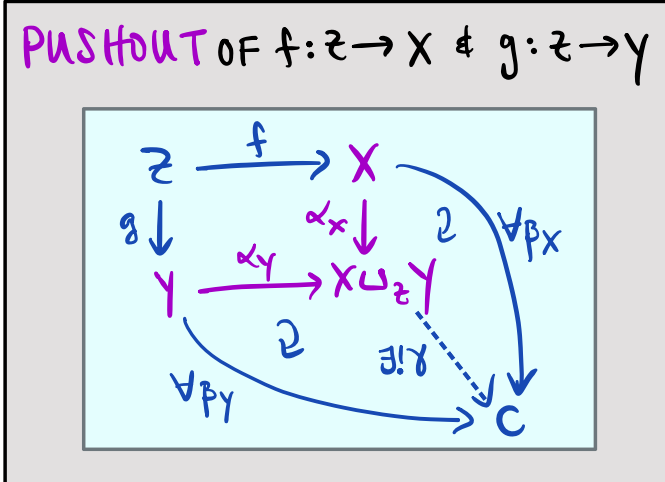
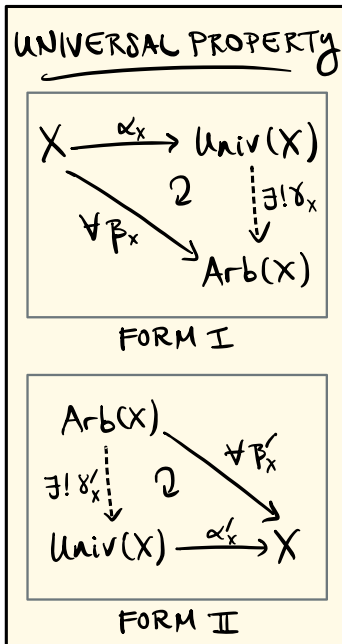
WITH FUNCTIONS $f: \emptyset \rightarrow X, g: \emptyset \rightarrow Y$ :	WITH FUNCTIONS $f: X \rightarrow z, g: Y \rightarrow z$ :
$X \sqcup_z Y \cong (X \uplus Y) / \sim$	$X \sqcap_z Y = \{(x, y) \in X \times Y \mid f(x) = g(y) \text{ in } z\}$
QUOTIENT SET OF $X \uplus Y$	SUBSET OF $X \times Y$
HERE, $f(z) \sim g(z)$ IN $X \uplus Y \quad \forall z \in z$	



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



EXERCISE 2.9 FOR  $X, Y, z \in \text{Set}$ , WE GET

<p>WITH FUNCTIONS  <math>f: \emptyset \rightarrow X, g: \emptyset \rightarrow Y</math>:</p> <p><math>X \sqcup_{\emptyset} Y \cong (X \oplus Y) / \sim \parallel</math>  <del>QUOTIENT SET OF <math>X \oplus Y</math></del></p> <p>HERE, <math>f(z) \sim g(z)</math> IN <math>X \oplus Y \forall z \in z</math></p>	<p>WITH FUNCTIONS  <math>f: X \rightarrow z, g: Y \rightarrow z</math>:</p> <p><math>X \sqcap_z Y = \{(x, y) \in X \times Y \mid f(x) = g(y) \text{ in } z\}</math>                  SUBSET OF <math>X \times Y</math></p>
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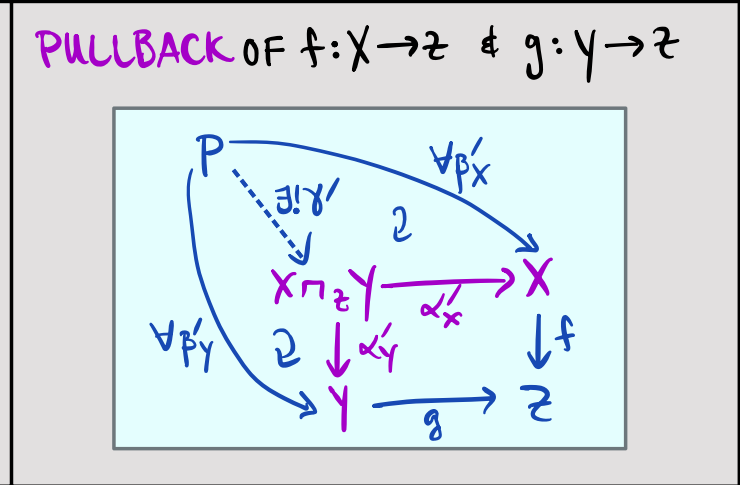
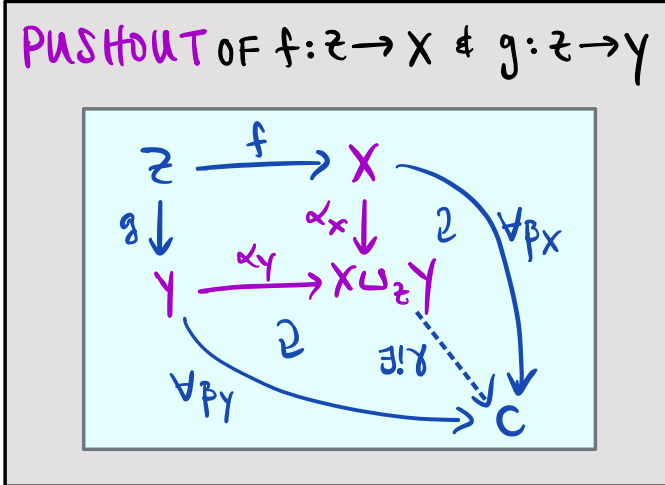
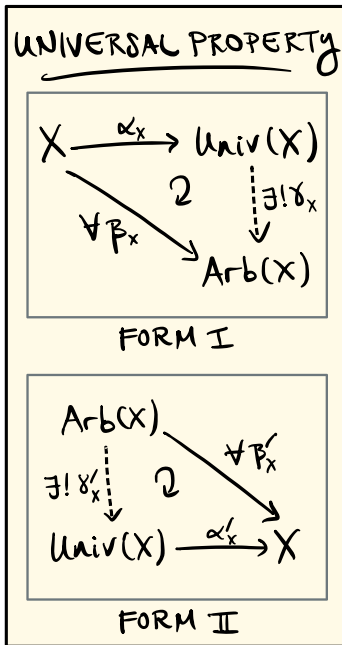
IN Set

$\emptyset$	$\emptyset$
$\{ \cdot \}$	$\{ \cdot \}$
$\sqcup$	$\oplus$ DISJOINT UNION
$\sqcap$	$\times$ CARTESIAN PRODUCT

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



**EXERCISE 2.9** FOR  $X, Y, z \in \text{Set}$ , WE GET

<p>WITH FUNCTIONS  <math>f: \emptyset \rightarrow X, g: \emptyset \rightarrow Y</math>:</p> <p><math>X \sqcup_{\emptyset} Y \cong (X \oplus Y) / \sim</math>  <del>QUOTIENT SET OF <math>X \oplus Y</math></del></p> <p>HERE, <math>f(z) \sim g(z)</math> IN <math>X \oplus Y \forall z \in z</math></p>	<p>WITH FUNCTIONS  <math>f: X \rightarrow \{ \cdot \}, g: Y \rightarrow \{ \cdot \}</math>:</p> <p><math>X \sqcap_{\{ \cdot \}} Y = \{ (x, y) \in X \times Y \mid f(x) = g(y) \text{ in } z \}</math>  <del>SUBSET OF <math>X \times Y</math></del></p>
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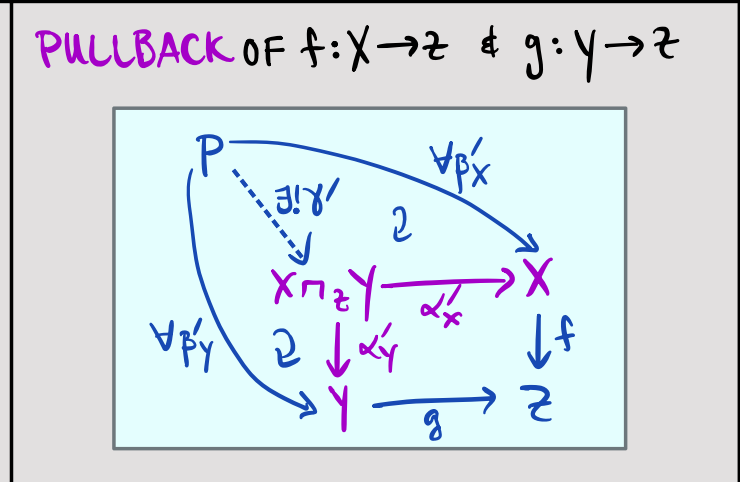
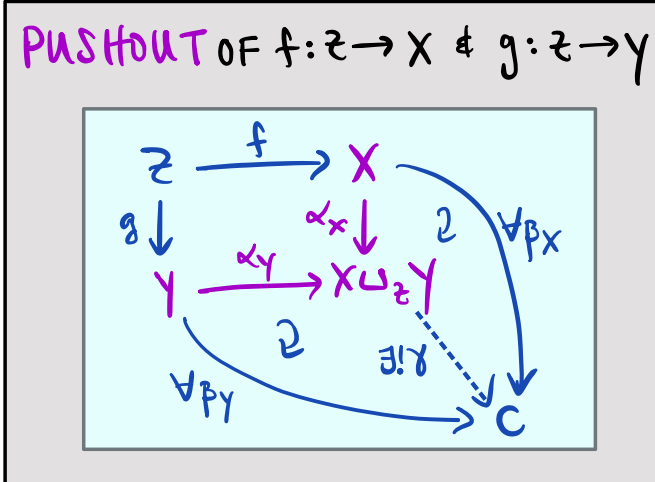
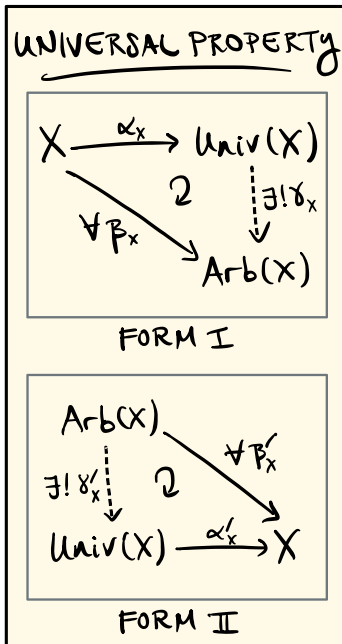
IN Set

$\emptyset$	$\emptyset$
$\{ \cdot \}$	$\{ \cdot \}$
$\sqcup$	$\oplus$ DISJOINT UNION
$\sqcap$	$\times$ CARTESIAN PRODUCT

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



EXERCISE 2.9 FOR  $X, Y, z \in \text{Set}$ , WE GET

<p>WITH FUNCTIONS  <math>f: \emptyset \rightarrow X, g: \emptyset \rightarrow Y</math>:</p> <p><math>X \sqcup_{\emptyset} Y \cong (X \oplus Y) / \sim</math></p> <p><del>QUOTIENT SET OF <math>X \oplus Y</math></del></p> <p>HERE, <math>f(z) \sim g(z)</math> IN <math>X \oplus Y \quad \forall z \in z</math></p>	<p>WITH FUNCTIONS  <math>f: X \rightarrow \{ \cdot \}, g: Y \rightarrow \{ \cdot \}</math>:</p> <p><math>X \sqcap_{\{ \cdot \}} Y = \{ (x, y) \in X \times Y \mid f(x) = g(y) \text{ in } z \}</math></p> <p><del>VACUOUS</del> <math>\rightarrow</math> <del>SUBSET OF <math>X \times Y</math></del></p>
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EXER 2.10

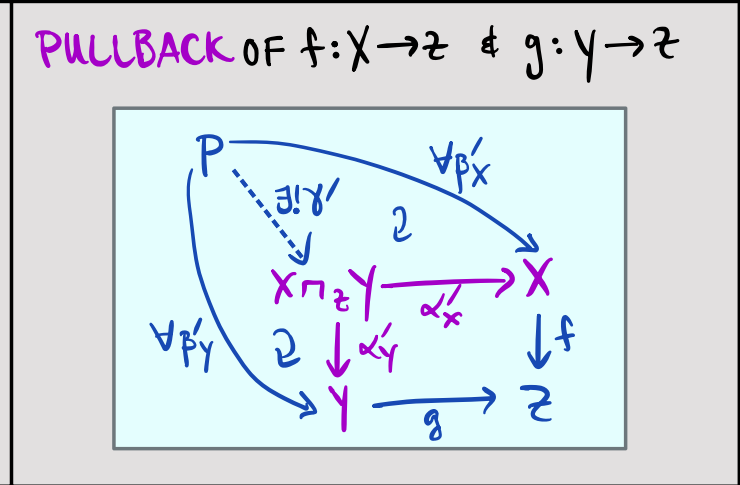
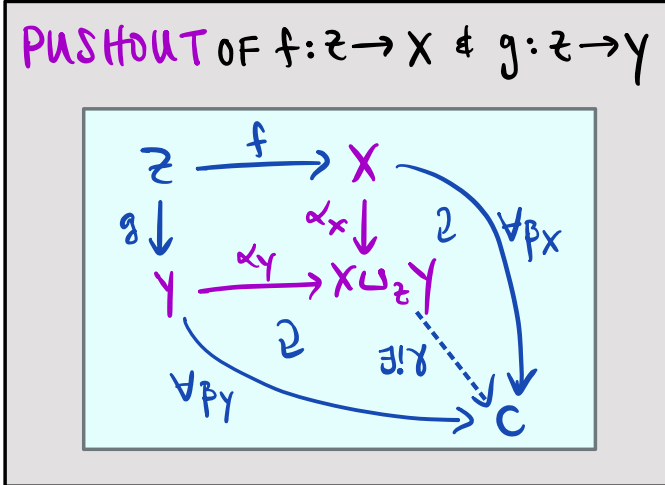
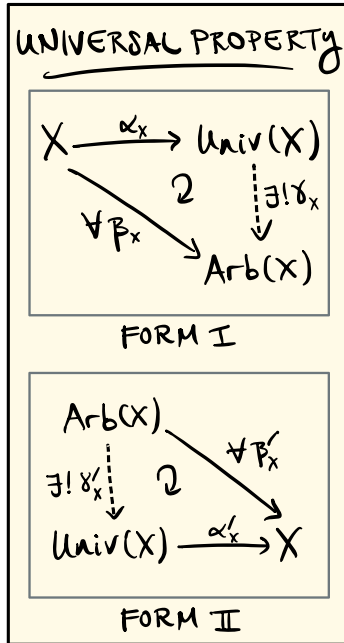
$X \sqcup_{I, \delta_x, \delta_y} Y$   
 $\cong$   
 $X \sqcup Y$

$X \sqcap_{T, \delta_x, \delta_y} Y$   
 $\cong$   
 $X \cap Y$

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



HINT FOR  $X \cup_I Y \cong X \cup Y$

EXER 2.10

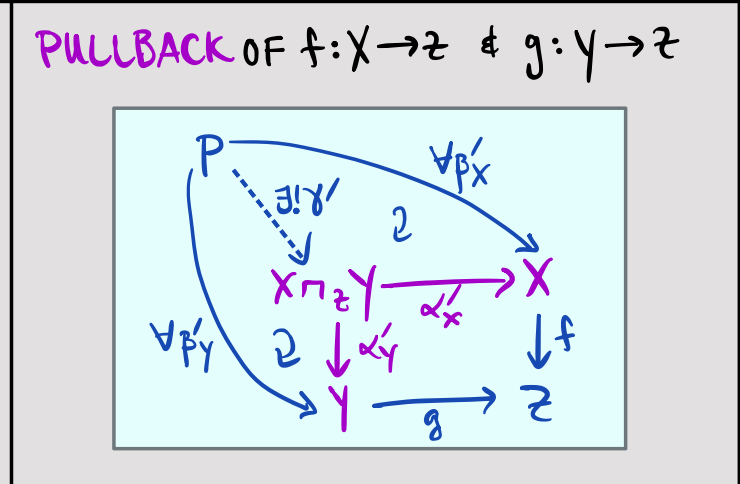
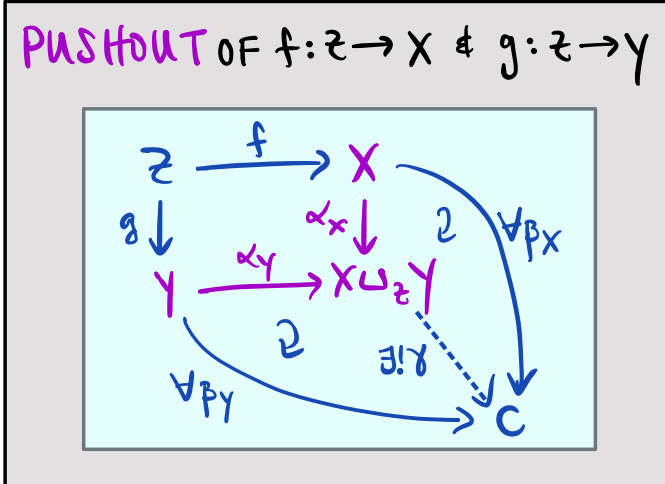
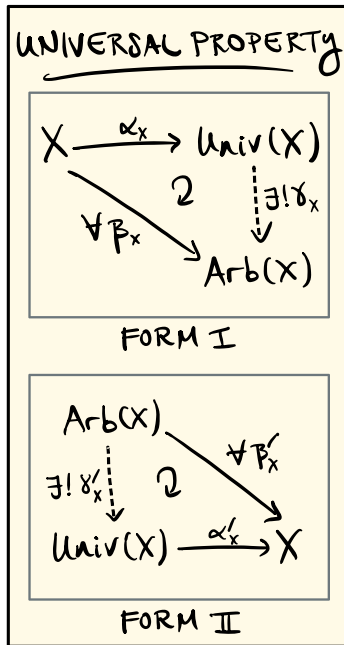
$X \cup_{I, \alpha_x, \alpha_y} Y$   
 $\cong$   
 $X \cup Y$

$X \cap_{T, \alpha_x, \alpha_y} Y$   
 $\cong$   
 $X \cap Y$

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

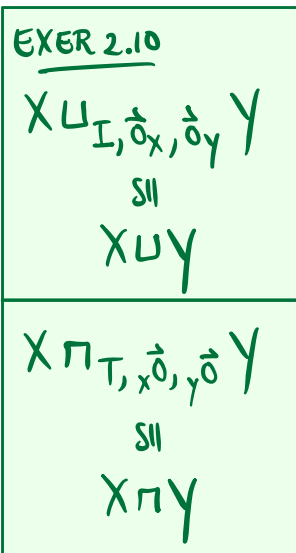
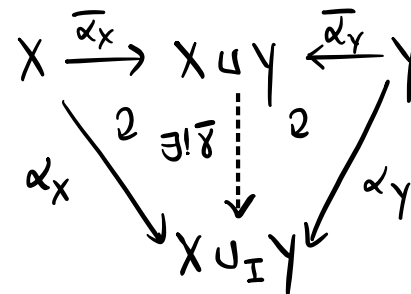
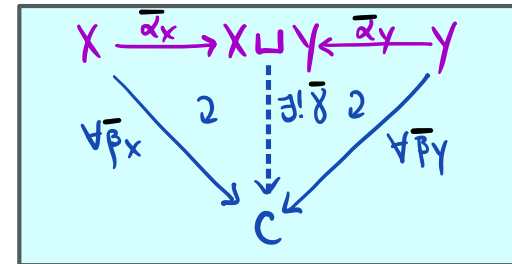
OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



HINT FOR  $X \cup_I Y \cong X \cup Y$

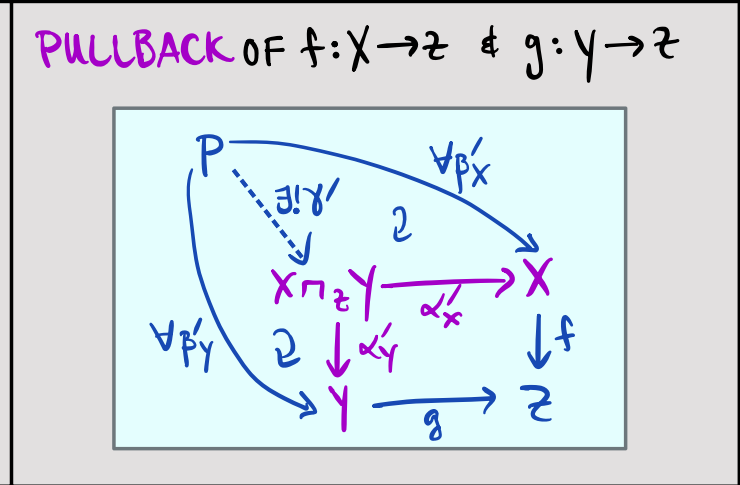
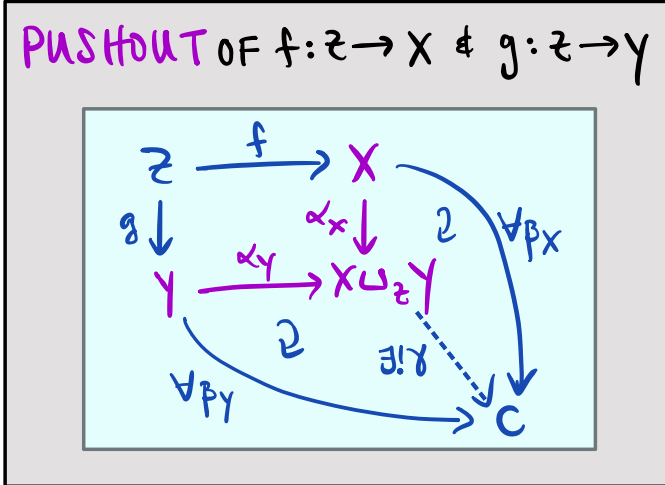
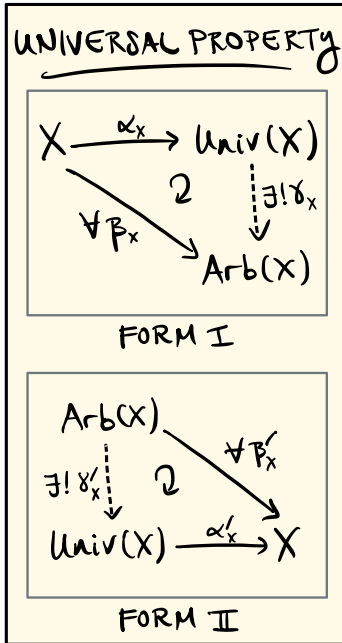
UNIV PROP OF  $X \cup Y$  YIELDS  
A MORPHISM  $\bar{\delta}: X \cup Y \rightarrow X \cup_I Y$   
VIA:



# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

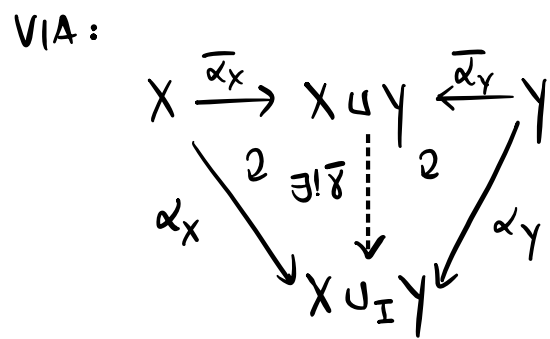
GIVEN A CATEGORY  $\mathcal{C}$ :



HINT FOR  $X \cup_I Y \cong X \cup Y$

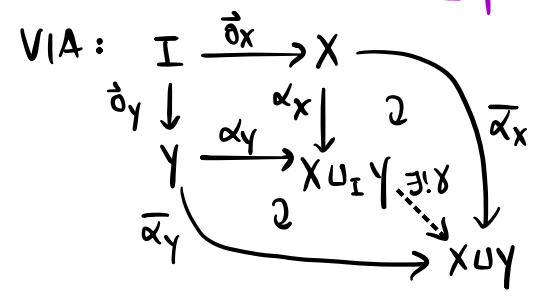
UNIV PROP OF  $X \cup Y$  YIELDS

A MORPHISM  $\bar{\gamma}: X \cup Y \rightarrow X \cup_I Y$



UNIV PROP OF  $X \cup_I Y$  YIELDS

A MORPHISM  $\gamma: X \cup_I Y \rightarrow X \cup Y$



EXER 2.10

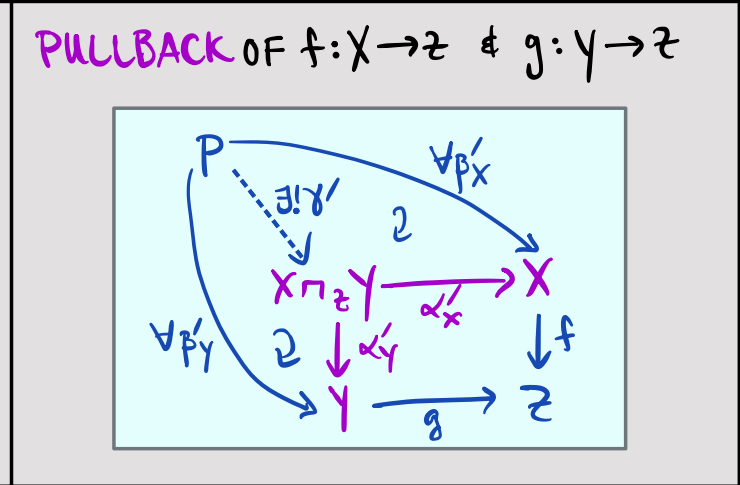
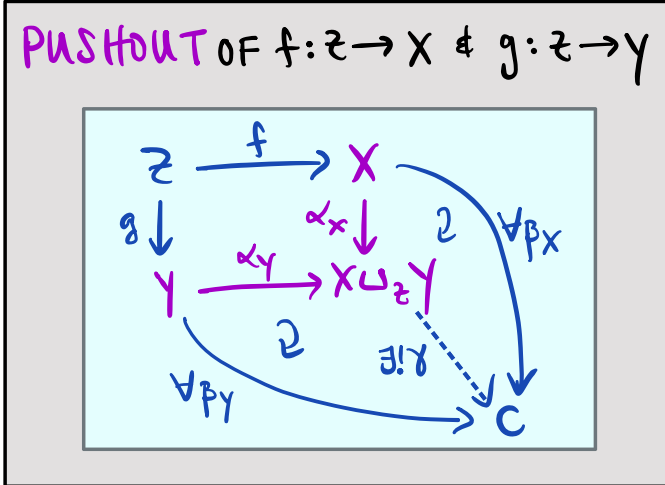
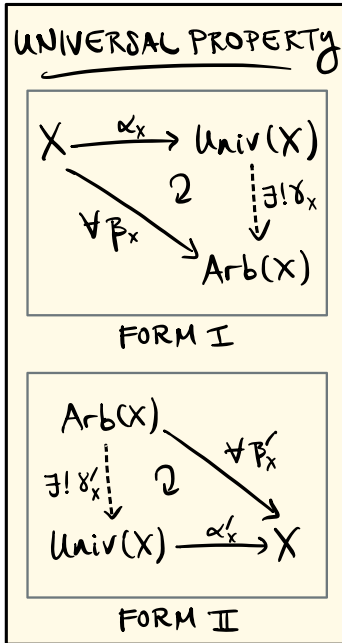
$X \cup_I, \alpha_x, \alpha_y Y$   
 $\cong$   
 $X \cup Y$

$X \cap_I, \alpha_x, \alpha_y Y$   
 $\cong$   
 $X \cap Y$

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :

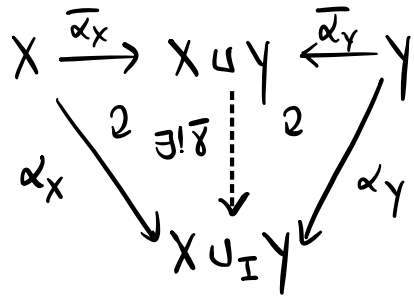


HINT FOR  $X \cup_I Y \cong X \cup Y$

UNIV PROP OF  $X \cup Y$  YIELDS

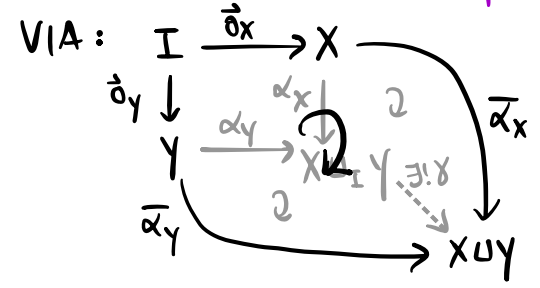
A MORPHISM  $\bar{\gamma}: X \cup Y \rightarrow X \cup_I Y$

VIA:



UNIV PROP OF  $X \cup_I Y$  YIELDS

A MORPHISM  $\bar{\gamma}: X \cup_I Y \rightarrow X \cup Y$



INDEED,  $\bar{\alpha}_x \bar{\delta}_x = \bar{\alpha}_y \bar{\delta}_y = \bar{\delta}_{X \cup Y}$

EXER 2.10

$X \cup_I, \bar{\delta}_x, \bar{\delta}_y Y$   
 $\cong$   
 $X \cup Y$

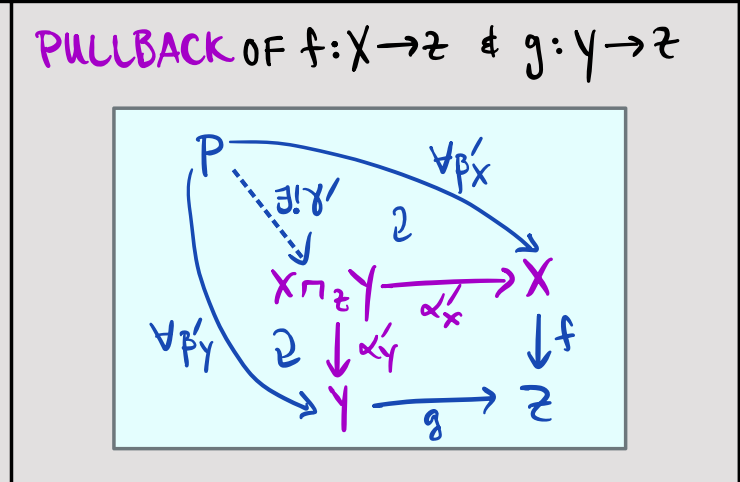
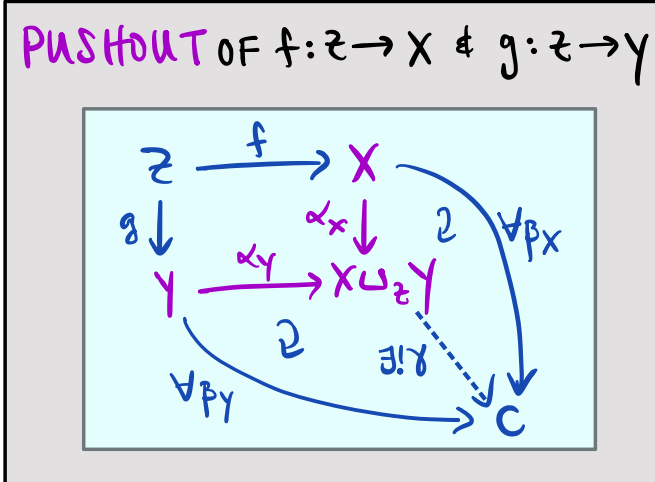
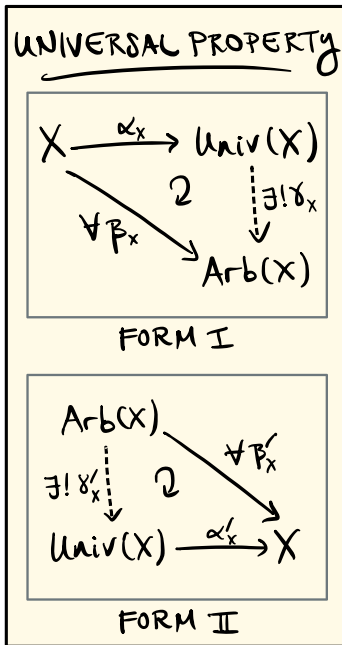
$X \cap_I, \bar{\delta}_x, \bar{\delta}_y Y$   
 $\cong$   
 $X \cap Y$



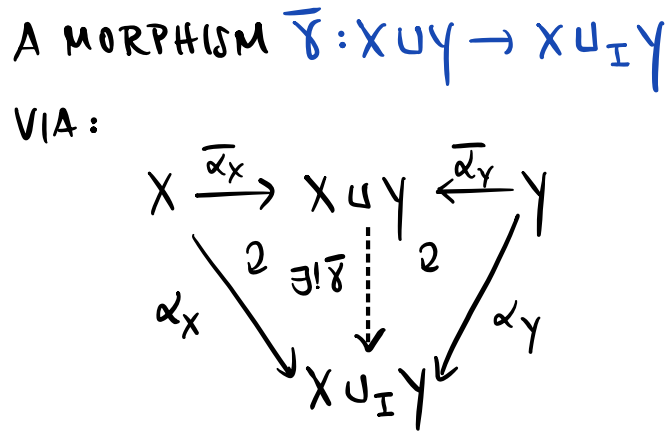
# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

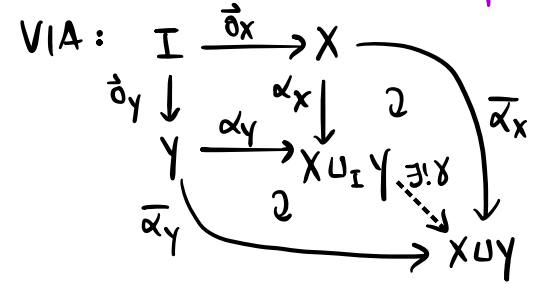
GIVEN A CATEGORY  $\mathcal{C}$ :



HINT FOR  $X \cup_I Y \cong X \cup Y$   
 UNIV PROP OF  $X \cup Y$  YIELDS  
 A MORPHISM  $\bar{\gamma}: X \cup Y \rightarrow X \cup_I Y$



UNIV PROP OF  $X \cup_I Y$  YIELDS  
 A MORPHISM  $\gamma: X \cup_I Y \rightarrow X \cup Y$



SHOW  $\gamma$  &  $\bar{\gamma}$  ARE MUTUALLY INV.

EXER 2.10

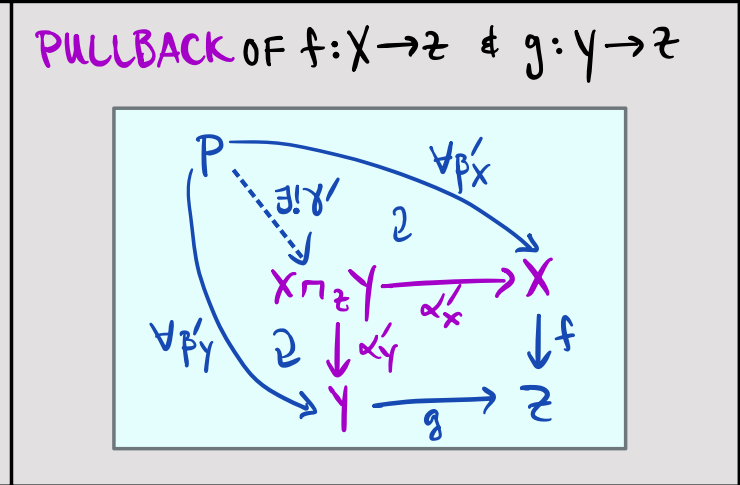
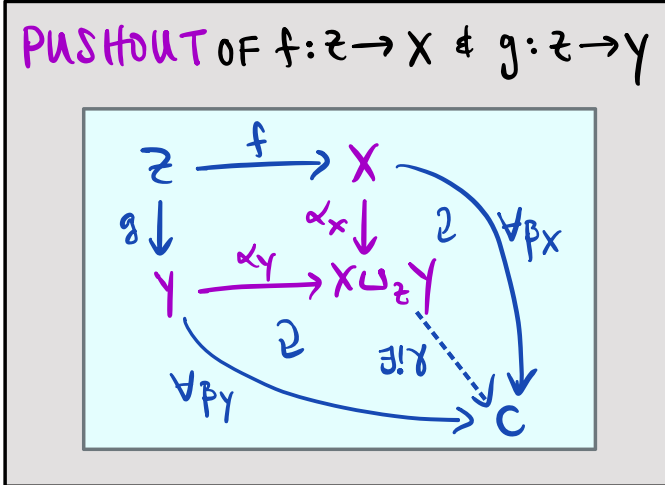
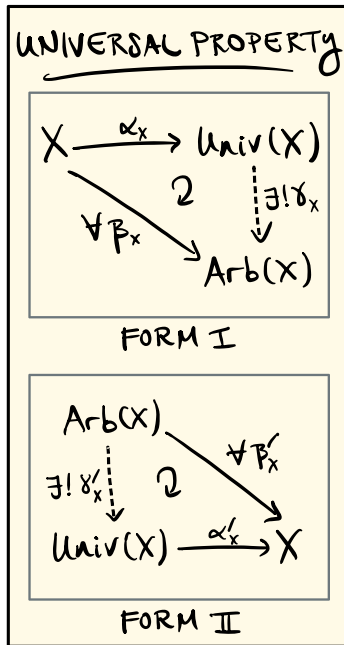
$X \cup_I, \bar{\alpha}_x, \bar{\alpha}_y Y$   
 $\cong$   
 $X \cup Y$

$X \cap_I, \bar{\alpha}_x, \bar{\alpha}_y Y$   
 $\cong$   
 $X \cap Y$

# I. UNIVERSAL CONSTRUCTIONS : PUSHOUTS AND PULLBACKS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



EXER 2.10

$X \sqcup_I Y$   
 $\parallel$   
 $X \sqcup Y$

---

$X \sqcap_T Y$   
 $\parallel$   
 $X \sqcap Y$

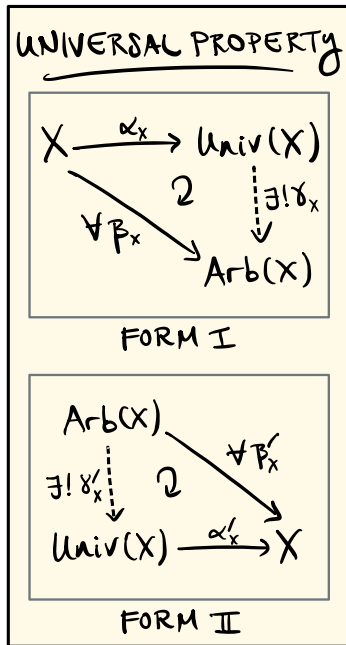
IF  $\left\{ \begin{array}{l} \text{PUSHOUTS} \\ \text{PULLBACKS} \end{array} \right\} \neq \left\{ \begin{array}{l} I \\ T \end{array} \right\}$  EXIST IN  $\mathcal{C}$ ,

SO DO  $\left\{ \begin{array}{l} \text{COPRODUCTS} \\ \text{PRODUCTS} \end{array} \right\}$ .

# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

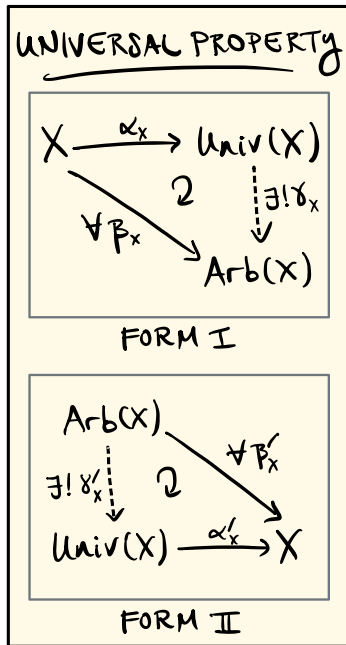
OPERATION ON PARALLEL MORPHISMS



# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

OPERATION ON PARALLEL MORPHISMS

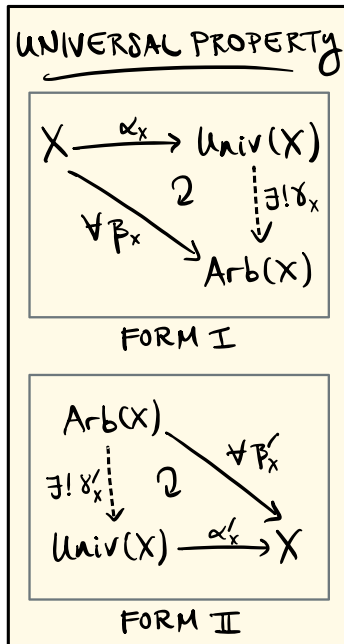


COEQUALIZER OF  $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

OPERATION ON PARALLEL MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



COEQUALIZER OF  $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

IS AN OBJECT  $\text{coeq}(f, g) \in \mathcal{C}$

EQUIPPED WITH A MORPHISM

$$\alpha: Y \rightarrow \text{coeq}(f, g)$$

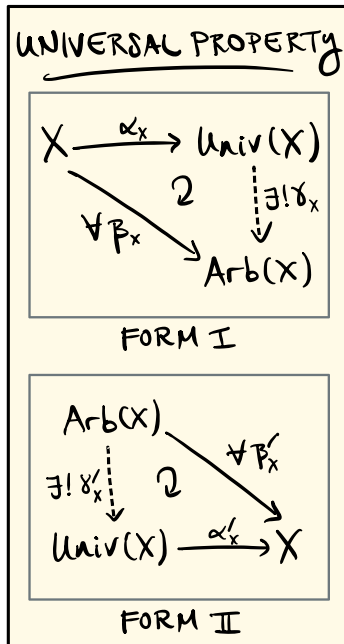
$$\Rightarrow \alpha f = \alpha g$$

$$X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y \xrightarrow{\alpha} \text{coeq}(f, g)$$

# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

OPERATION ON PARALLEL MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



COEQUALIZER OF  $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

IS AN OBJECT  $\text{coeq}(f, g) \in \mathcal{C}$

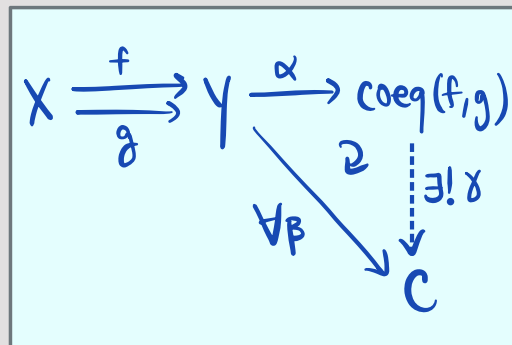
EQUIPPED WITH A MORPHISM

$$\alpha: Y \rightarrow \text{coeq}(f, g)$$

$$\Rightarrow \alpha f = \alpha g$$

WHERE  $\forall \beta: Y \rightarrow C \Rightarrow \beta f = \beta g$

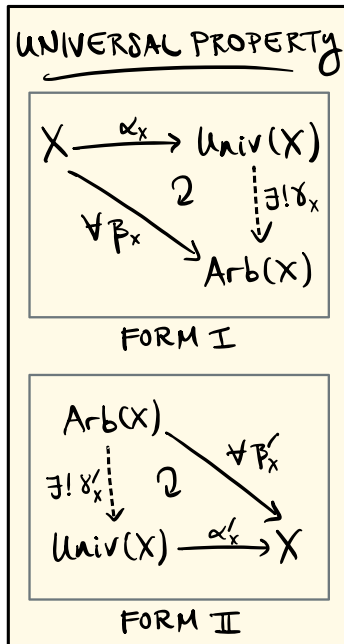
WE GET:



# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

OPERATION ON PARALLEL MORPHISMS



COEQUALIZER OF  $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

IS AN OBJECT  $coeq(f,g) \in \mathcal{C}$

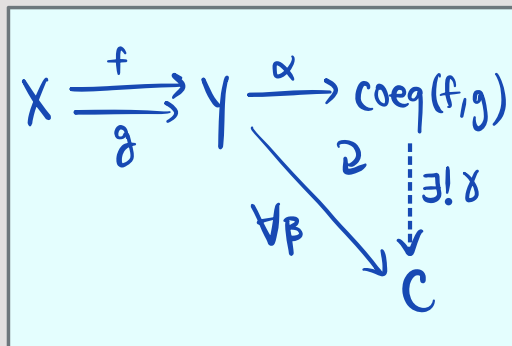
EQUIPPED WITH A MORPHISM

$$\alpha: Y \rightarrow coeq(f,g)$$

$$\Rightarrow \alpha f = \alpha g$$

WHERE  $\forall \beta: Y \rightarrow C \Rightarrow \beta f = \beta g$

WE GET:



EQUALIZER OF  $X \begin{matrix} \xrightarrow{f} \\ \xrightarrow{g} \end{matrix} Y$

IS AN OBJECT  $eq(f,g) \in \mathcal{C}$

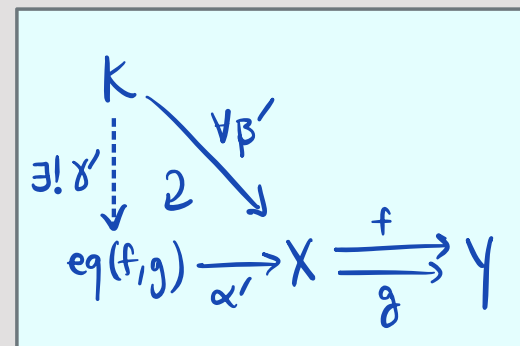
EQUIPPED WITH A MORPHISM

$$\alpha': eq(f,g) \rightarrow X$$

$$\Rightarrow f \alpha' = g \alpha'$$

WHERE  $\forall \beta': K \rightarrow X \Rightarrow f \beta' = g \beta'$

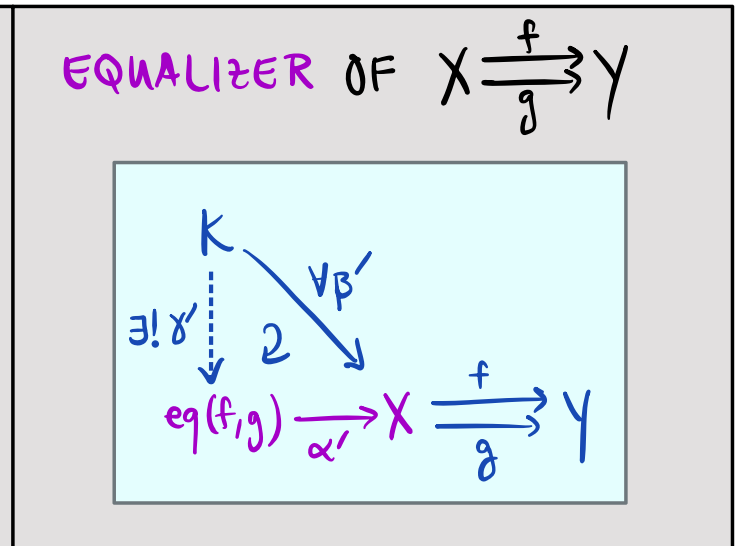
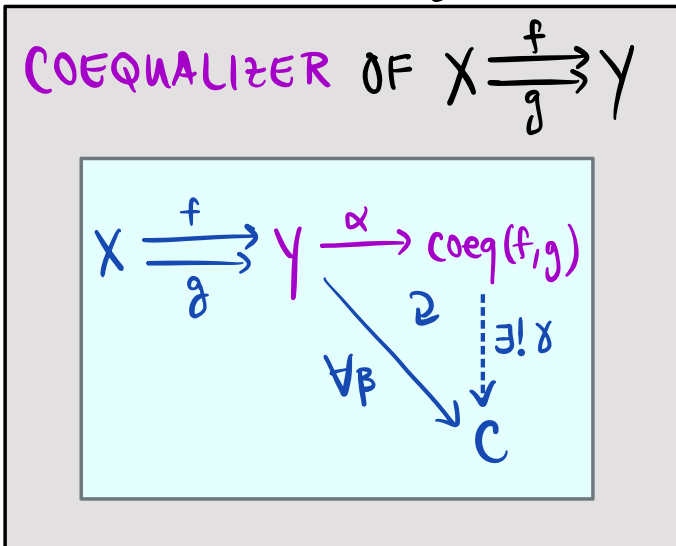
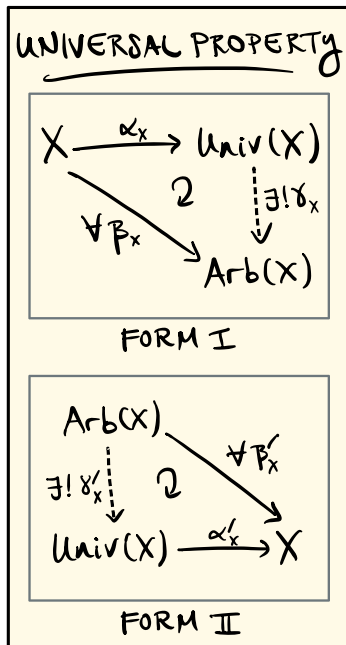
WE GET:



# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

OPERATION ON PARALLEL MORPHISMS



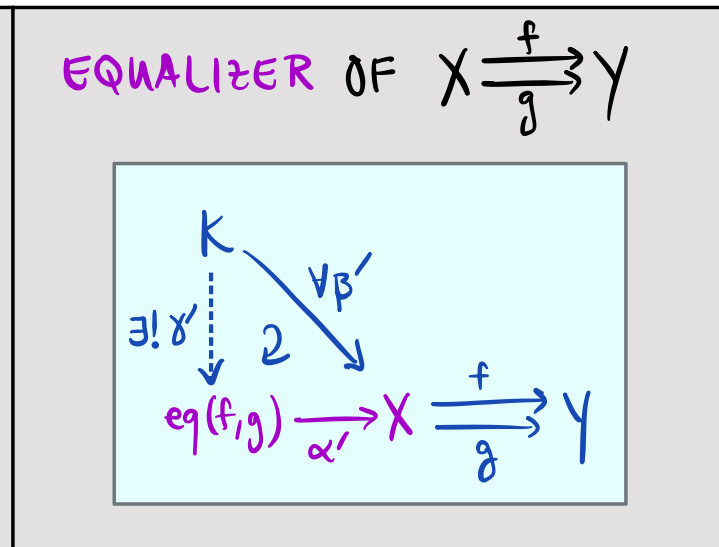
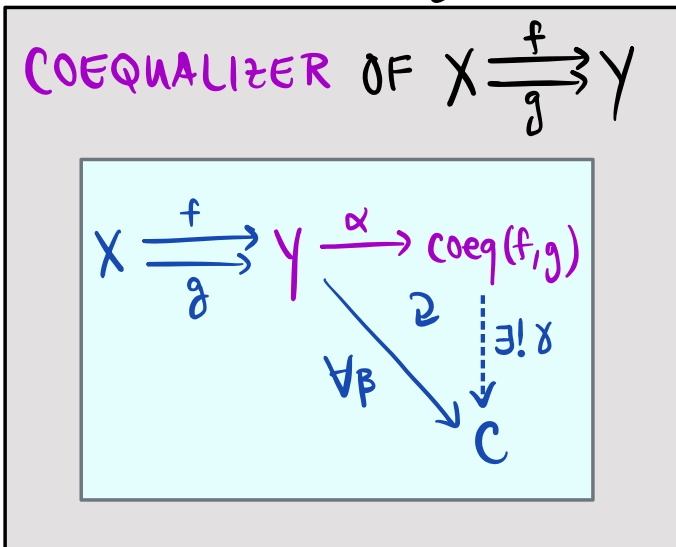
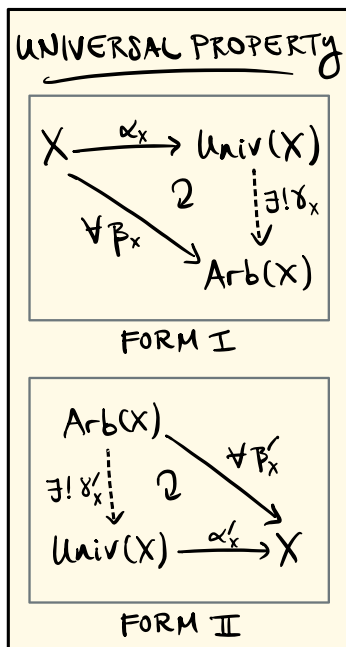
SOME  
SPECIAL  
CASES



# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

OPERATION ON PARALLEL MORPHISMS



SOME SPECIAL CASES

∞

EXERCISE 2.11

LET  $f, g : X \rightarrow Y$   
BE FUNCTIONS IN Set.

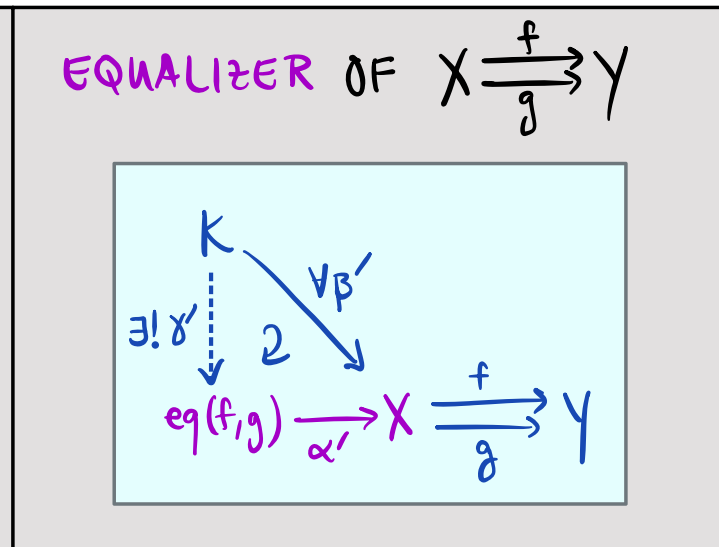
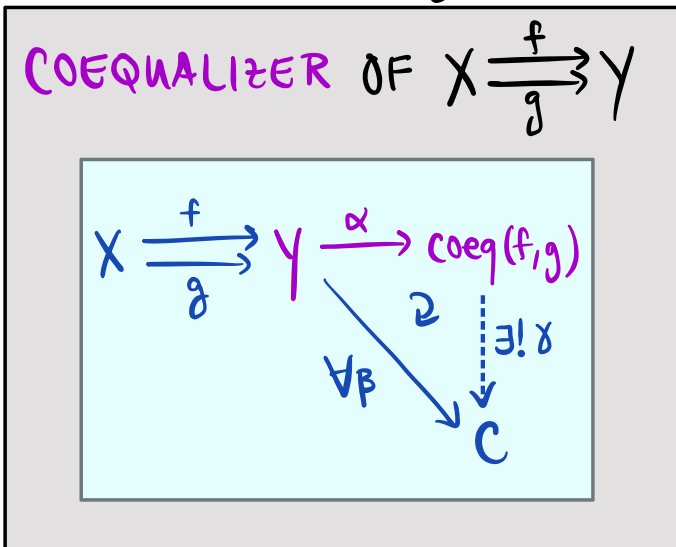
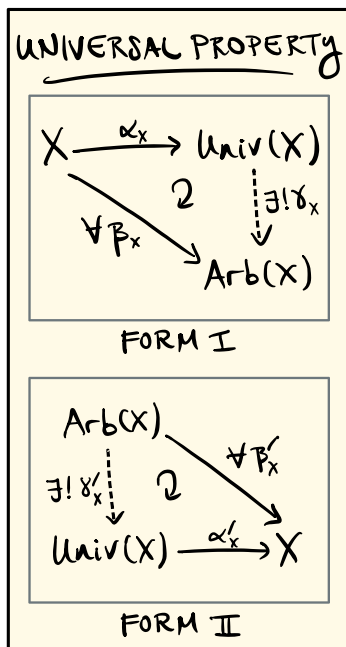
THEN:

$eq(f, g) = \{x \in X \mid f(x) = g(x)\}$   
SUBSET OF X

# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

OPERATION ON PARALLEL MORPHISMS



RECALL FOR  $\mathbb{R}$ -ALG.  $A$   
 W/ MODULES  $V_A \neq A W$   
 GET:

$$V \otimes_A W \cong (V \otimes_{\mathbb{R}} W) / R$$

$$\text{SPAN}_{\mathbb{R}} \langle (v+ a) \otimes w - v \otimes (a \cdot w) \rangle$$

$v \in V, w \in W, a \in A$

## EXERCISE 2.11

LET  $f, g : X \rightarrow Y$   
 BE FUNCTIONS IN Set.

THEN:

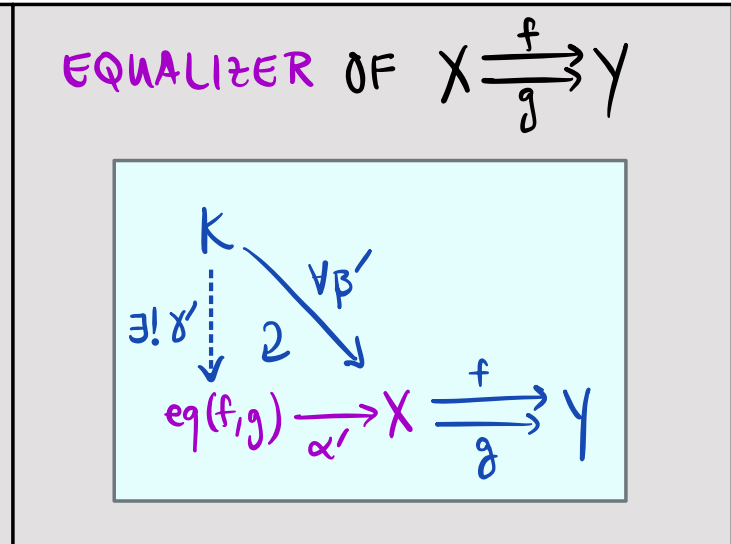
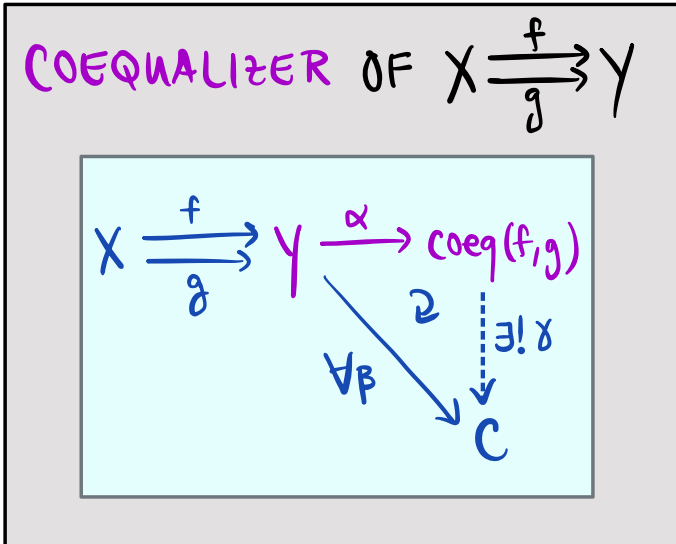
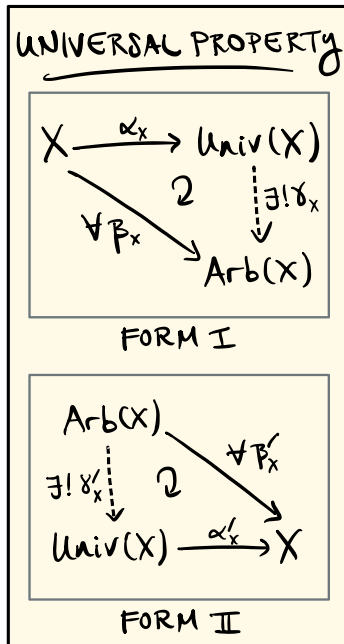
$$\text{eq}(f, g) = \{x \in X \mid f(x) = g(x)\}$$

SUBSET OF  $X$

# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

GIVEN A CATEGORY  $\mathcal{C}$ :

OPERATION ON PARALLEL MORPHISMS



RECALL FOR  $\mathbb{R}$ -ALG.  $A$   
 $W$ / MODULES  $V_A \neq A W$   
 GET:

$$V \otimes_A W \cong (V \otimes_{\mathbb{R}} W) / \mathcal{R}$$

SPAN  $\langle (v+aw) \otimes w - v \otimes (aw) \mid v \in V, w \in W, a \in A \rangle$

EXERCISE 2.11 WE GET

$$V \otimes_A W \cong \text{Coeq}(f, g)$$

FOR

$$f: V \otimes_{\mathbb{R}} A \otimes_{\mathbb{R}} W \xrightarrow{\triangleleft \otimes_{\mathbb{R}} \text{id}_W} V \otimes_{\mathbb{R}} W$$

$$g: V \otimes_{\mathbb{R}} A \otimes_{\mathbb{R}} W \xrightarrow{\text{id}_V \otimes_{\mathbb{R}} \triangleright} V \otimes_{\mathbb{R}} W$$

EXERCISE 2.11

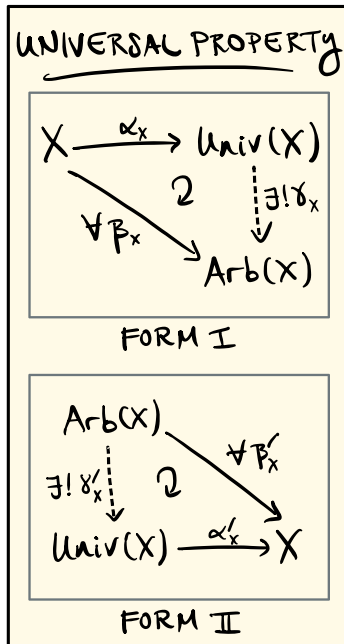
LET  $f, g: X \rightarrow Y$   
 BE FUNCTIONS IN SET.

THEN:

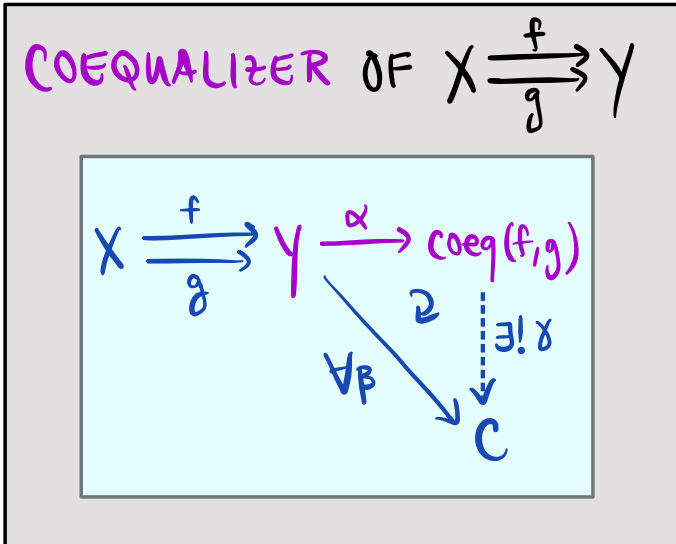
$$\text{eq}(f, g) = \{x \in X \mid f(x) = g(x)\}$$

SUBSET OF  $X$

# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS

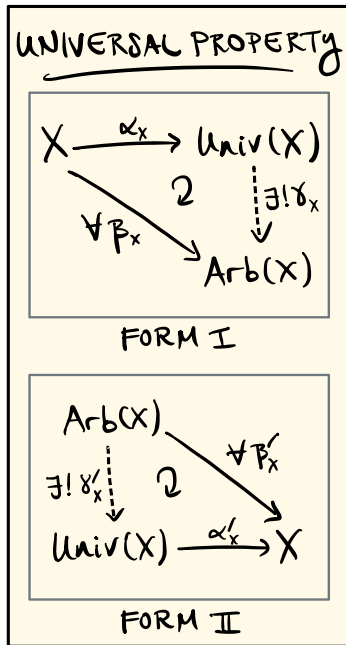


GIVEN A CATEGORY  $\mathcal{C}$ :

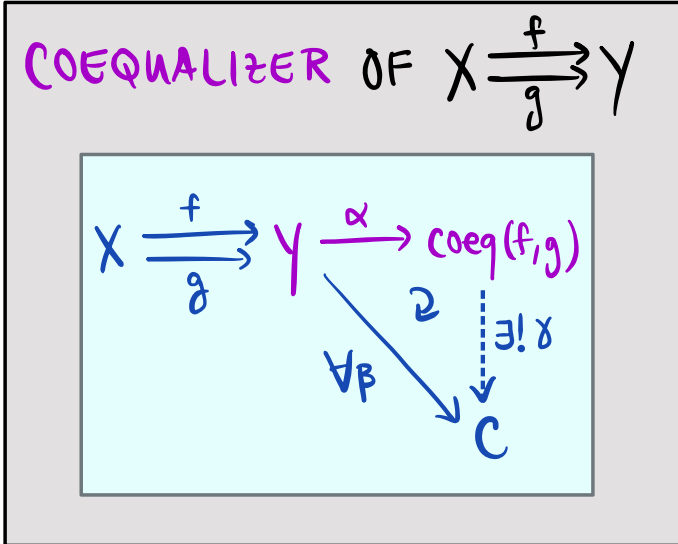


ARISES FROM OTHER  
UNIV. CONSTRUCTIONS

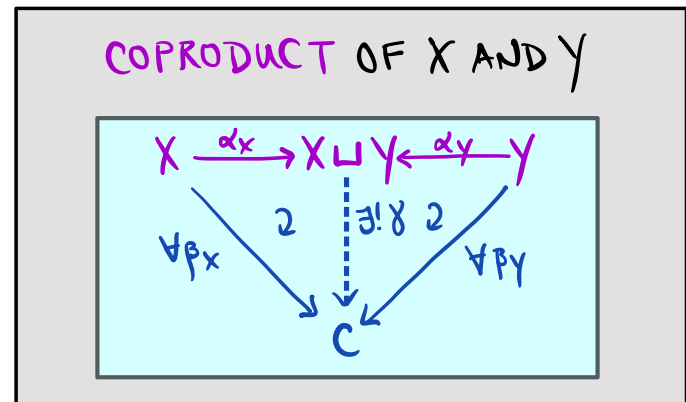
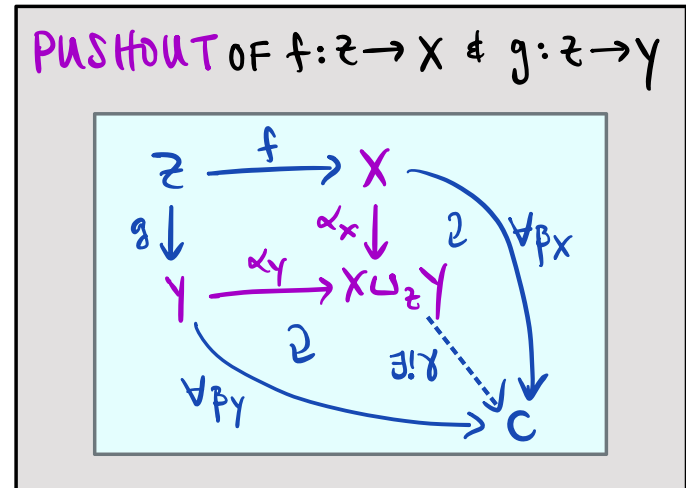
# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS



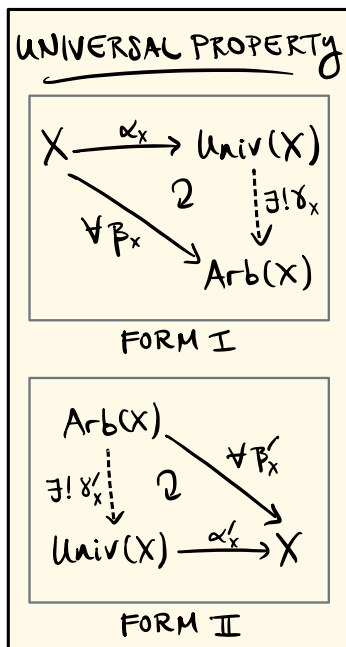
GIVEN A CATEGORY  $\mathcal{C}$ :



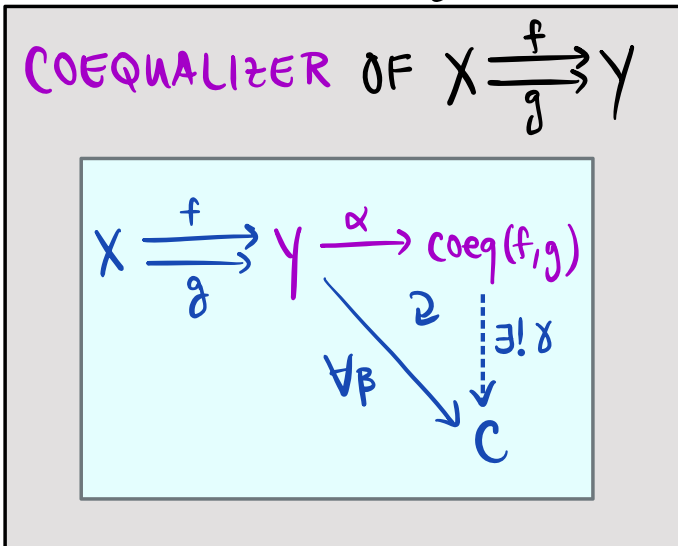
ARISES FROM OTHER UNIV. CONSTRUCTIONS



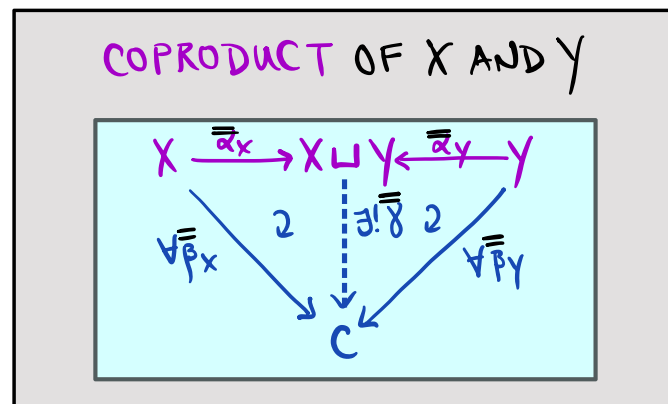
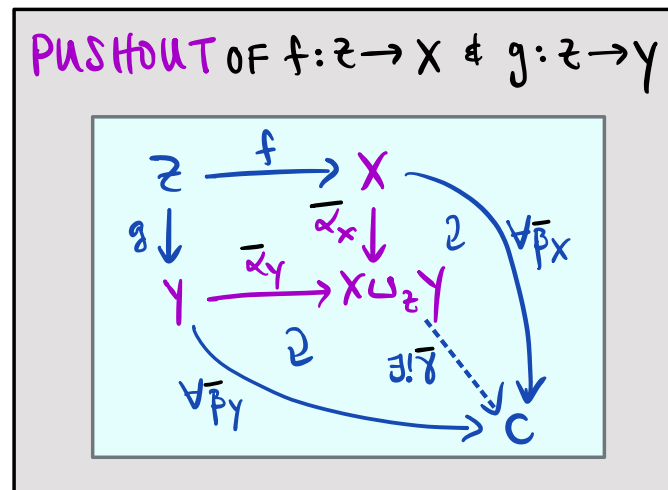
# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS



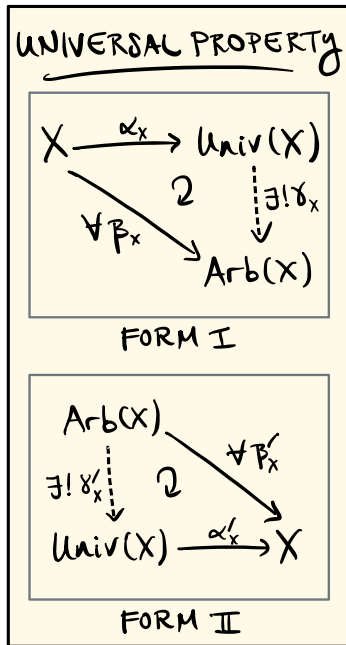
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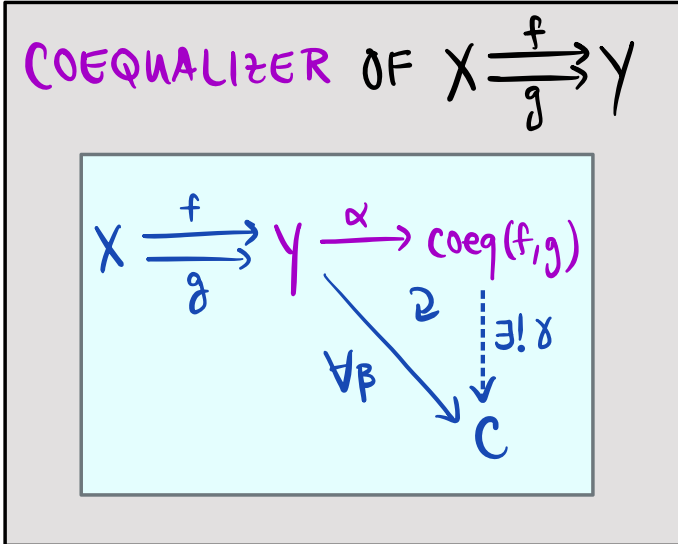
ARISES FROM OTHER UNIV. CONSTRUCTIONS



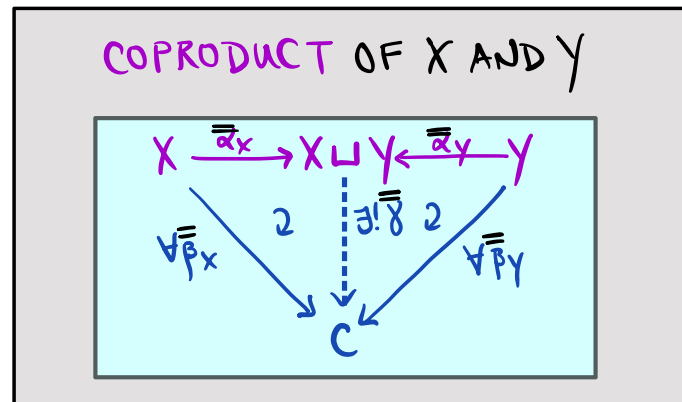
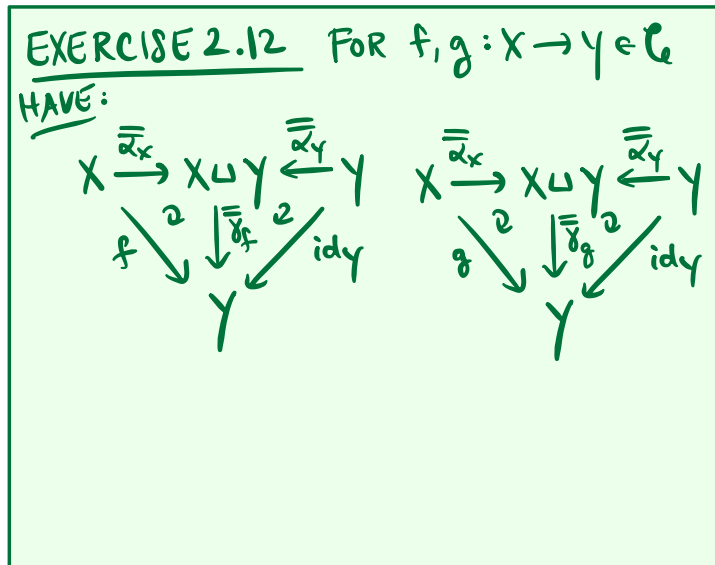
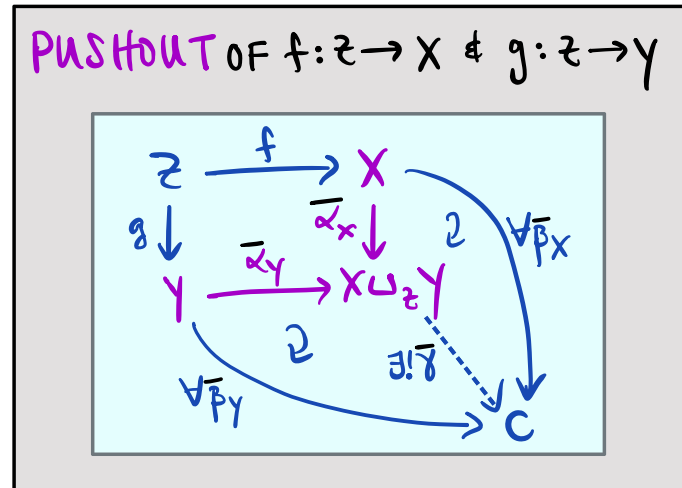
# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS



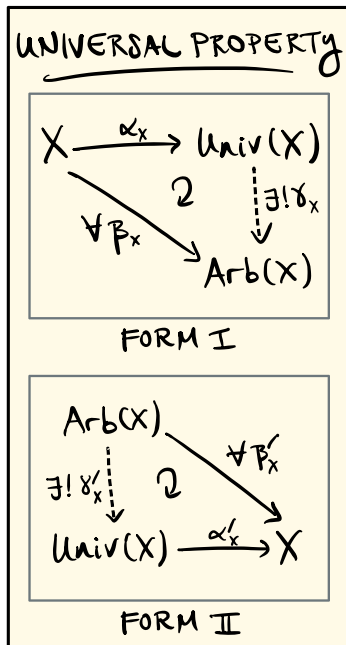
GIVEN A CATEGORY  $\mathcal{C}$ :



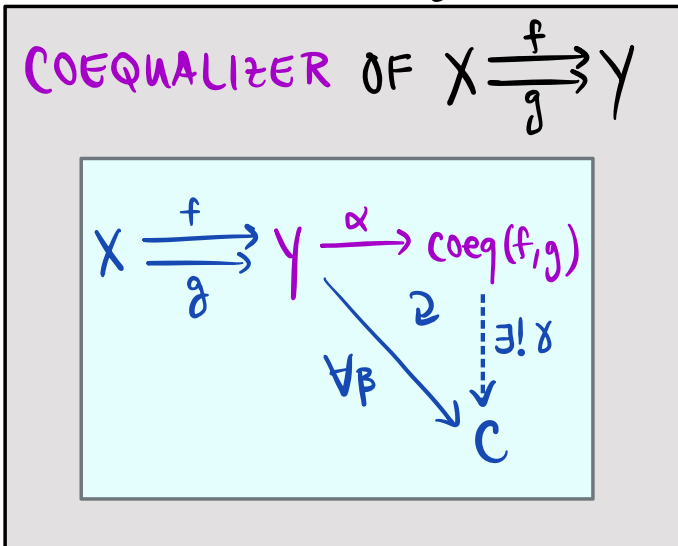
ARISES FROM OTHER UNIV. CONSTRUCTIONS



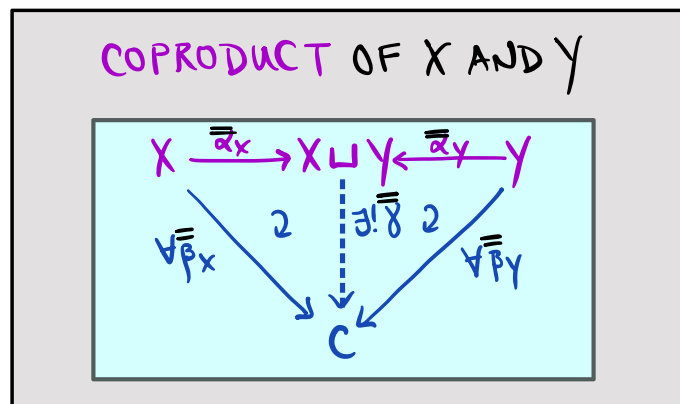
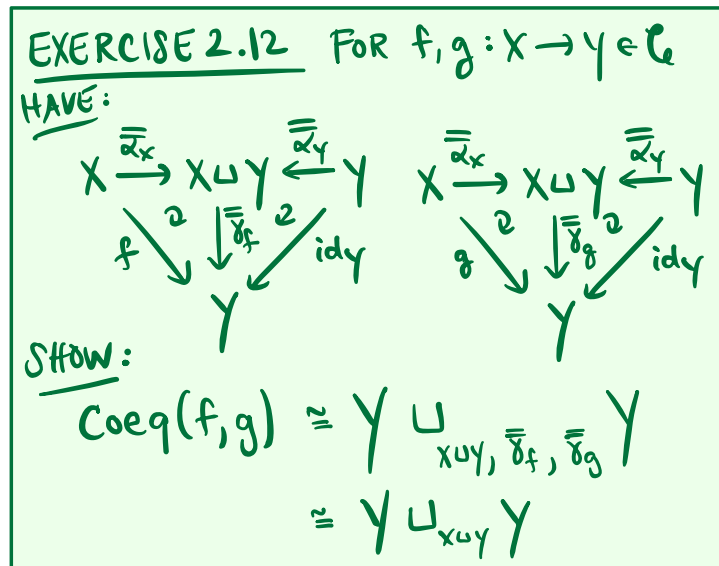
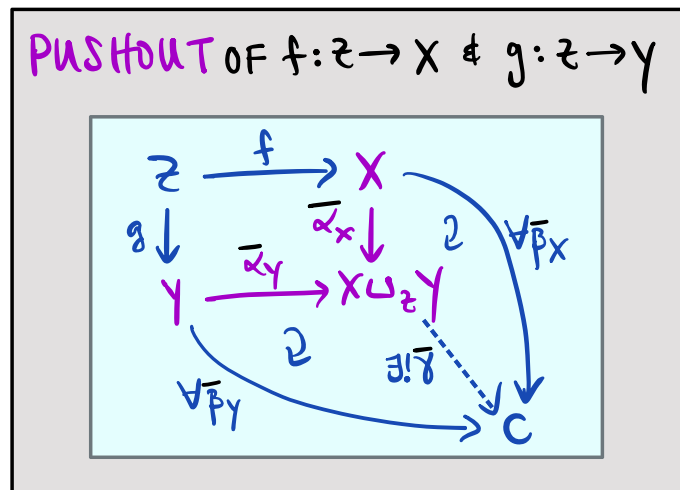
# I. UNIVERSAL CONSTRUCTIONS : COEQUALIZERS & EQUALIZERS



GIVEN A CATEGORY  $\mathcal{C}$ :

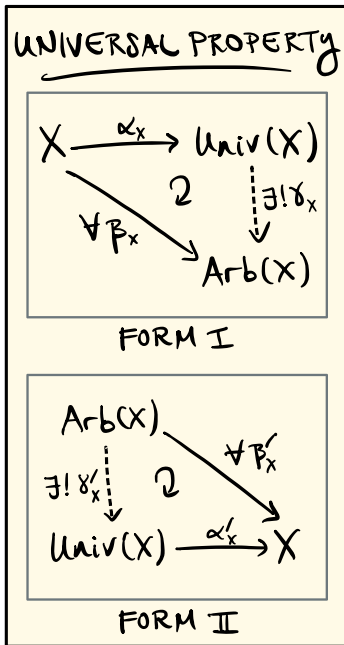


ARISES FROM OTHER UNIV. CONSTRUCTIONS

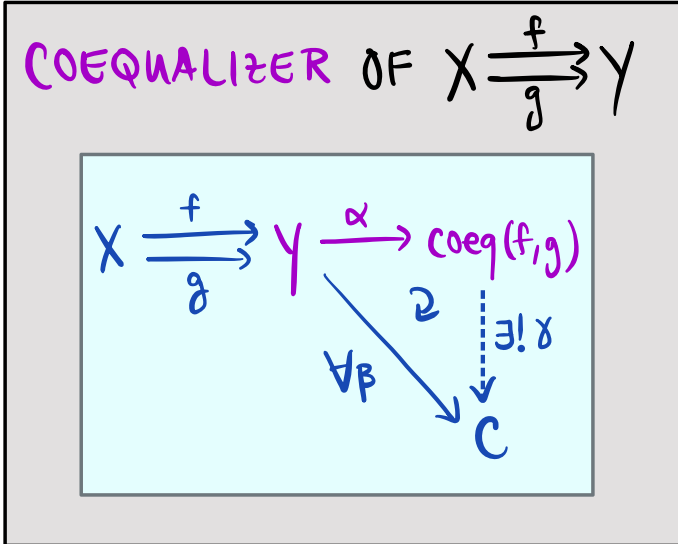




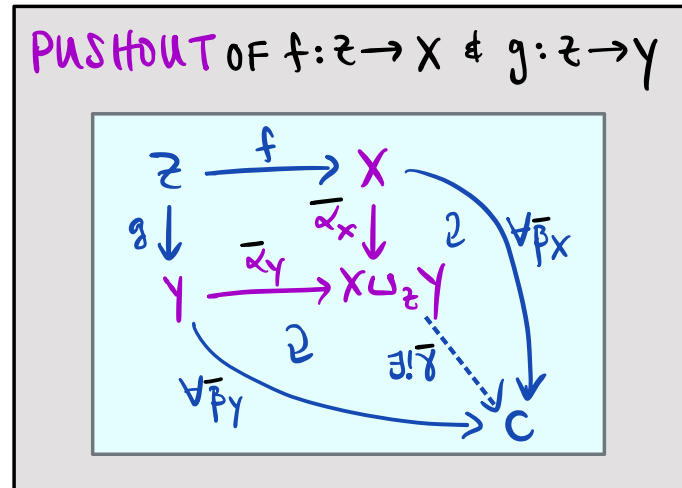
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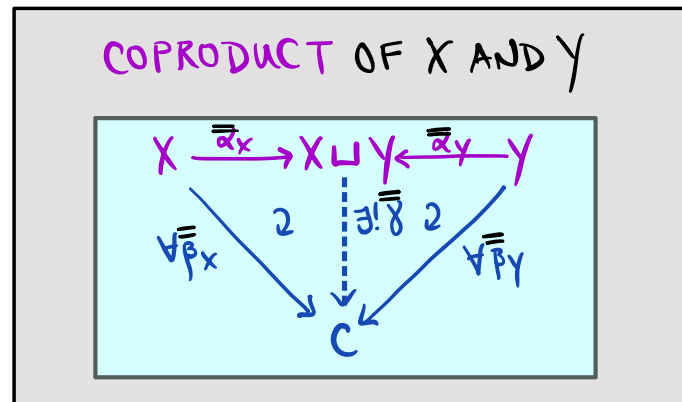
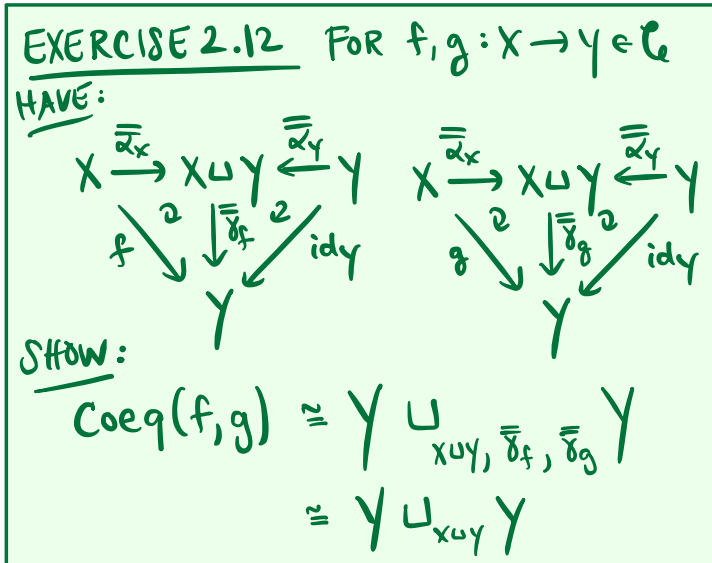
GIVEN A CATEGORY  $\mathcal{C}$ :



ARISES FROM OTHER UNIV. CONSTRUCTIONS

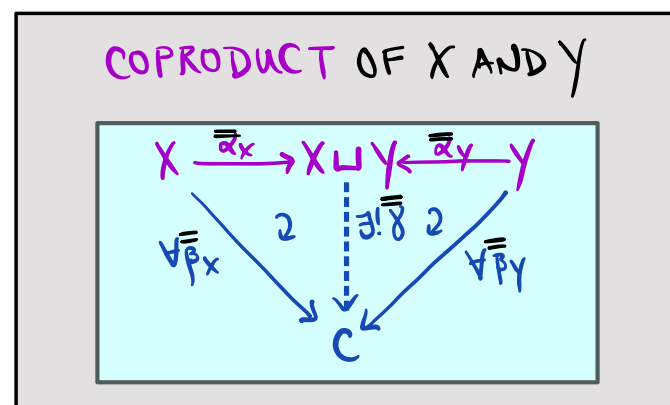
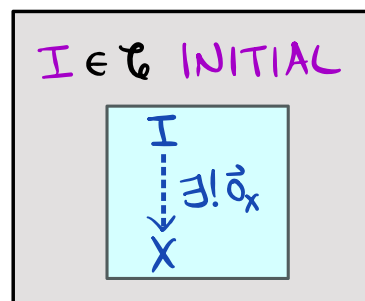
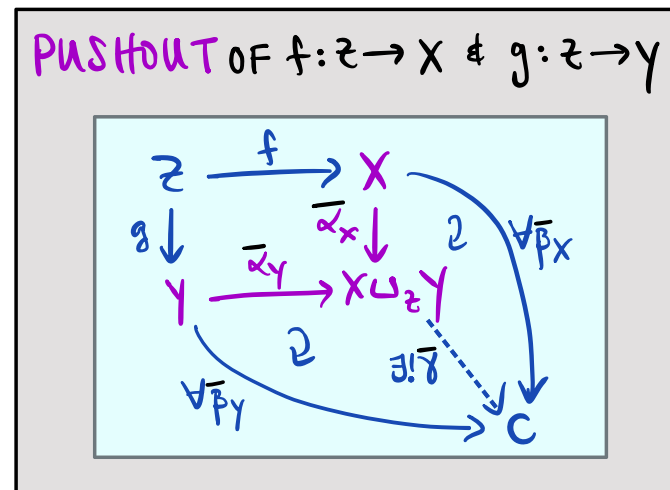
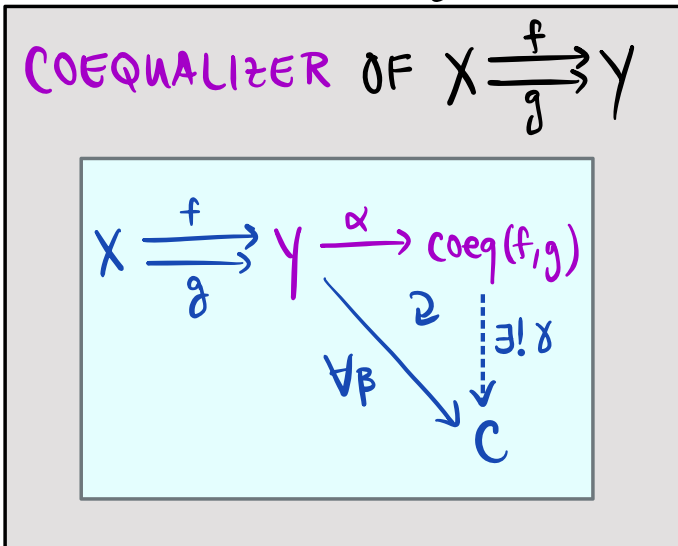
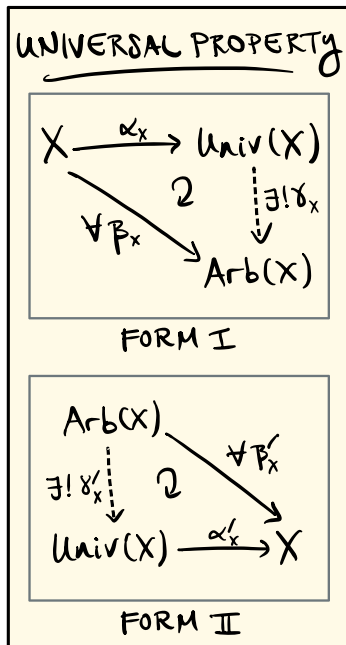


PUSHOUTS  
&  
COPRODUCTS  
EXIST  
 $\Downarrow$   
COEQUALIZERS  
EXIST



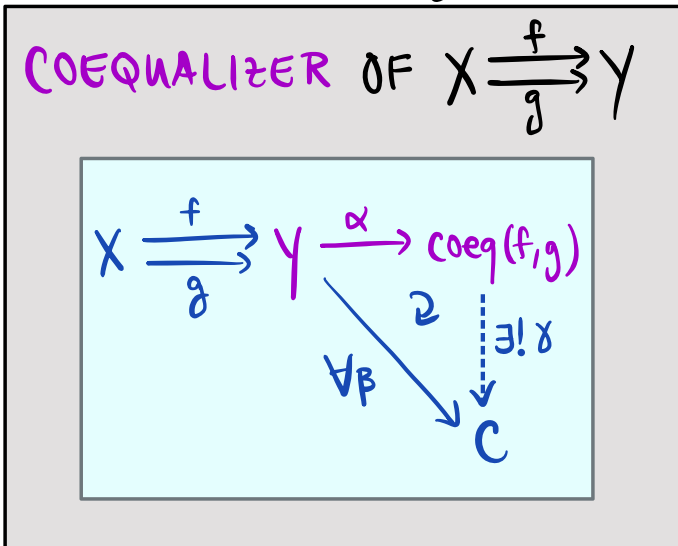
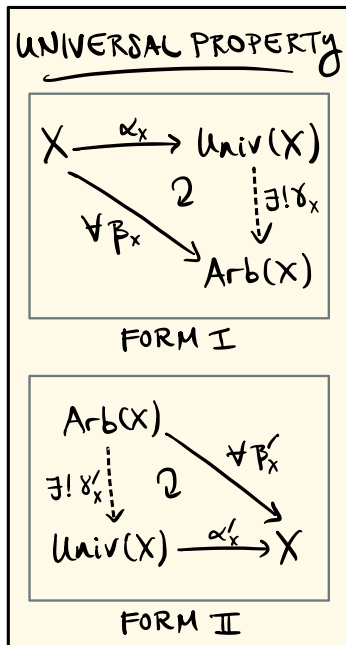
# I. UNIVERSAL CONSTRUCTIONS : RECAP

GIVEN A CATEGORY  $\mathcal{C}$ :

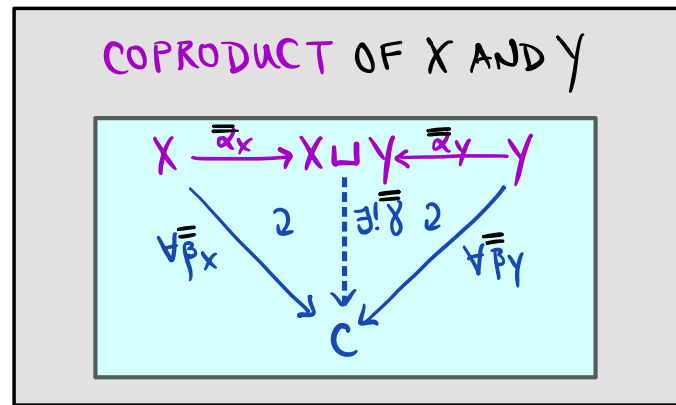
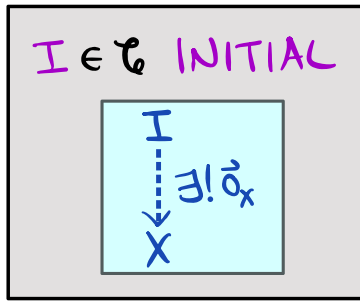
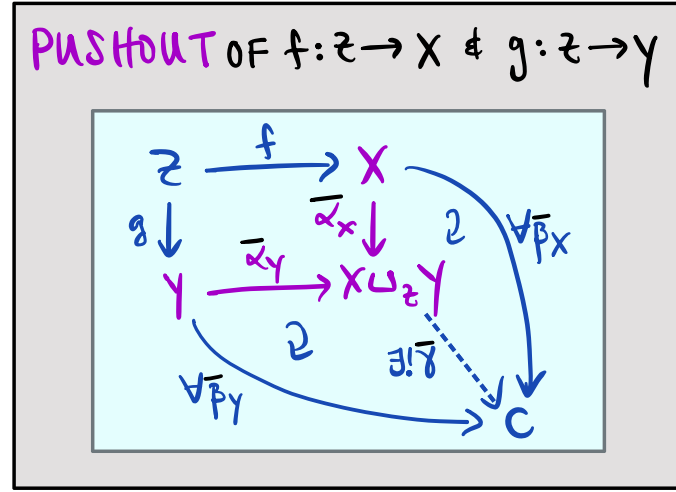


# I. UNIVERSAL CONSTRUCTIONS : RECAP

GIVEN A CATEGORY  $\mathcal{C}$ :

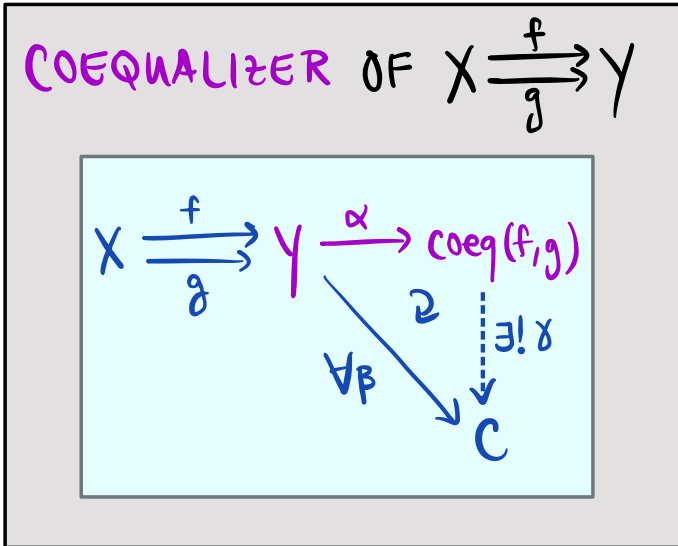
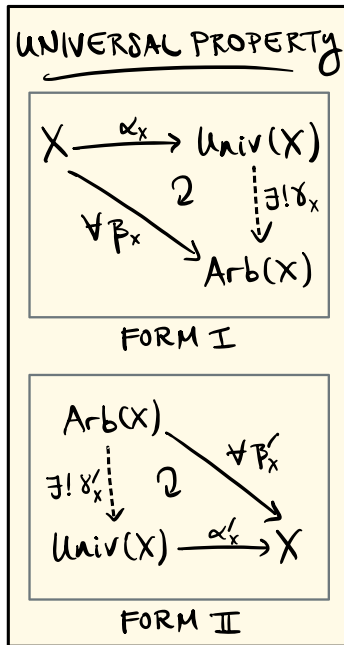


HAVE THE FOLLOWING EXISTENCE IMPLICATIONS

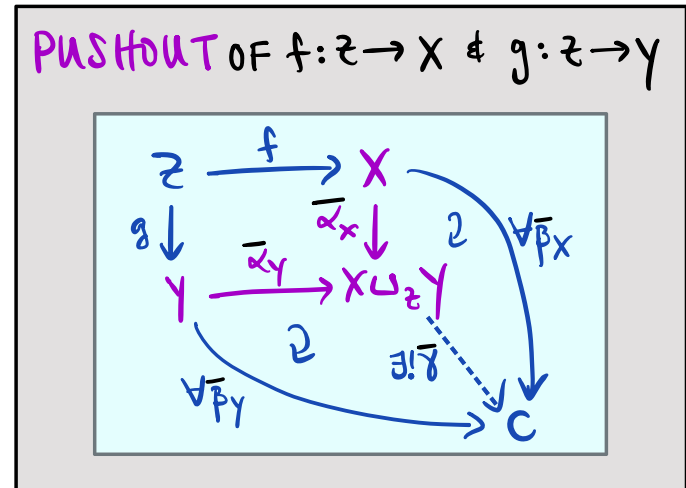


# I. UNIVERSAL CONSTRUCTIONS : RECAP

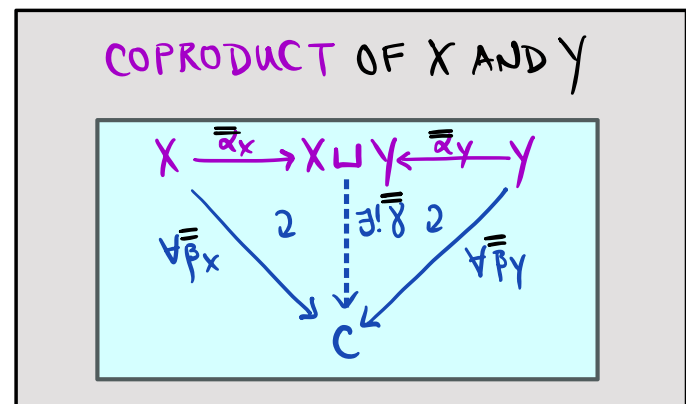
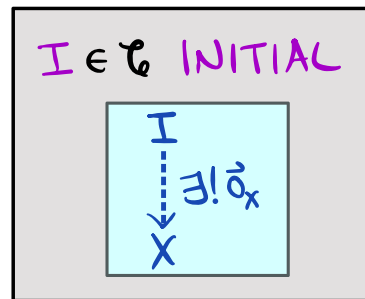
GIVEN A CATEGORY  $\mathcal{C}$ :



HAVE THE FOLLOWING EXISTENCE IMPLICATIONS



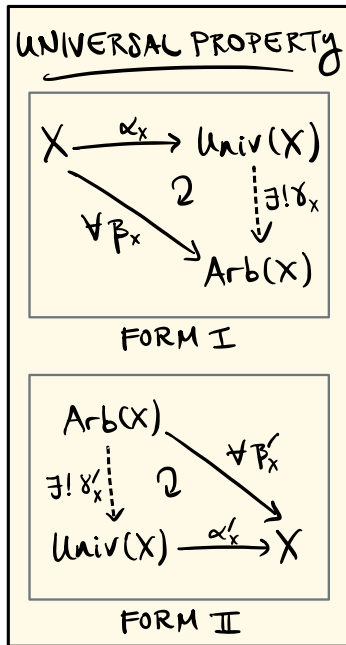
‡ LIKEWISE FOR  
 EQUALIZERS  
 PULLBACKS  
 TERMINAL OBJECT  
 PRODUCTS



# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

OPERATION ON MORPHISMS

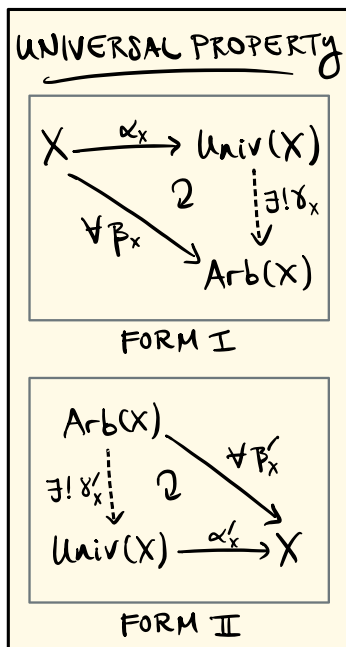
GIVEN A CATEGORY  $\mathcal{C}$ :



# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



A MORPHISM  $g: X \rightarrow Y$  IS

A ZERO MORPHISM IF

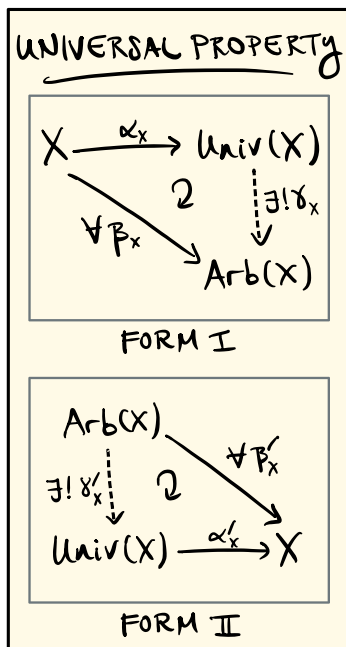
$$gf = gf' \quad \& \quad hg = h'g$$

$$\forall f, f': W \rightarrow X \in \mathcal{C} \quad \forall h, h': Y \rightarrow Z \in \mathcal{C}$$

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

OPERATION ON MORPHISMS

GIVEN A CATEGORY  $\mathcal{C}$ :



A MORPHISM  $g: X \rightarrow Y$  IS  
 A ZERO MORPHISM IF

$$gf = gf' \quad \& \quad hg = h'g$$

$$\forall f, f': W \rightarrow X \in \mathcal{C} \quad \forall h, h': Y \rightarrow Z \in \mathcal{C}$$

Ex.

IF  $\exists$  ZERO OBJECT IN  $\mathcal{C}$ , THEN:

$$\vec{0}_{X,Y} : X \xrightarrow{x\vec{0}} 0 \xrightarrow{\vec{0}_Y} Y$$

IS A ZERO MORPHISM.

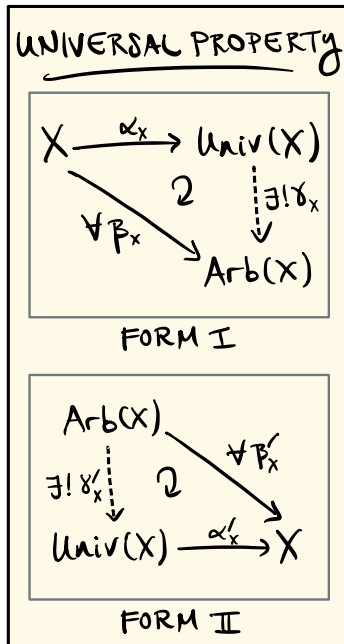


$\forall$  MORPHISMS  $f \in \mathcal{C}$ :

$$f \circ \vec{0} = \vec{0} \quad \& \quad \vec{0} \circ f = \vec{0}$$

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS



$g$  IS A ZERO MORPH.

IF  $gf = gf' \quad \forall f, f'$   
 $\neq hg = h'g \quad \forall h, h'$

Ex.

$$\vec{0}_{x,y} : X \xrightarrow{x\vec{0}} 0 \xrightarrow{\vec{0}_y} Y$$

$$\downarrow$$

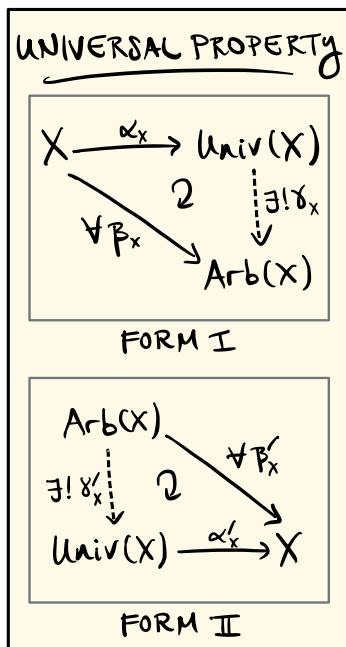
$$f \cdot \vec{0} = \vec{0} \neq \vec{0} \cdot f = \vec{0}$$

$\forall f \in \mathcal{C}$



# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS



COKERNEL OF  $f: X \rightarrow Y$

IS AN OBJECT  $\text{coker}(f) \in \mathcal{C}$

EQUIPPED WITH A MORPHISM

$$\alpha: Y \rightarrow \text{coker}(f)$$

$$\rightarrow \alpha f = \vec{0}_{X, \text{coker}(f)}$$

$$X \xrightarrow{f} Y \xrightarrow{\alpha} \text{coker}(f)$$

g IS A ZERO MORPH.

$$\text{IF } gf = gf' \quad \forall f, f'$$

$$\neq hg = h'g \quad \forall h, h'$$

Ex.

$$\vec{0}_{X, Y}: X \xrightarrow{\times \vec{0}} 0 \xrightarrow{\vec{0}_Y} Y$$

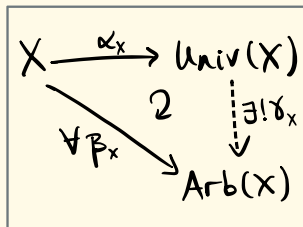
↓

$$f \circ \vec{0} = \vec{0} \neq \vec{0} \circ f = \vec{0} \quad \forall f \in \mathcal{C}$$

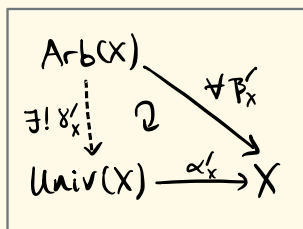
# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS

UNIVERSAL PROPERTY



FORM I



FORM II

$g$  IS A ZERO MORPH.

$$\text{IF } gf = gf' \quad \forall f, f'$$

$$\neq hg = h'g \quad \forall h, h'$$

Ex.

$$\vec{0}_{X,Y} : X \xrightarrow{\vec{0}_X} 0 \xrightarrow{\vec{0}_Y} Y$$



$$f \cdot \vec{0} = \vec{0} \neq \vec{0} \cdot f = \vec{0} \quad \forall f \in \mathcal{C}$$

COKERNEL OF  $f: X \rightarrow Y$

IS AN OBJECT  $\text{coker}(f) \in \mathcal{C}$

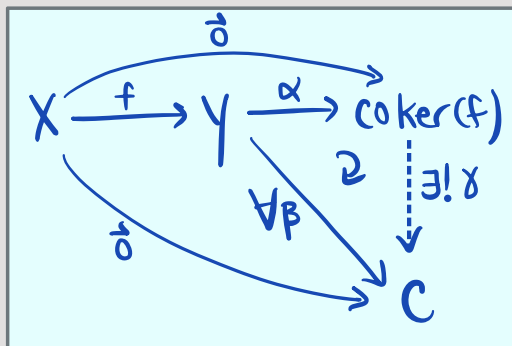
EQUIPPED WITH A MORPHISM

$$\alpha: Y \rightarrow \text{coker}(f)$$

$$\rightarrow \alpha f = \vec{0}_{X, \text{coker}(f)}$$

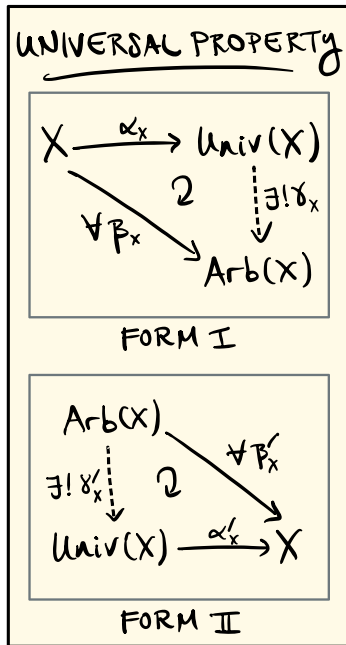
WHERE  $\forall \beta: Y \rightarrow C \quad \exists \beta f = \vec{0}_{X,C}$

WE GET:



# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.:  $\uparrow$  OPERATION ON MORPHISMS



COKERNEL OF  $f: X \rightarrow Y$

IS AN OBJECT  $\text{coker}(f) \in \mathcal{C}$

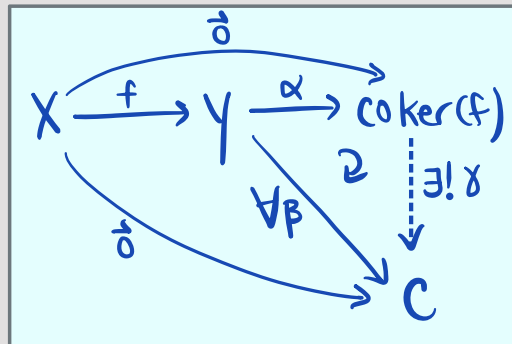
EQUIPPED WITH A MORPHISM

$$\alpha: Y \rightarrow \text{coker}(f)$$

$$\rightarrow \alpha f = \vec{0}_{X, \text{coker}(f)}$$

WHERE  $\forall \beta: Y \rightarrow C \exists \beta f = \vec{0}_{X, C}$

WE GET:



KERNEL OF  $f: X \rightarrow Y$

IS AN OBJECT  $\text{ker}(f) \in \mathcal{C}$

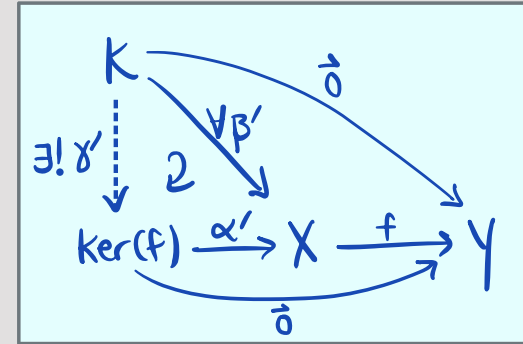
EQUIPPED WITH A MORPHISM

$$\alpha': \text{ker}(f) \rightarrow X$$

$$\rightarrow f \alpha' = \vec{0}_{\text{ker}(f), Y}$$

WHERE  $\forall \beta': K \rightarrow X \exists f \beta' = \vec{0}_{K, Y}$

WE GET:



$g$  IS A ZERO MORPH.

$$\text{IF } gf = gf' \quad \forall f, f'$$

$$\neq hg = h'g \quad \forall h, h'$$

Ex.

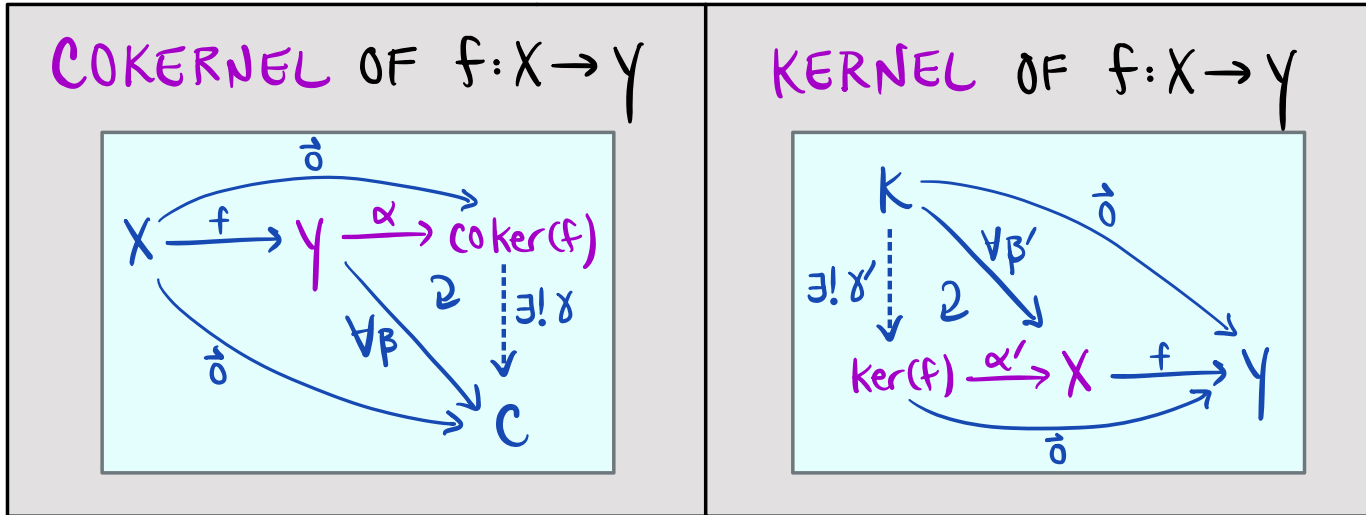
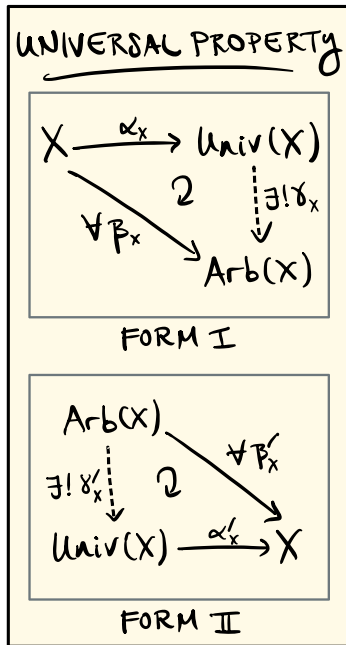
$$\vec{0}_{X, Y}: X \xrightarrow{\alpha} 0 \xrightarrow{\beta} Y$$

$\downarrow$

$$f \cdot \vec{0} = \vec{0} \neq \vec{0} \cdot f = \vec{0} \quad \forall f \in \mathcal{C}$$

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS



$g$  IS A ZERO MORPH.

IF  $gf = gf' \quad \forall f, f'$   
 $\neq hg = h'g \quad \forall h, h'$

Ex.

$$\vec{0}_{x,y}: X \xrightarrow{\vec{0}_x} 0 \xrightarrow{\vec{0}_y} Y$$

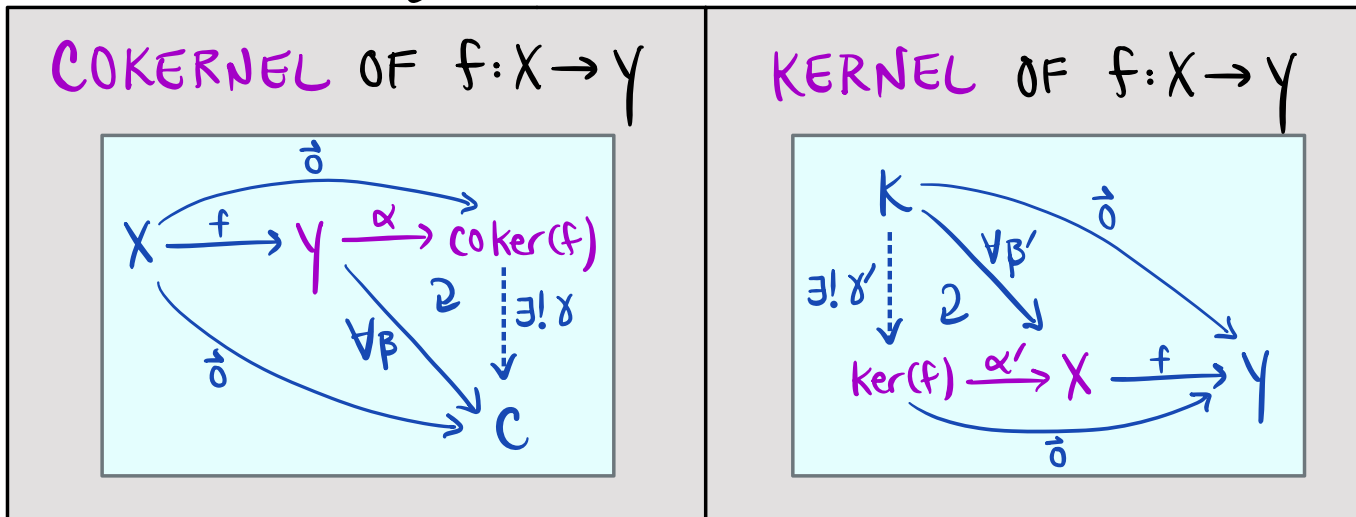
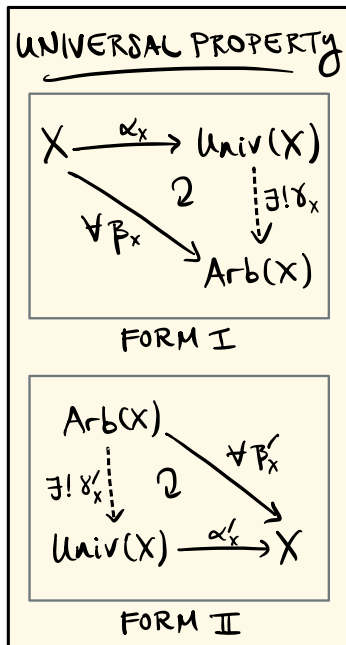
$$\downarrow$$

$$f \cdot \vec{0} = \vec{0} \neq \vec{0} \cdot f = \vec{0}$$

$\forall f \in \mathcal{C}$

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS



↑ RECOVERS THE USUAL NOTION OF (CO)KERNELS  
IN MANY CATEGORIES EXERCISE 2.15

$g$  IS A ZERO MORPH.

IF  $gf = gf' \forall f, f'$   
 $\neq hg = h'g \forall h, h'$

Ex.

$$\vec{0}_{X,Y}: X \xrightarrow{\vec{0}_X} 0 \xrightarrow{\vec{0}_Y} Y$$

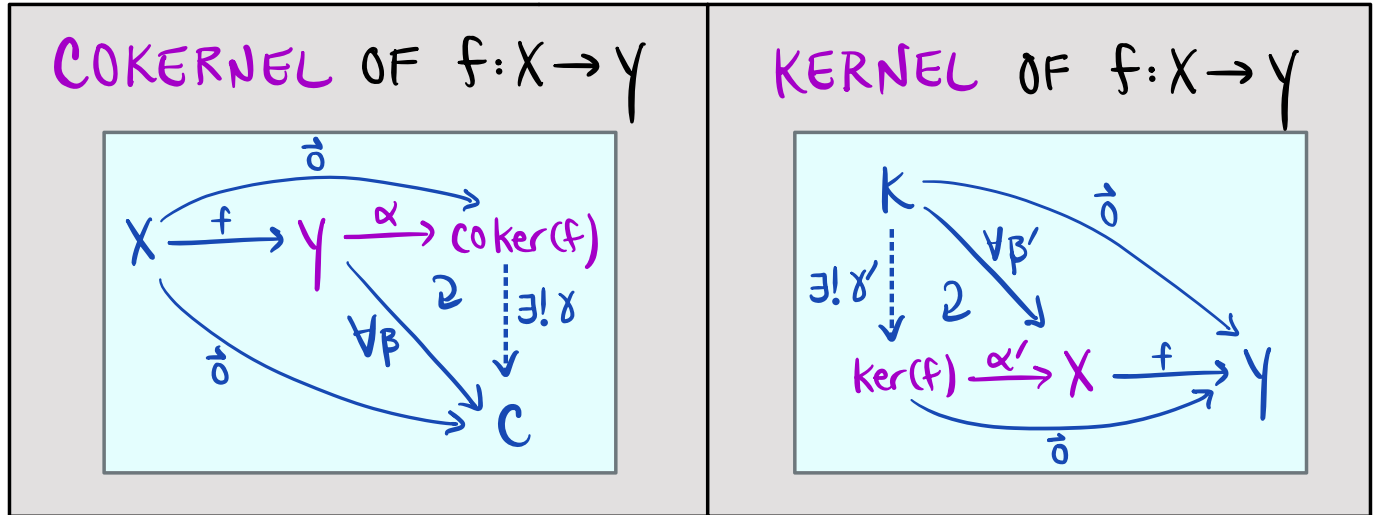
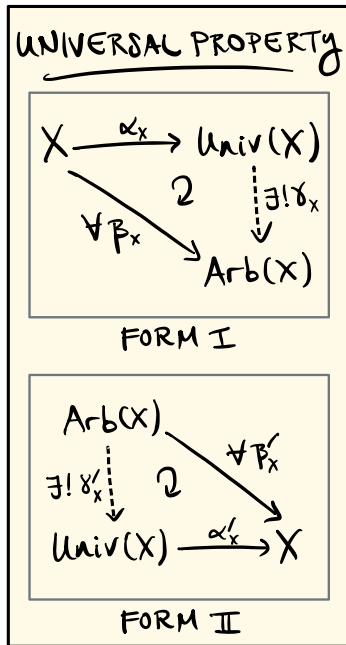
$$\downarrow$$

$$f \cdot \vec{0} = \vec{0} \neq \vec{0} \cdot f = \vec{0}$$

$\forall f \in \mathcal{C}$

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS



↖ RECOVERS THE USUAL NOTION OF (CO)KERNELS  
IN MANY CATEGORIES EXERCISE 2.15

↖ IS A SPECIAL CASE OF A (CO)EQUALIZER EXERCISE 2.13

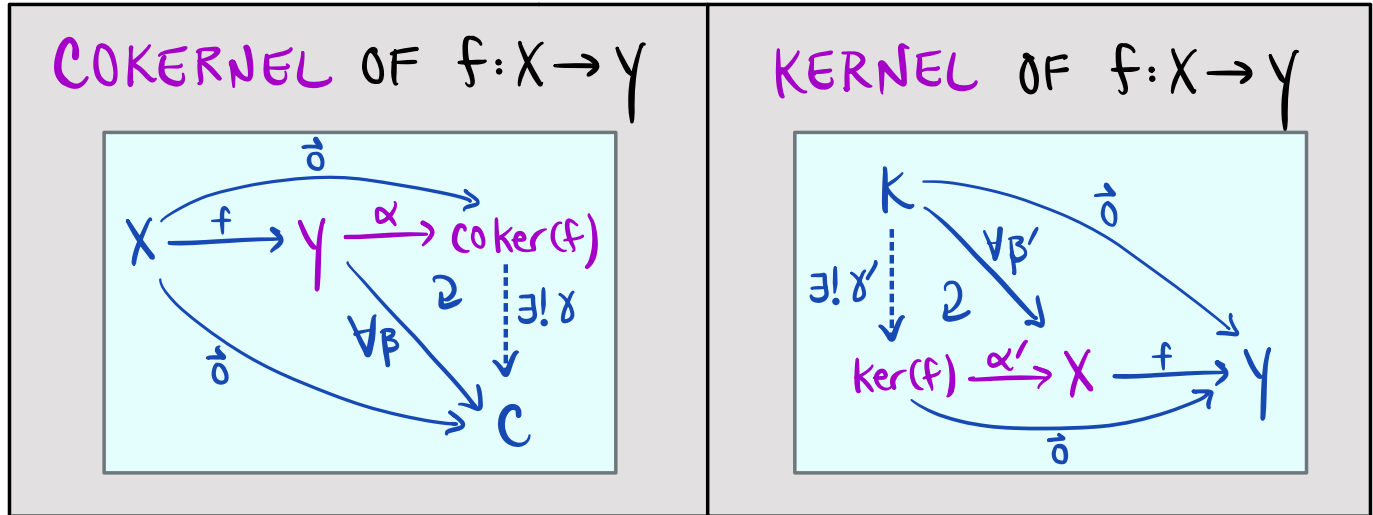
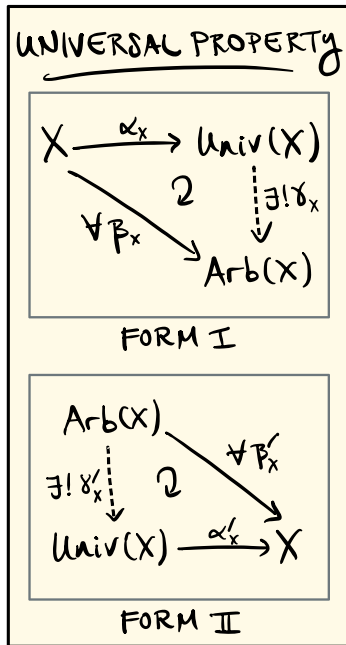
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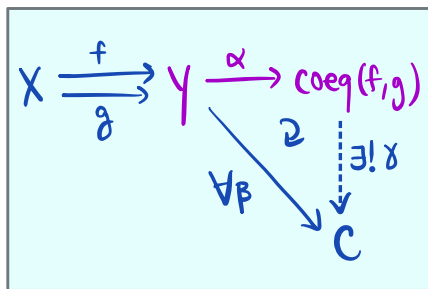
Ex.  
 $\vec{0}_{X,Y}: X \xrightarrow{\vec{0}_X} 0 \xrightarrow{\vec{0}_Y} Y$   
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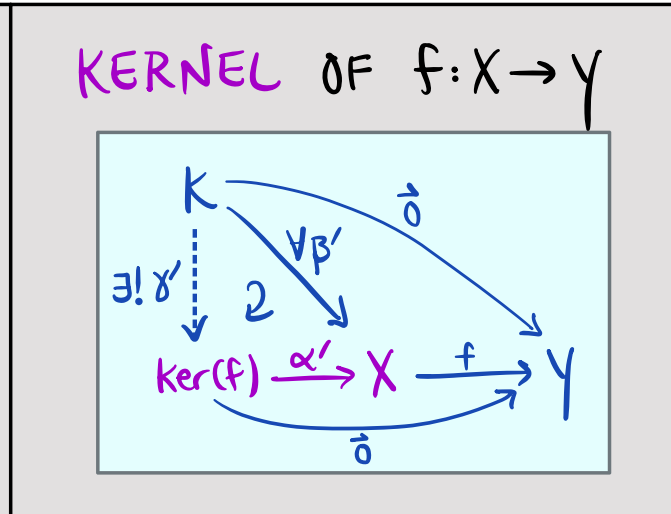
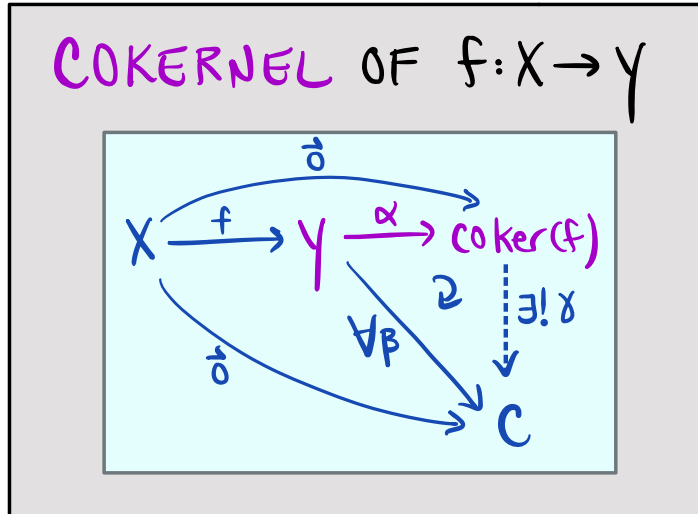
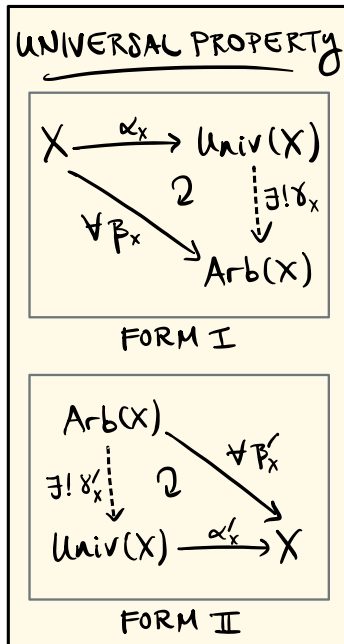
↖ IS A SPECIAL CASE OF A (CO)EQUALIZER

↖ THINK HOW TO GET A COKERNEL FROM THIS

EXERCISE 2.13

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

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Ex.

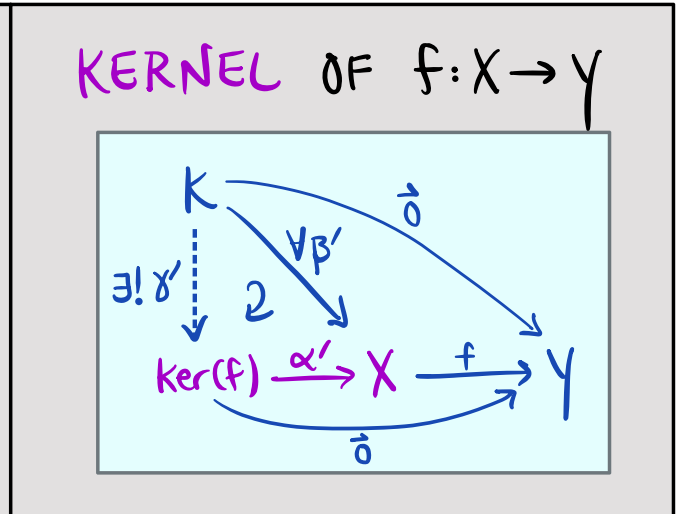
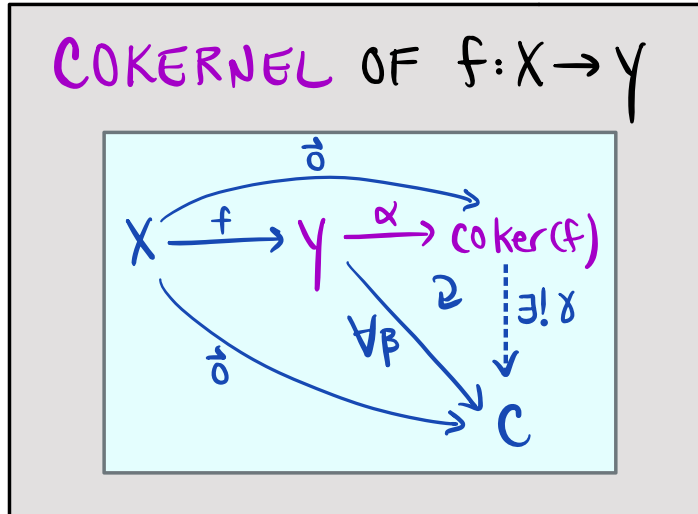
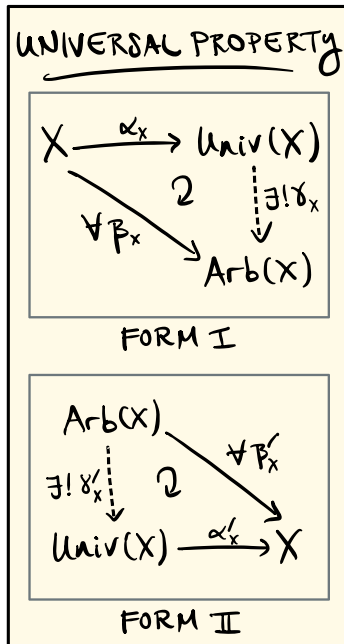
$$\begin{array}{ccc}
 \vec{0}_{x,y}: X & \xrightarrow{\vec{0}_x} & 0 & \xrightarrow{\vec{0}_y} & Y \\
 \downarrow & & & & \\
 f \cdot \vec{0} = \vec{0} & \neq & \vec{0} \cdot f = \vec{0} & & \\
 & & \forall f \in \mathcal{C} & & 
 \end{array}$$

EXERCISE 2.13 COKERNELS ARE EPIC & KERNELS ARE MONIC } RIGHT CANCELLATIVE  
{ LEFT CANCELLATIVE



# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

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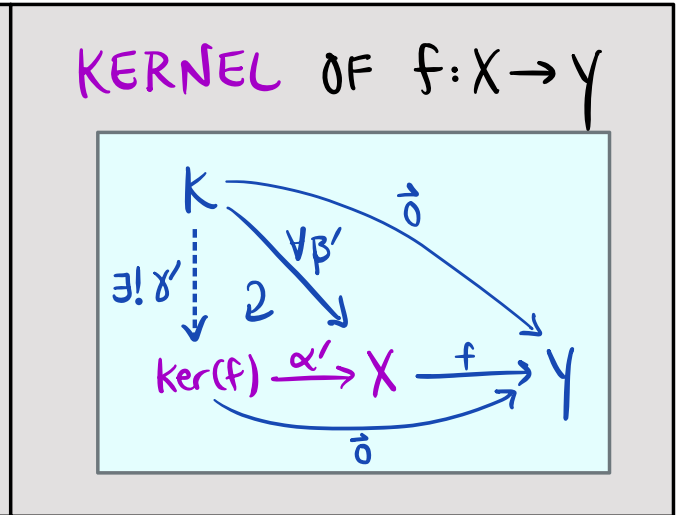
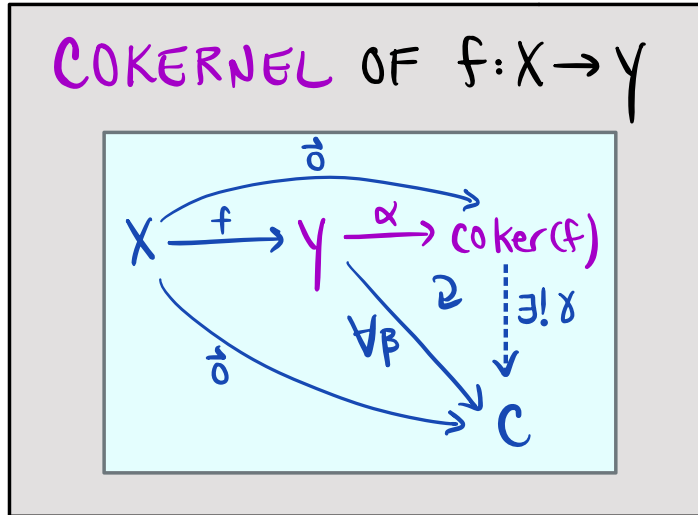
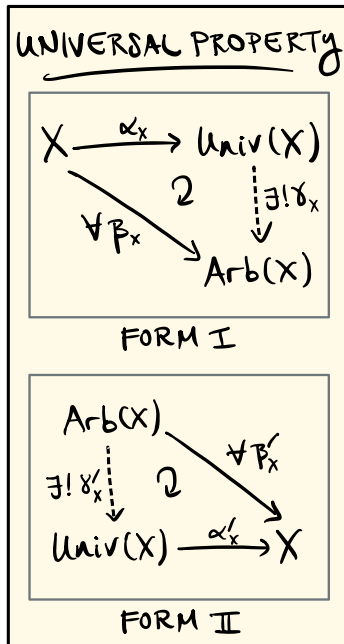
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**EXERCISE 2.13 COKERNELS ARE EPIC** ← RIGHT CANCELLATIVE

you do! ↗

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

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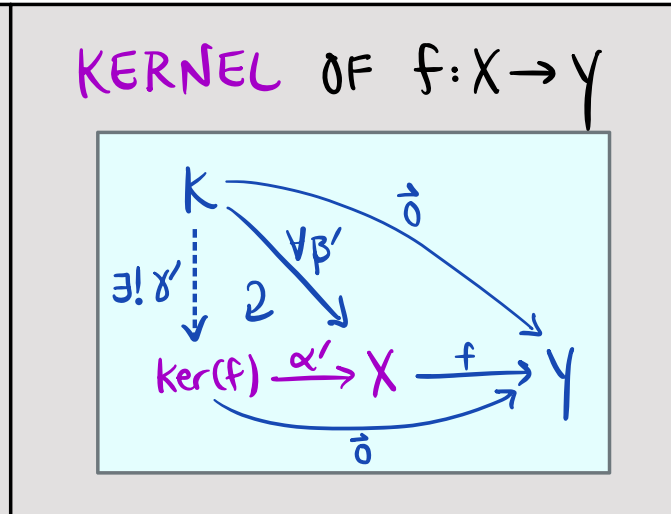
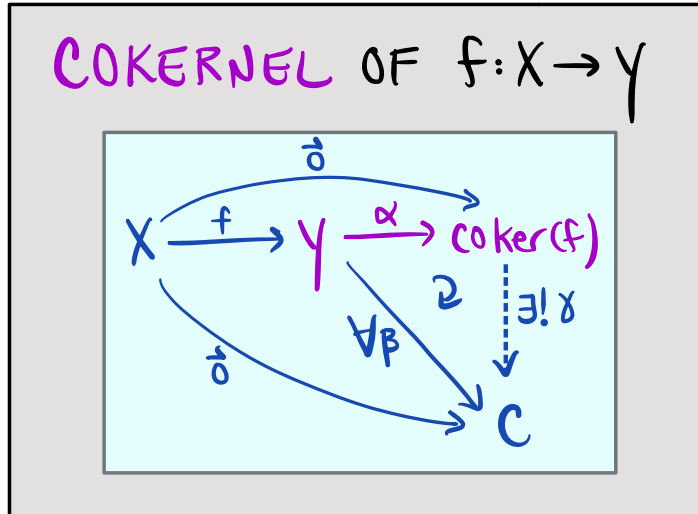
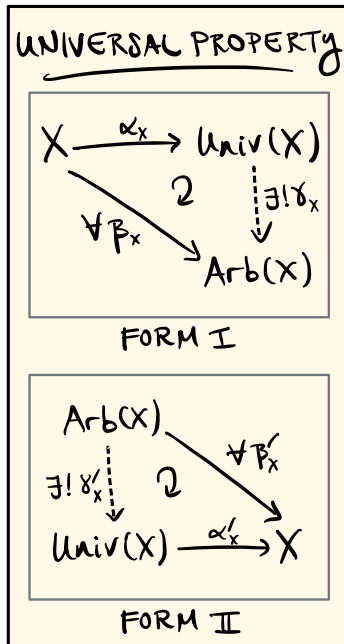
$\forall f \in \mathcal{C}$

EXERCISE 2.13 COKERNELS ARE EPIC,  $\leftarrow$  RIGHT CANCELLATIVE

Pf/ TAKE  $h, h' : \text{coker}(f) \rightarrow Z \Rightarrow h\alpha = h'\alpha$  AS MORPHISMS  $Y \rightarrow Z$

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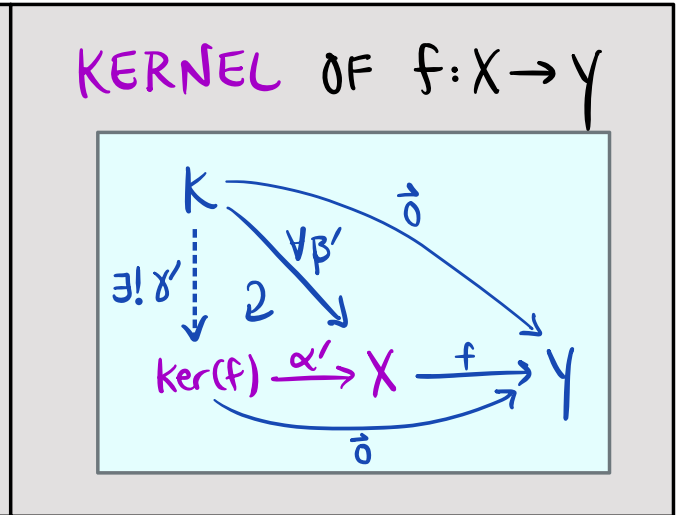
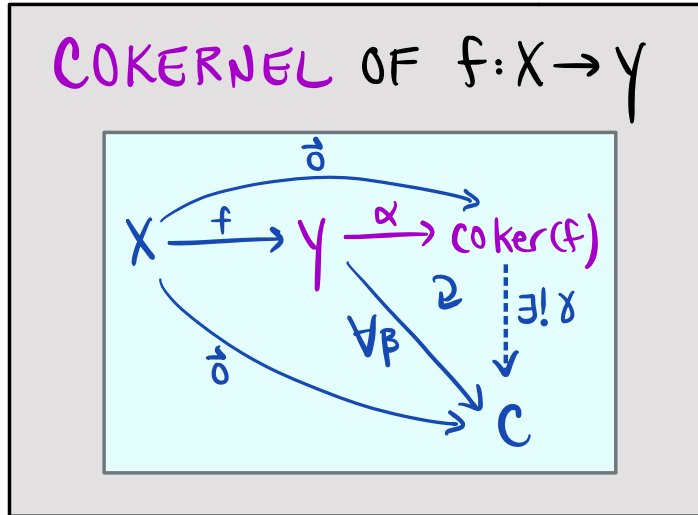
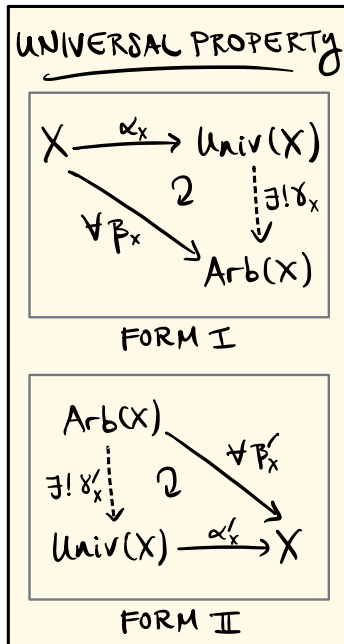
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NOW

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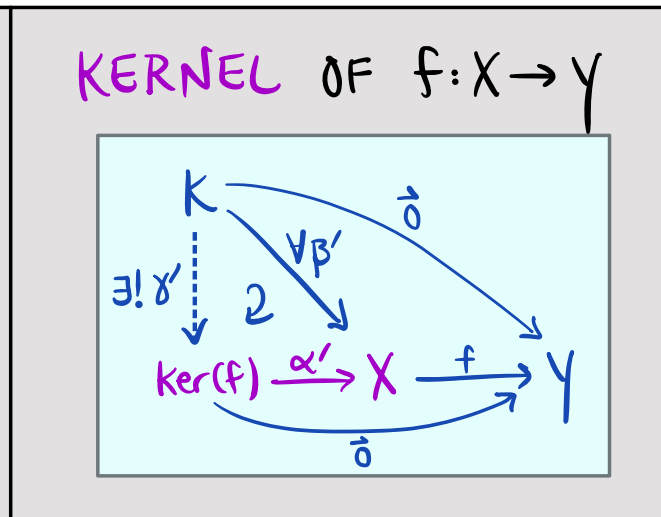
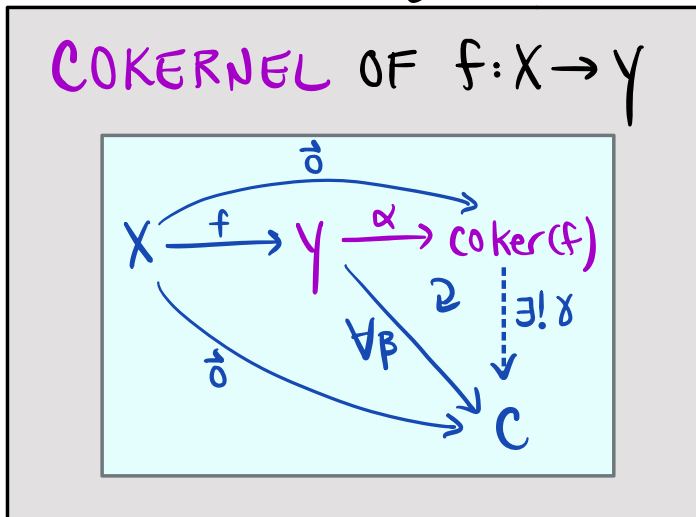
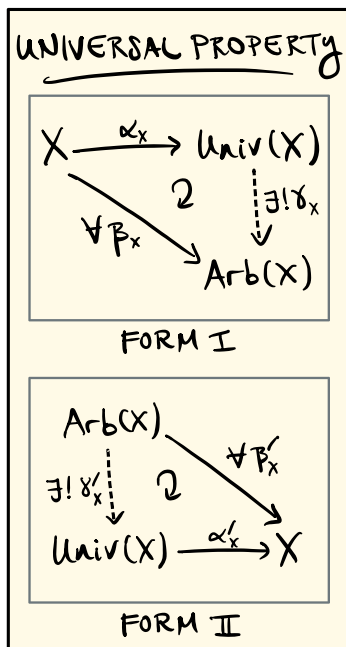
Pf/ TAKE  $h, h' : \text{coker}(f) \rightarrow Z \Rightarrow h\alpha = h'\alpha$  AS MORPHISMS  $Y \rightarrow Z$

NOW

$[(h\alpha)f = h(\alpha f) = h\vec{0} = \vec{0}]$

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Ex.  
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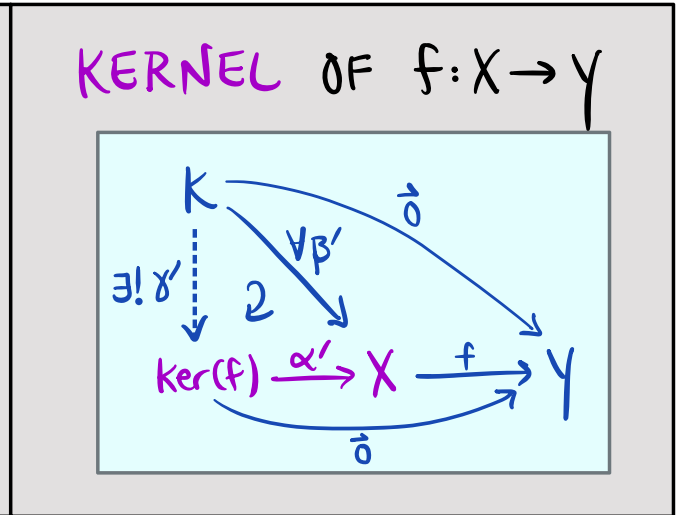
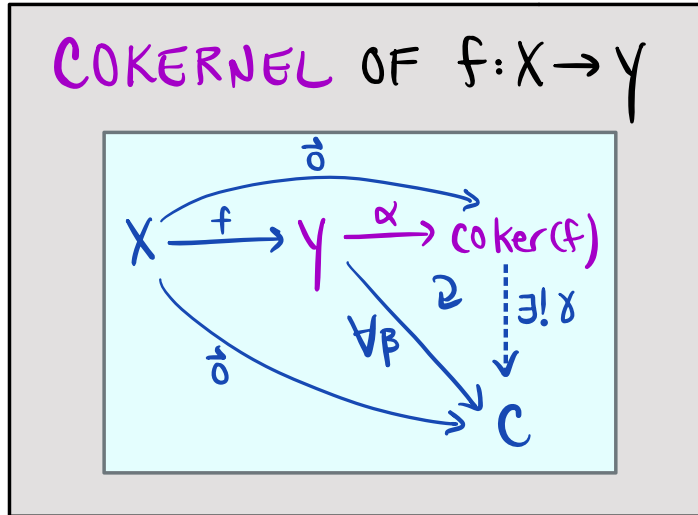
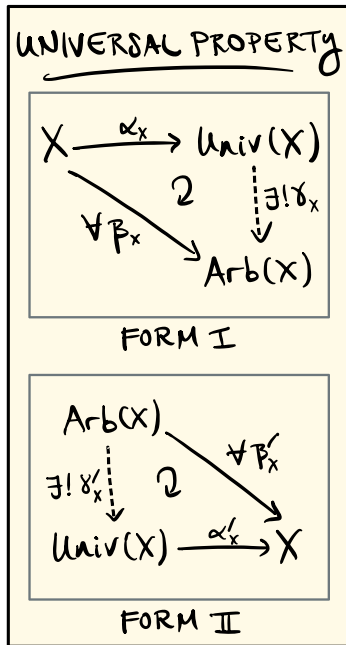
NOW

$[(h\alpha)f = h(\alpha f) = h\vec{0} = \vec{0}]$

By UNIQUENESS OF  $\delta$ ,  
 WE MUST HAVE  $h = h'$ .  $\parallel$

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

GIVEN A CATEGORY  $\mathcal{C}$  w/ ZERO OBJ.: ↗ OPERATION ON MORPHISMS



$\Downarrow$   
 $\alpha$  EPIC

$\Downarrow$   
 $\alpha'$  MONIC

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Ex.

$$\vec{0}_{x,y}: X \xrightarrow{x \vec{0}} 0 \xrightarrow{\vec{0}_y} Y$$

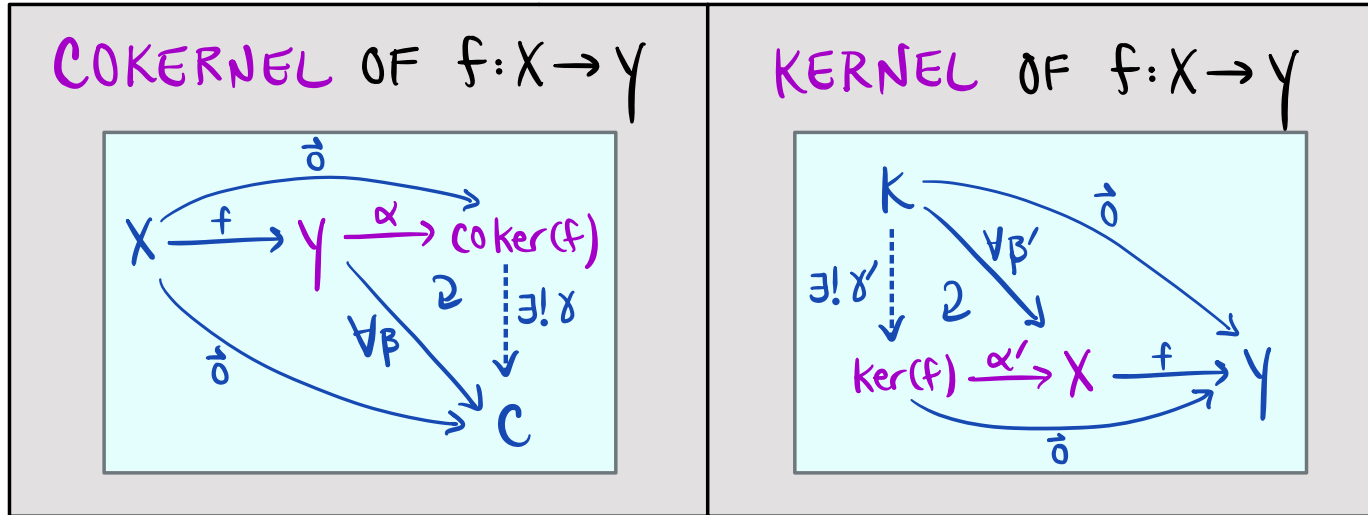
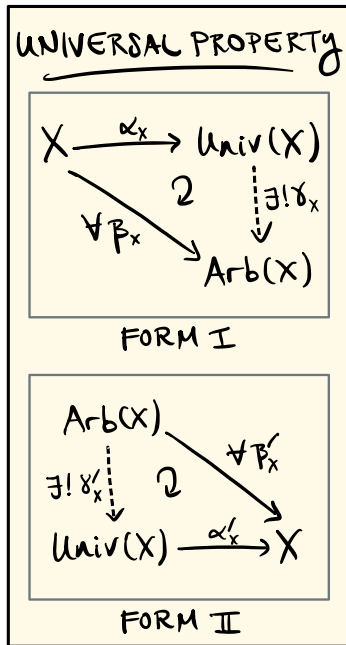
$$\Downarrow$$

$$f \cdot \vec{0} = \vec{0} \neq \vec{0} \cdot f = \vec{0}$$

$\forall f \in \mathcal{C}$

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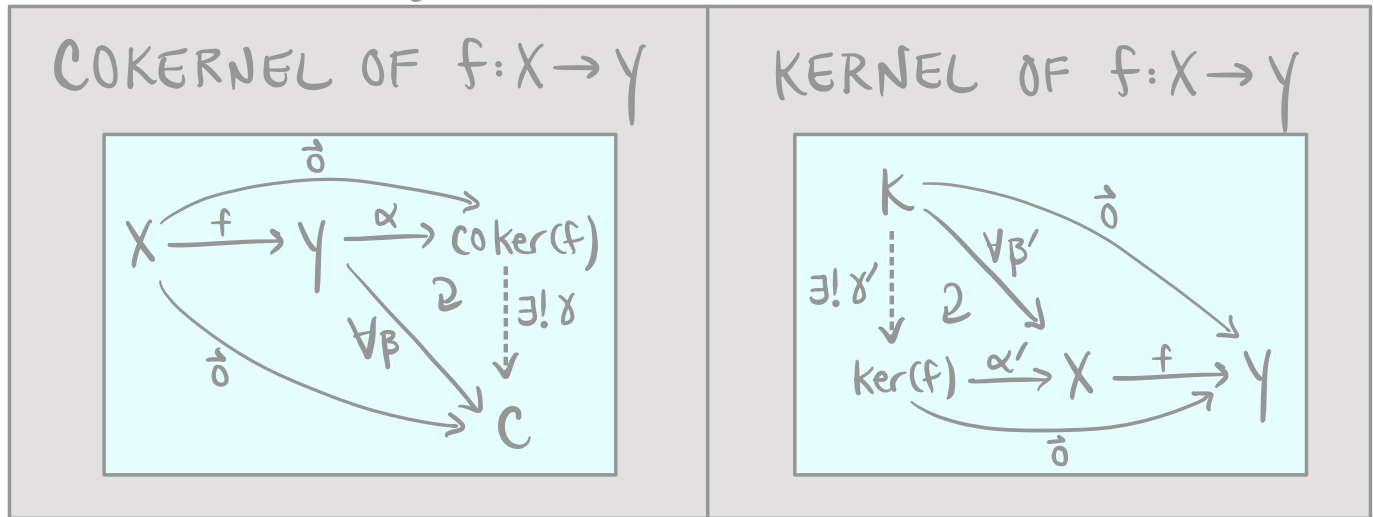
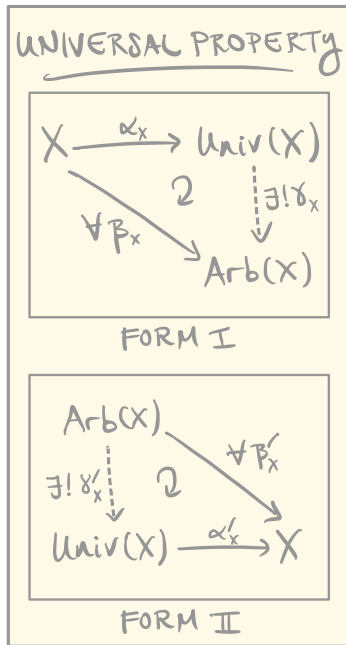
$\forall f \in \mathcal{C}$

} AN EPI IS CALLED NORMAL IF IT ARISES  
{ A MONO

AS THE { COKERNEL  
 KERNEL } OF A MORPHISM.

# I. UNIVERSAL CONSTRUCTIONS : COKERNELS & KERNELS

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{ AN EPI IS CALLED NORMAL IF IT ARISES  
A MONO }  
AS THE { COKERNEL  
KERNEL } OF A MORPHISM.

≡ NEXT WE DISCUSS CATEGORIES IN WHICH ≡  
ALL OF THE ABOVE EXISTS



## II. ABELIAN CATEGORIES : OVERVIEW

TAKE  $\mathcal{C}$  A CATEGORY

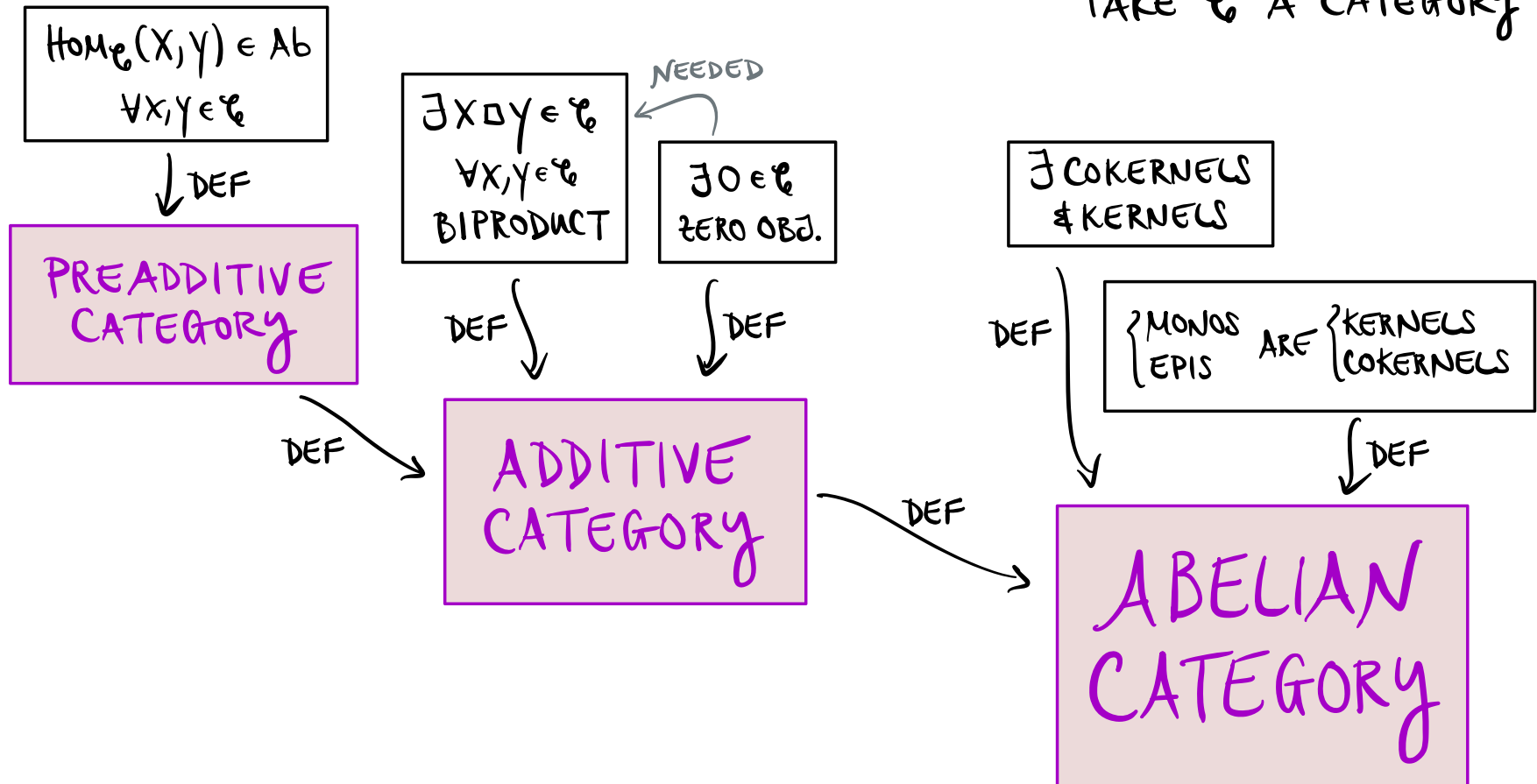
PREADDITIVE  
CATEGORY

ADDITIVE  
CATEGORY

ABELIAN  
CATEGORY

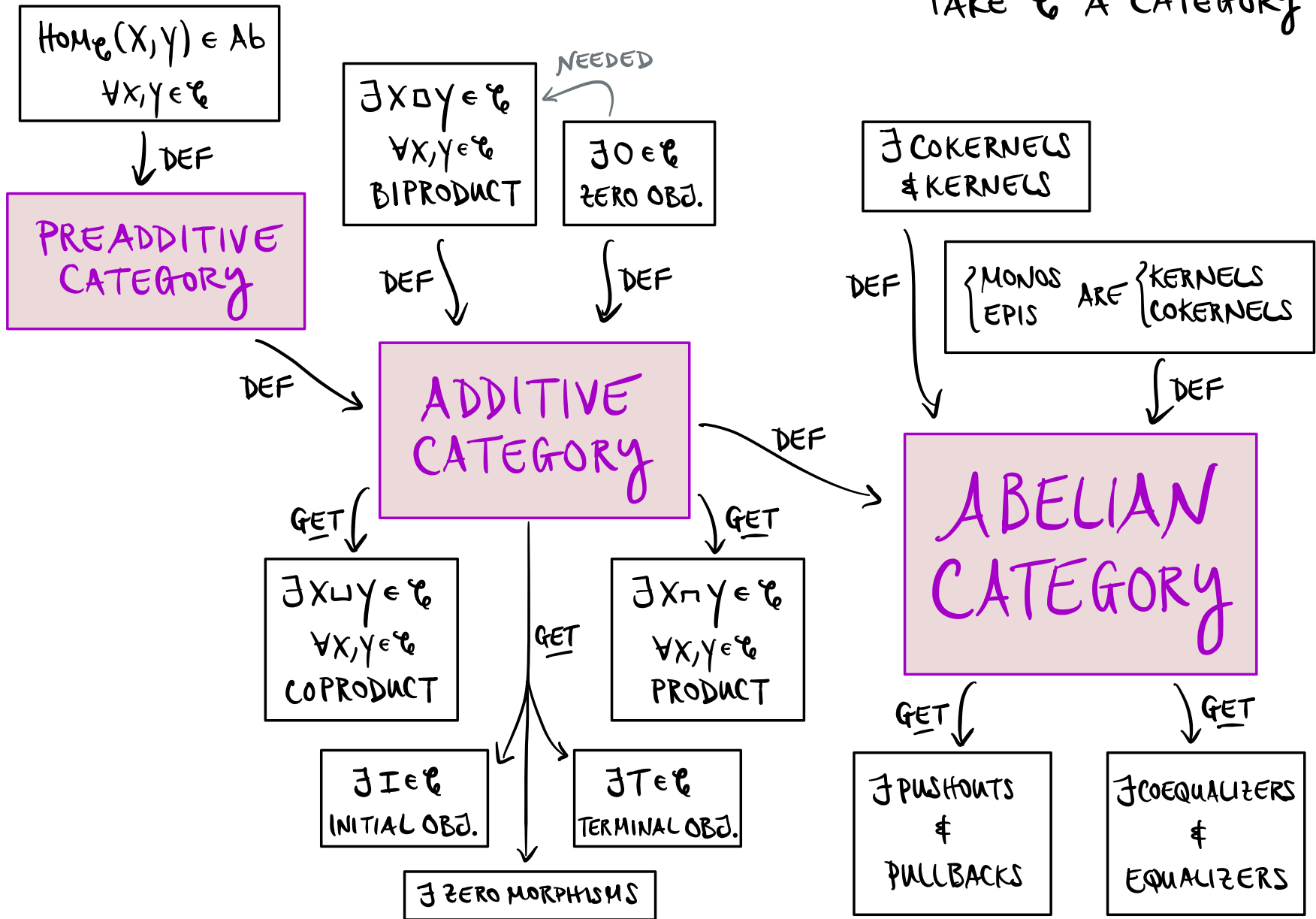
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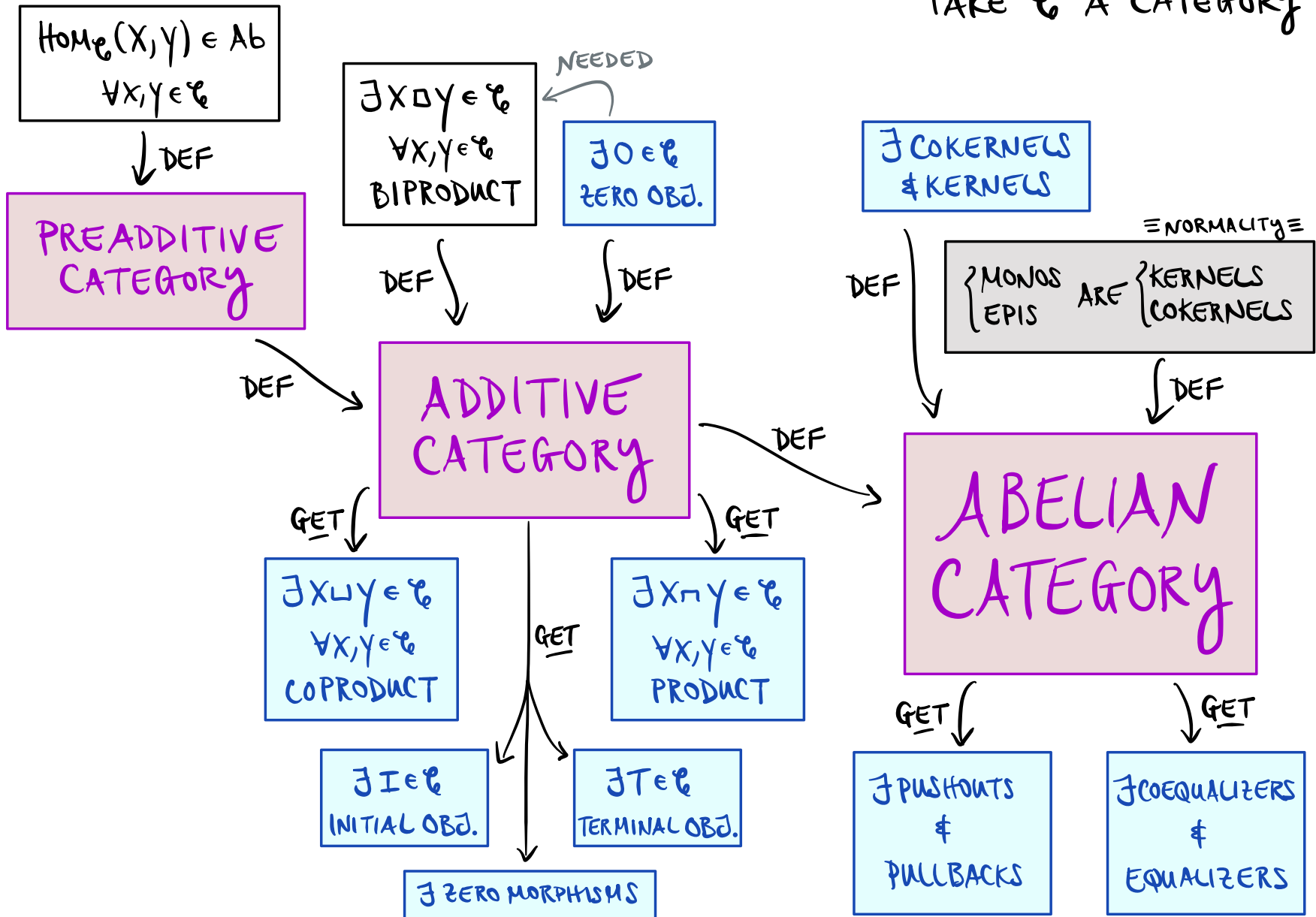
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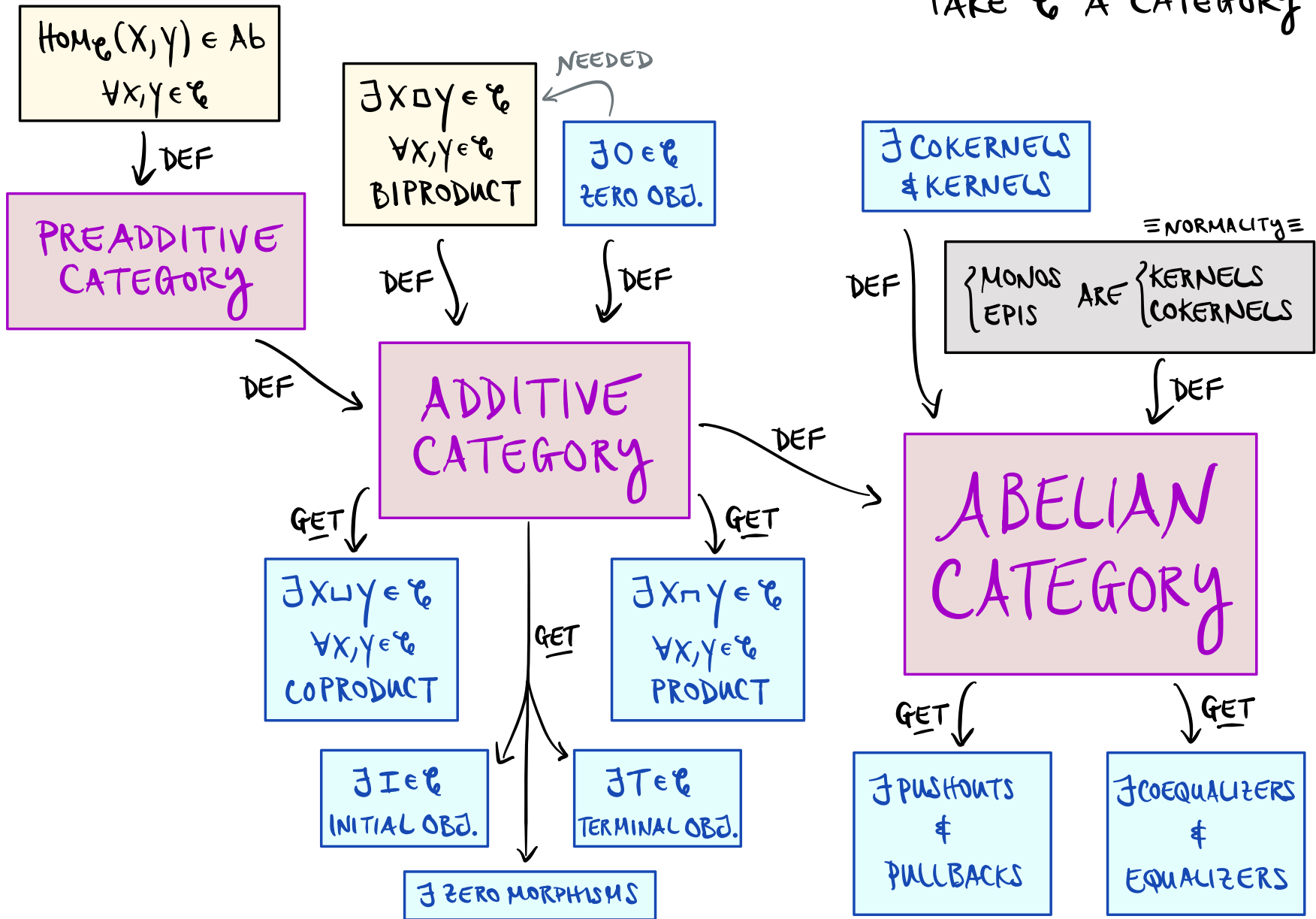
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## II. ABELIAN CATEGORIES : OVERVIEW

TAKE  $\mathcal{C}$  A CATEGORY



## II. ABELIAN CATEGORIES : PREADDITIVE & LINEAR CATEGORIES

$$\begin{array}{l} \text{Hom}_{\mathcal{C}}(X, Y) \in \text{Ab} \\ \forall X, Y \in \mathcal{C} \end{array}$$

↓ DEF

PREADDITIVE  
CATEGORY

## II. ABELIAN CATEGORIES : PREADDITIVE & LINEAR CATEGORIES

$$\text{Hom}_{\mathcal{C}}(X, Y) \in \text{Ab} \\ \forall X, Y \in \mathcal{C}$$

↓ DEF

PREADDITIVE  
CATEGORY

A CATEGORY  $\mathcal{C}$  IS PREADDITIVE

(OR AN AB-CATEGORY,  
OR ENRICHED OVER AB)

IF  $\text{Hom}_{\mathcal{C}}(X, Y)$  IS AN ABELIAN GROUP  $\forall X, Y \in \mathcal{C}$

w/ } OPERATION +

ADDITIVE IDENTITY  $\vec{0}$  (WHEN  $\exists 0 \in \mathcal{C}$ )

ADDITIVE INVERSE OF  $f: X \rightarrow Y$  DENOTED BY  $-f: X \rightarrow Y$

ALSO REQUIRE • DISTRIBUTES OVER +

## II. ABELIAN CATEGORIES : PREADDITIVE & LINEAR CATEGORIES

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PREADDITIVE  
CATEGORY

Ex. Ab

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ALSO REQUIRE  $\circ$  DISTRIBUTES OVER +

FIX  
GROUND  
FIELD  
 $\mathbb{R}$   
↳

A CATEGORY  $\mathcal{C}$  IS LINEAR (OR ENRICHED OVER  $\text{Vec}$ )

IF  $\mathcal{C}$  IS PREADDITIVE

$\&$   $\text{Hom}_{\mathcal{C}}(X, Y)$  IS A VECTOR SPACE  $\forall X, Y \in \mathcal{C}$

$\&$   $\circ$  DISTRIBUTES OVER  $*$ .

## II. ABELIAN CATEGORIES : PREADDITIVE & LINEAR CATEGORIES

$$\text{Hom}_{\mathcal{C}}(X, Y) \in \text{Ab}$$

$$\forall X, Y \in \mathcal{C}$$

↓ DEF

PREADDITIVE  
CATEGORY

Ex. Ab

Ex.  $\text{Vec}_{\mathbb{R}}$   
A-Mod

FIX  
GROUND  
FIELD  
 $\mathbb{R}$

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ALSO REQUIRE  $\circ$  DISTRIBUTES OVER +

A CATEGORY  $\mathcal{C}$  IS LINEAR (OR ENRICHED OVER  $\text{Vec}$ )

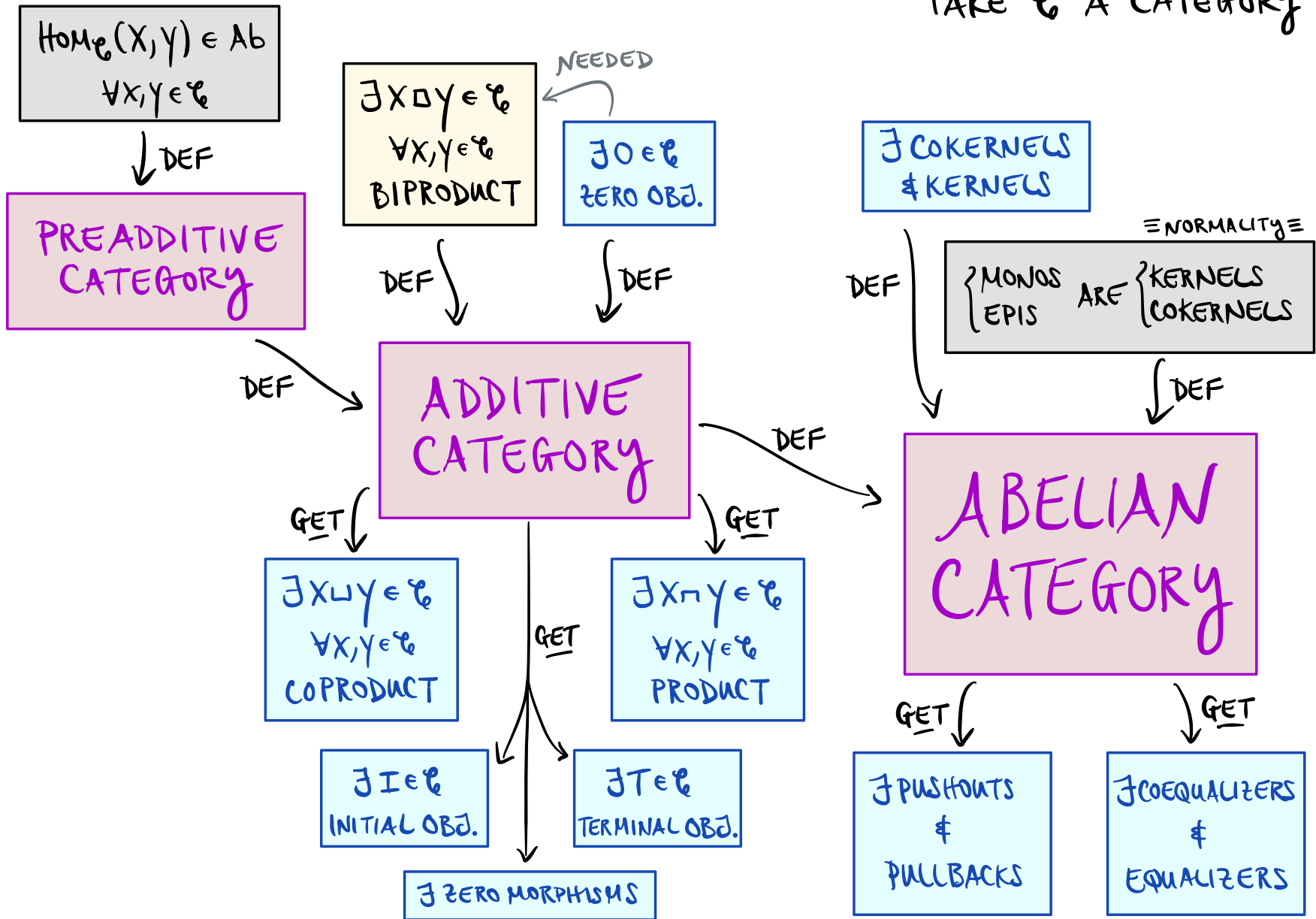
IF  $\mathcal{C}$  IS PREADDITIVE

$\&$   $\text{Hom}_{\mathcal{C}}(X, Y)$  IS A VECTOR SPACE  $\forall X, Y \in \mathcal{C}$

$\&$   $\circ$  DISTRIBUTES OVER  $*$ .

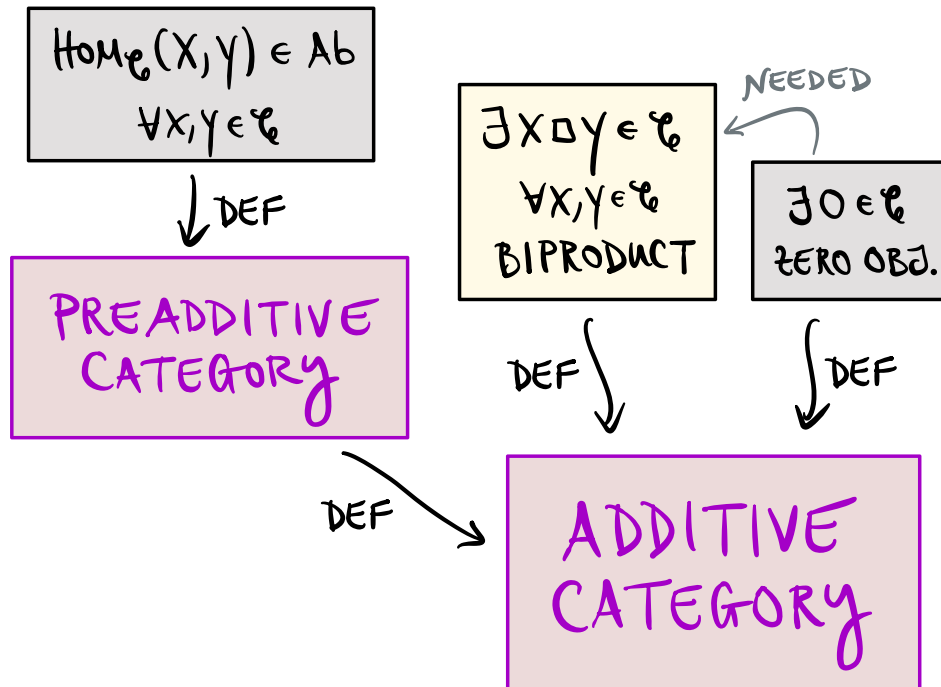
## II. ABELIAN CATEGORIES : OVERVIEW

TAKE  $\mathcal{C}$  A CATEGORY

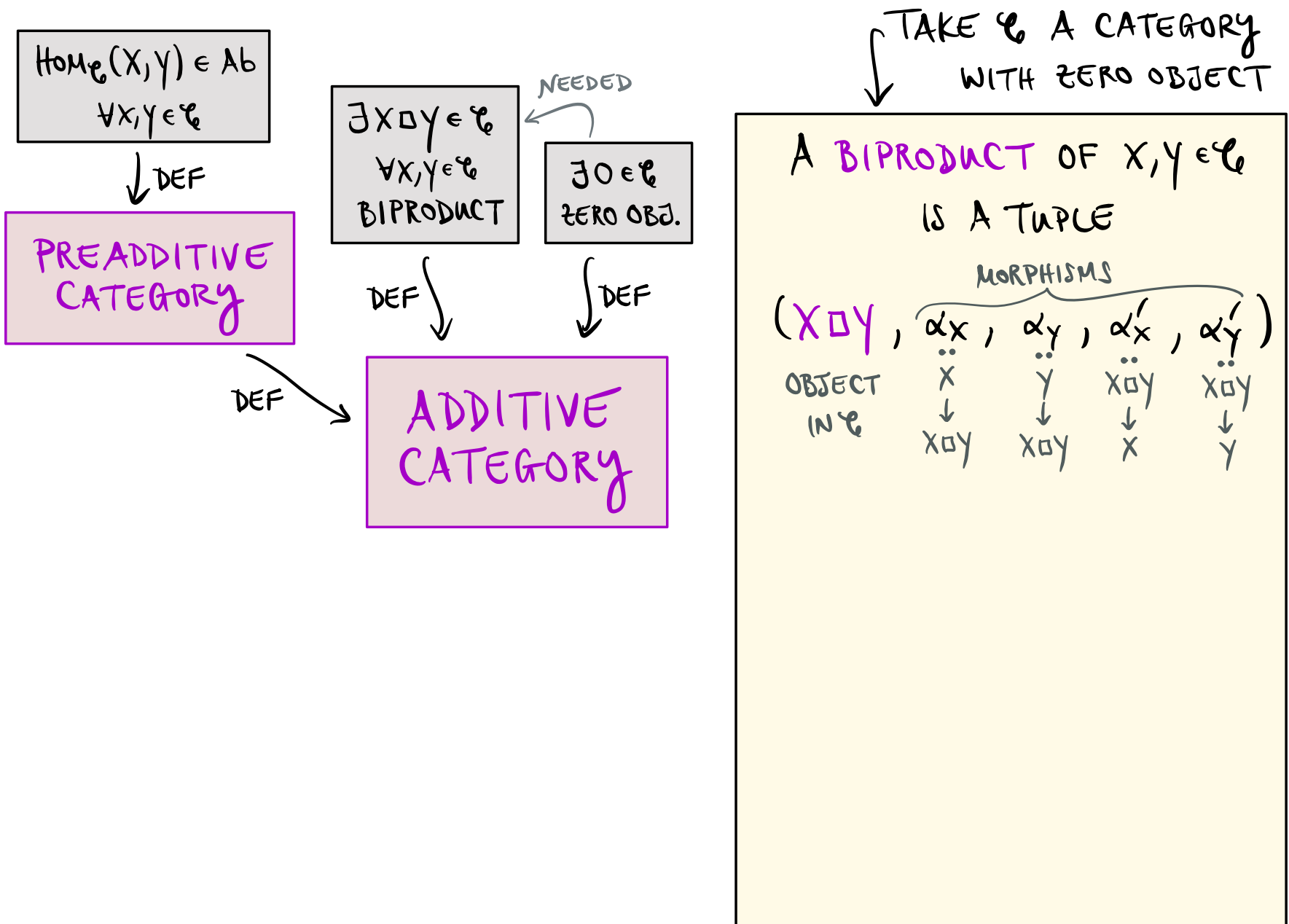


## II. ABELIAN CATEGORIES : OVERVIEW

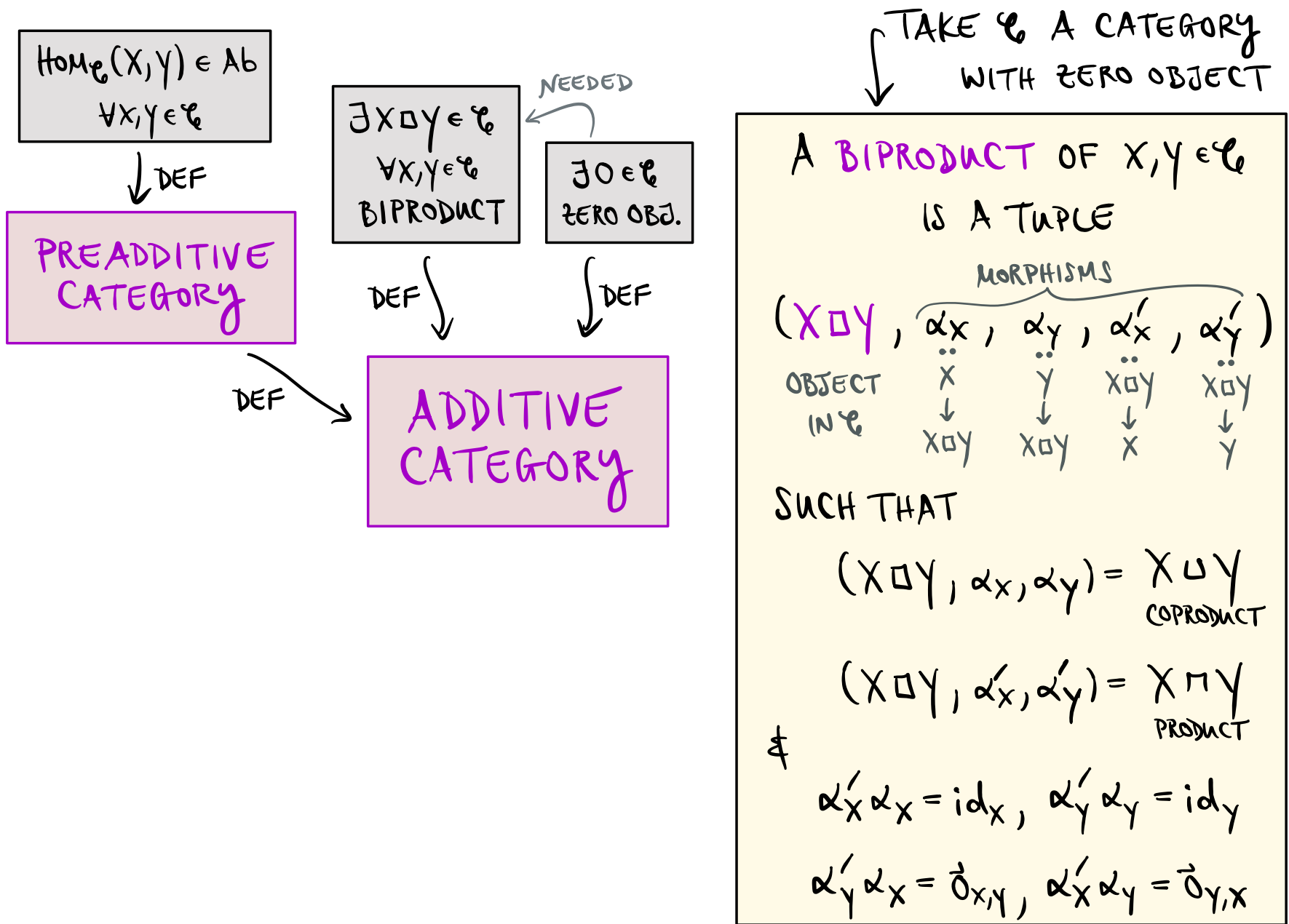
TAKE  $\mathcal{C}$  A CATEGORY



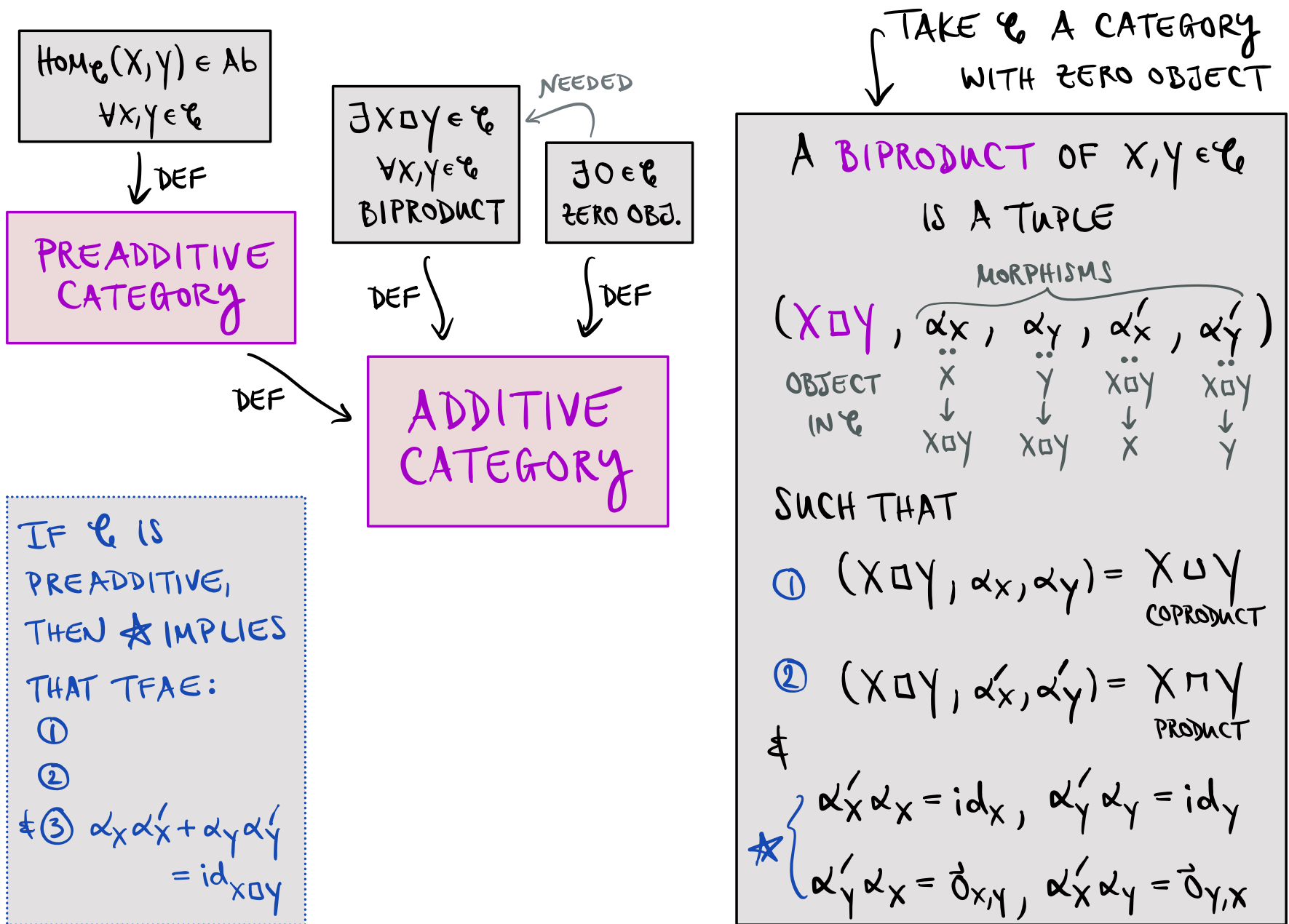
## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES



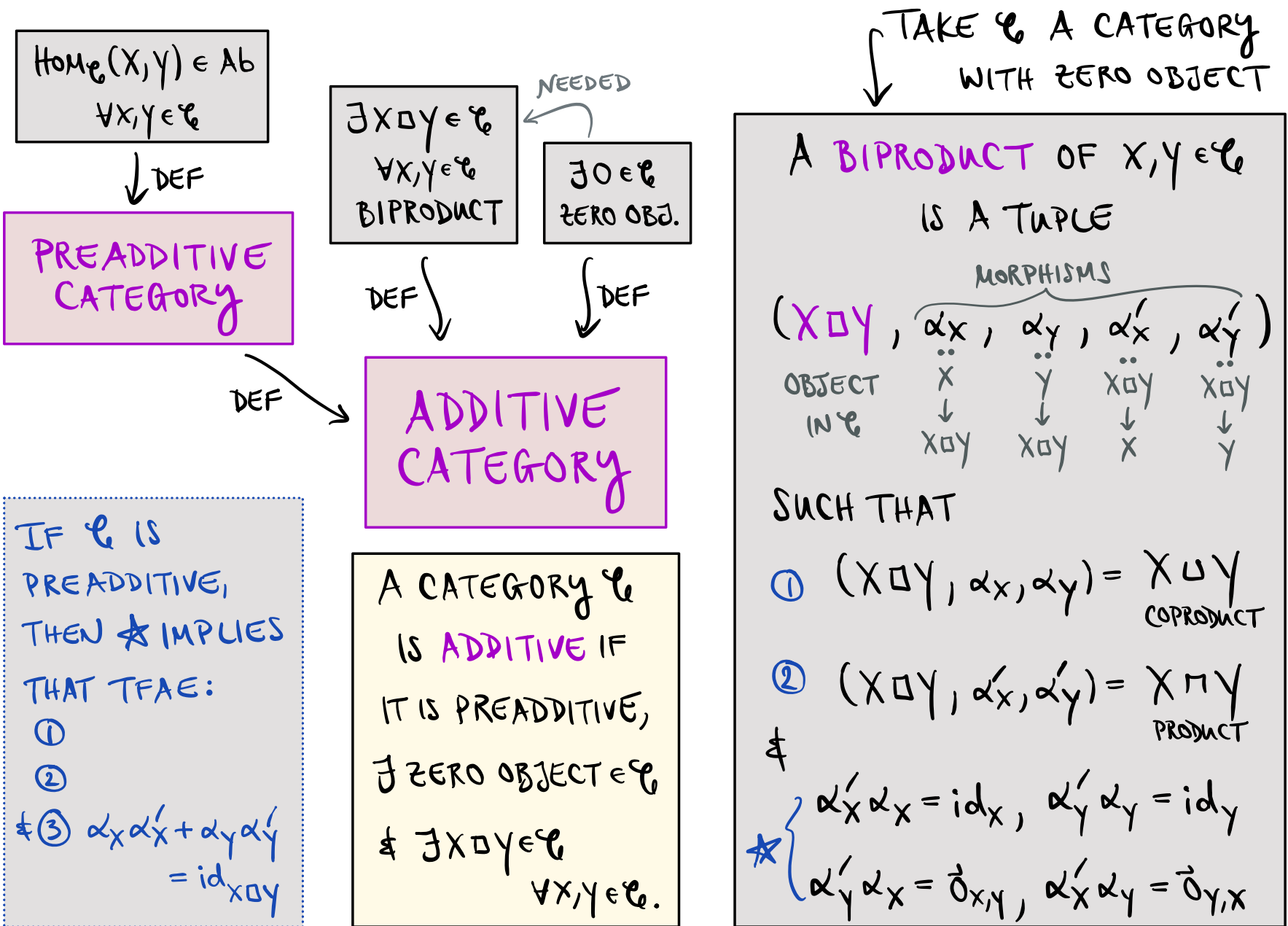
## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES



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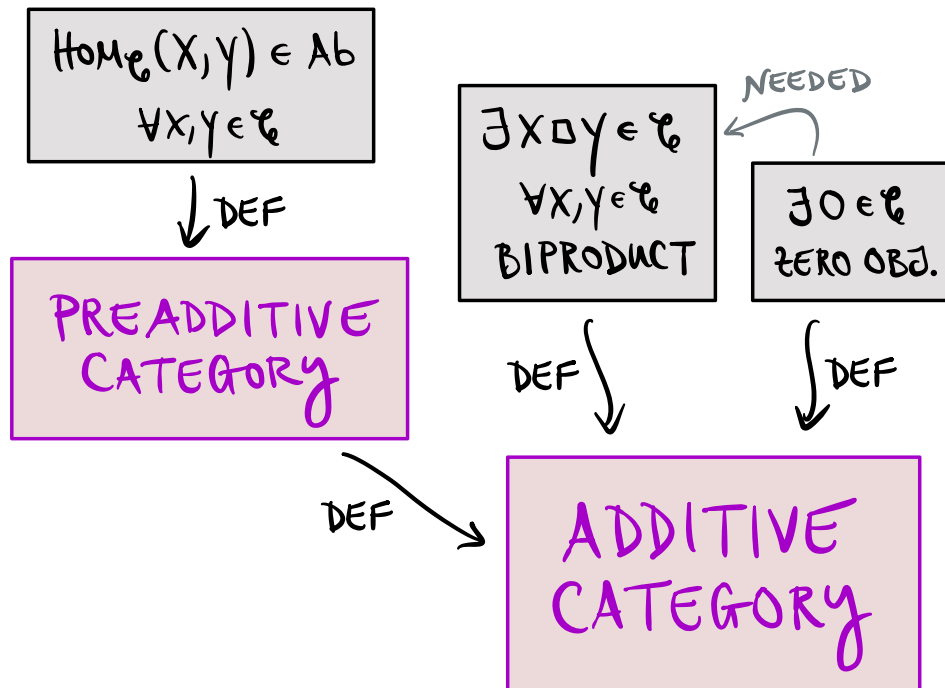


## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES





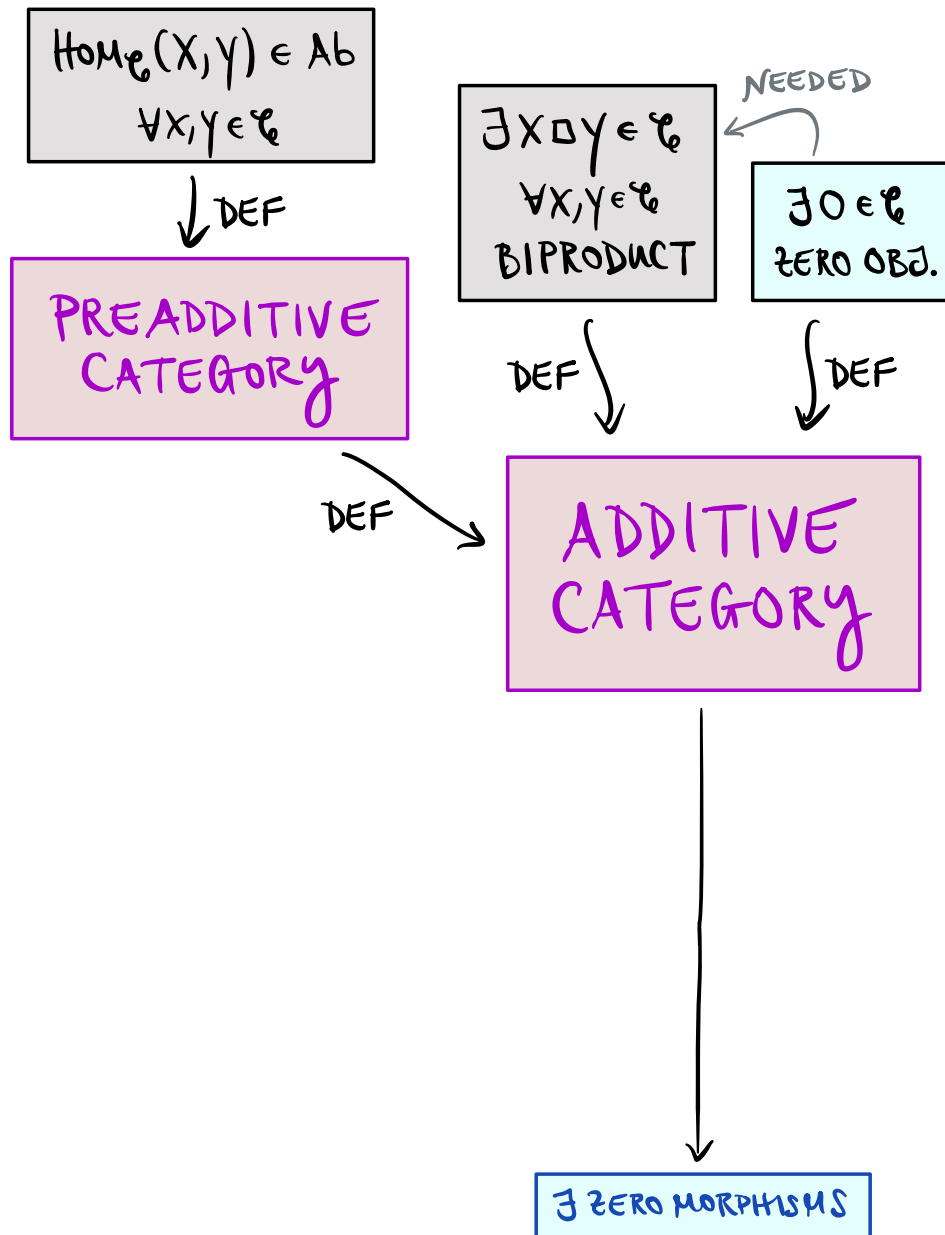
## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES



TAKE  $\mathcal{C}$  A CATEGORY  
 WITH ZERO OBJECT

A **BIPRODUCT** OF  $X, Y \in \mathcal{C}$   
 IS A TUPLE  
 $(X \sqcup Y, \alpha_X, \alpha_Y, \alpha'_X, \alpha'_Y)$   
 $\exists$ .  
 $(X \sqcup Y, \alpha_X, \alpha_Y) = X \cup Y$   
 $(X \sqcup Y, \alpha'_X, \alpha'_Y) = X \cap Y$   
 $\&$   
 $\alpha'_X \alpha_X = \text{id}_X, \alpha'_Y \alpha_Y = \text{id}_Y$   
 $\alpha'_Y \alpha_X = \vec{0}_{X, Y}, \alpha'_X \alpha_Y = \vec{0}_{Y, X}$

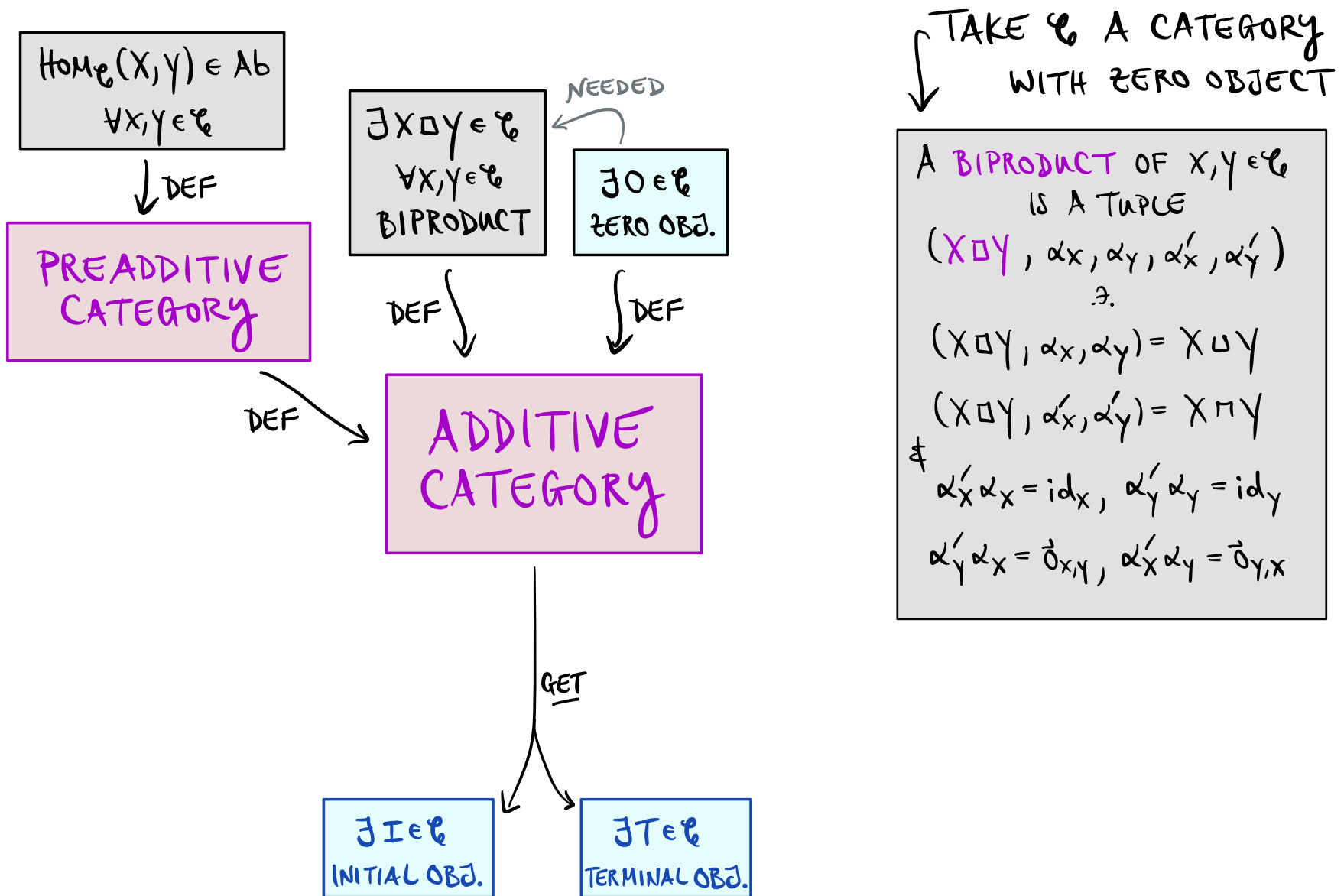
## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES



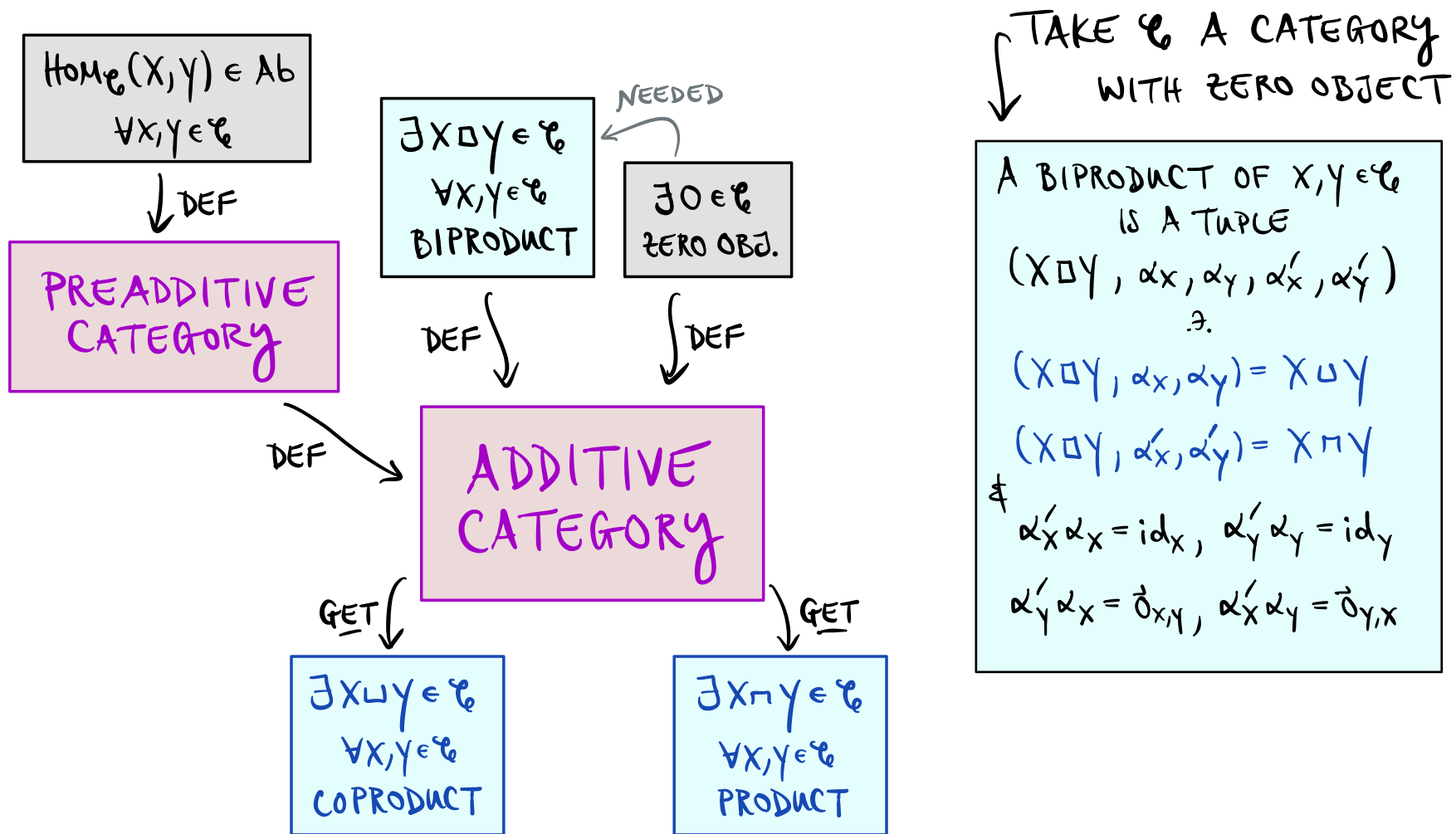
TAKE  $\mathcal{C}$  A CATEGORY  
 WITH ZERO OBJECT

A **BIPRODUCT** OF  $X, Y \in \mathcal{C}$   
 IS A TUPLE  
 $(X \square Y, \alpha_X, \alpha_Y, \alpha'_X, \alpha'_Y)$   
 $\exists$ .  
 $(X \square Y, \alpha_X, \alpha_Y) = X \cup Y$   
 $(X \square Y, \alpha'_X, \alpha'_Y) = X \cap Y$   
 $\&$   
 $\alpha'_X \alpha_X = \text{id}_X, \alpha'_Y \alpha_Y = \text{id}_Y$   
 $\alpha'_Y \alpha_X = \vec{0}_{X, Y}, \alpha'_X \alpha_Y = \vec{0}_{Y, X}$

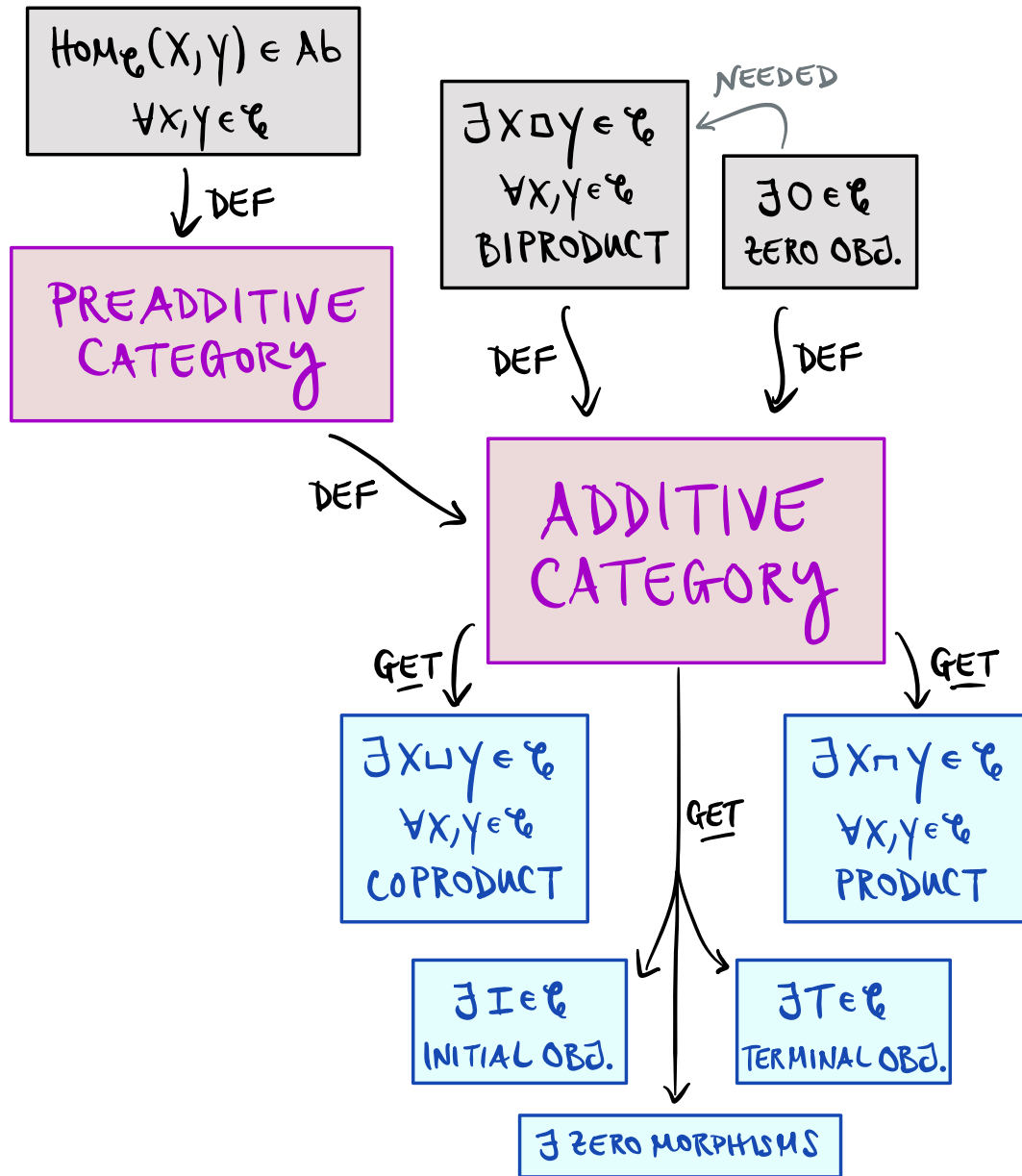
## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES



## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES

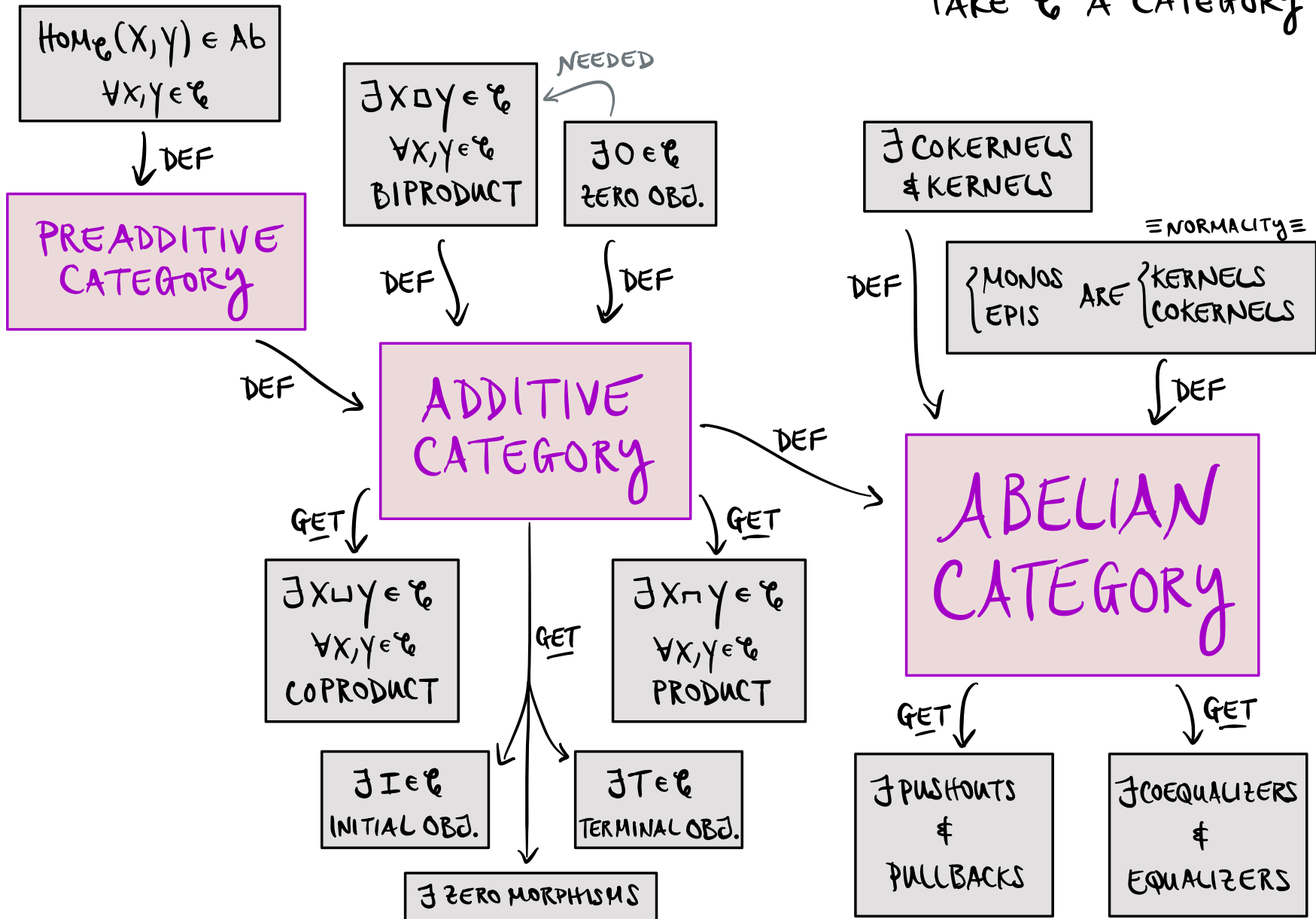


## II. ABELIAN CATEGORIES : ADDITIVE CATEGORIES



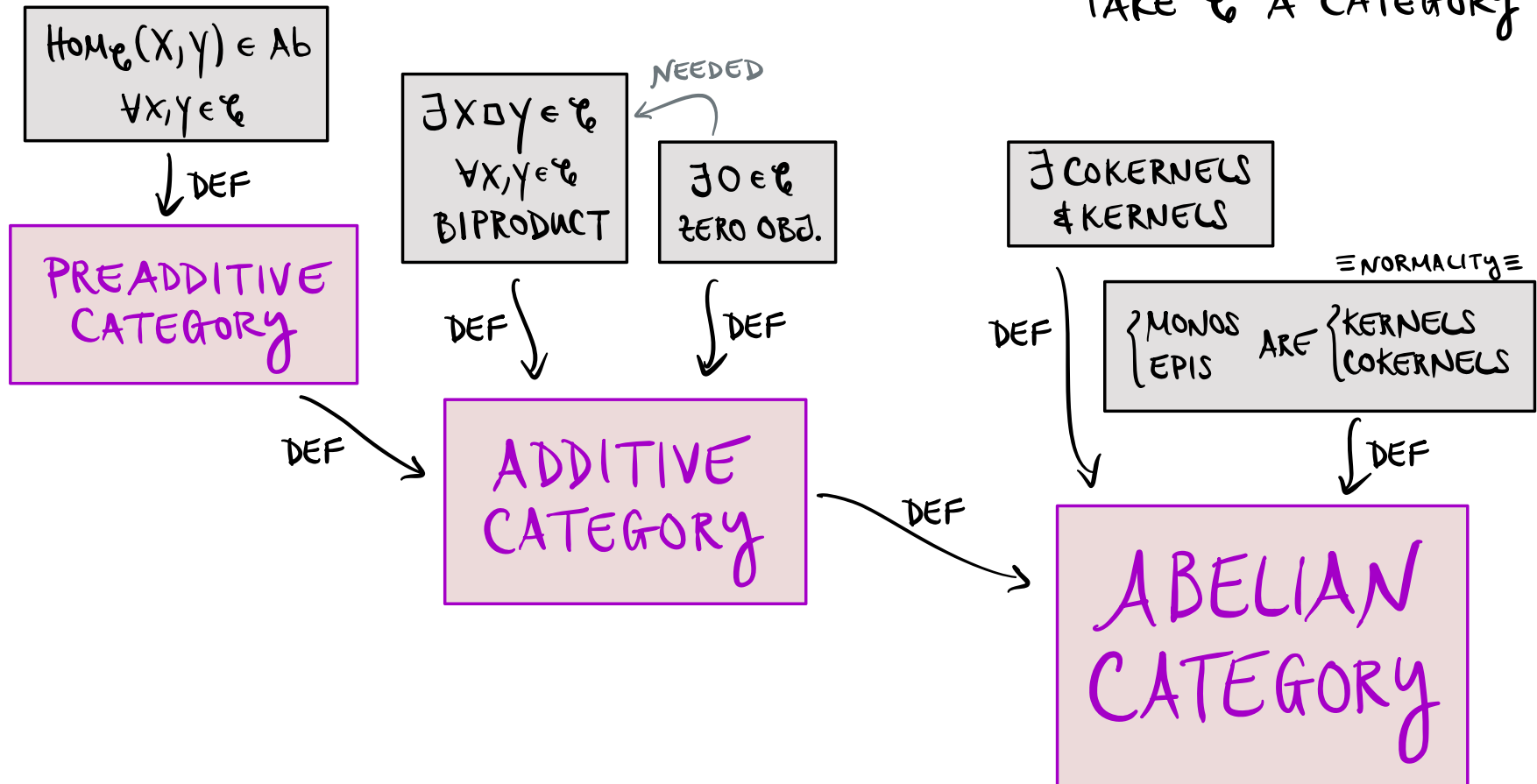
## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY



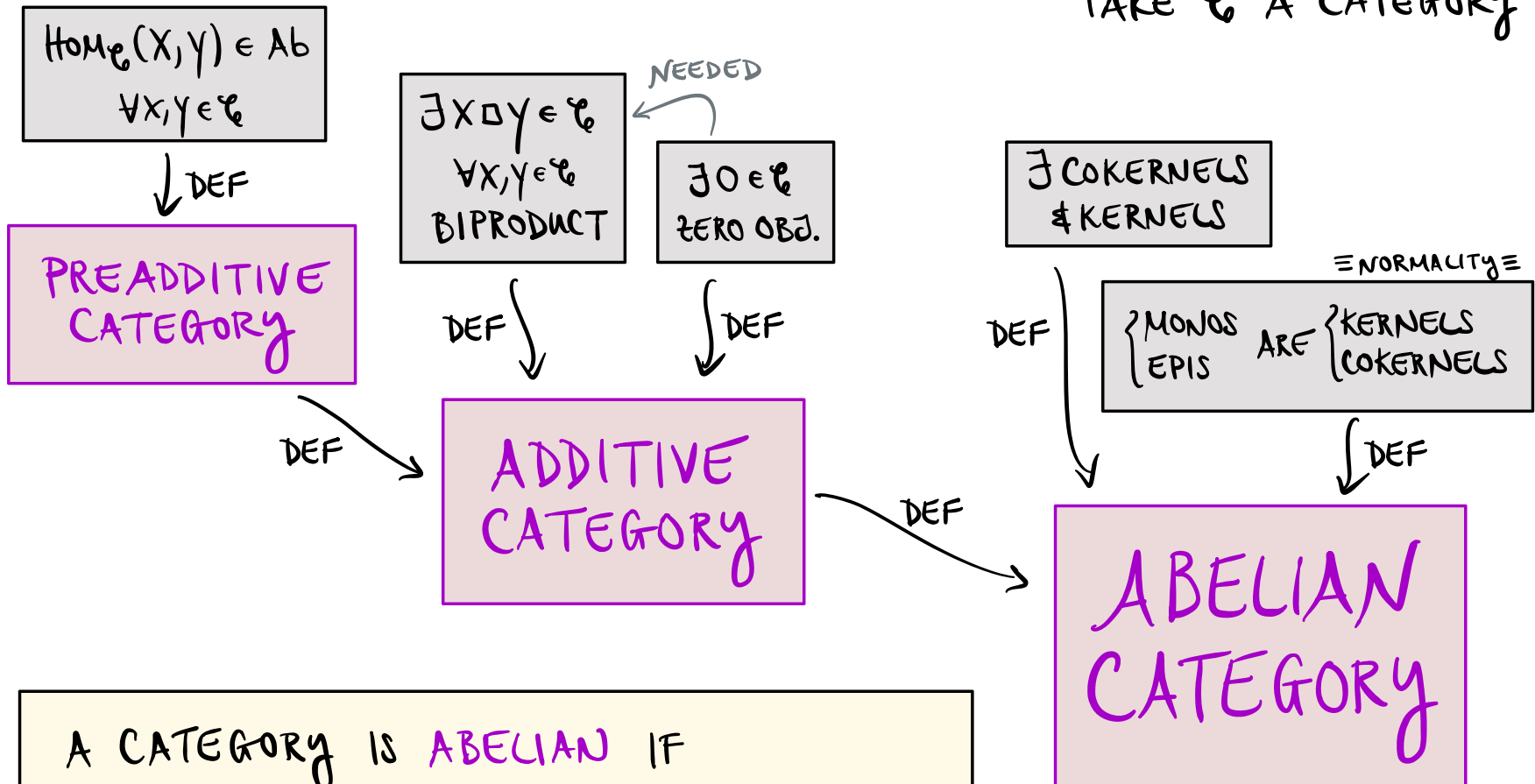
## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY



## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY

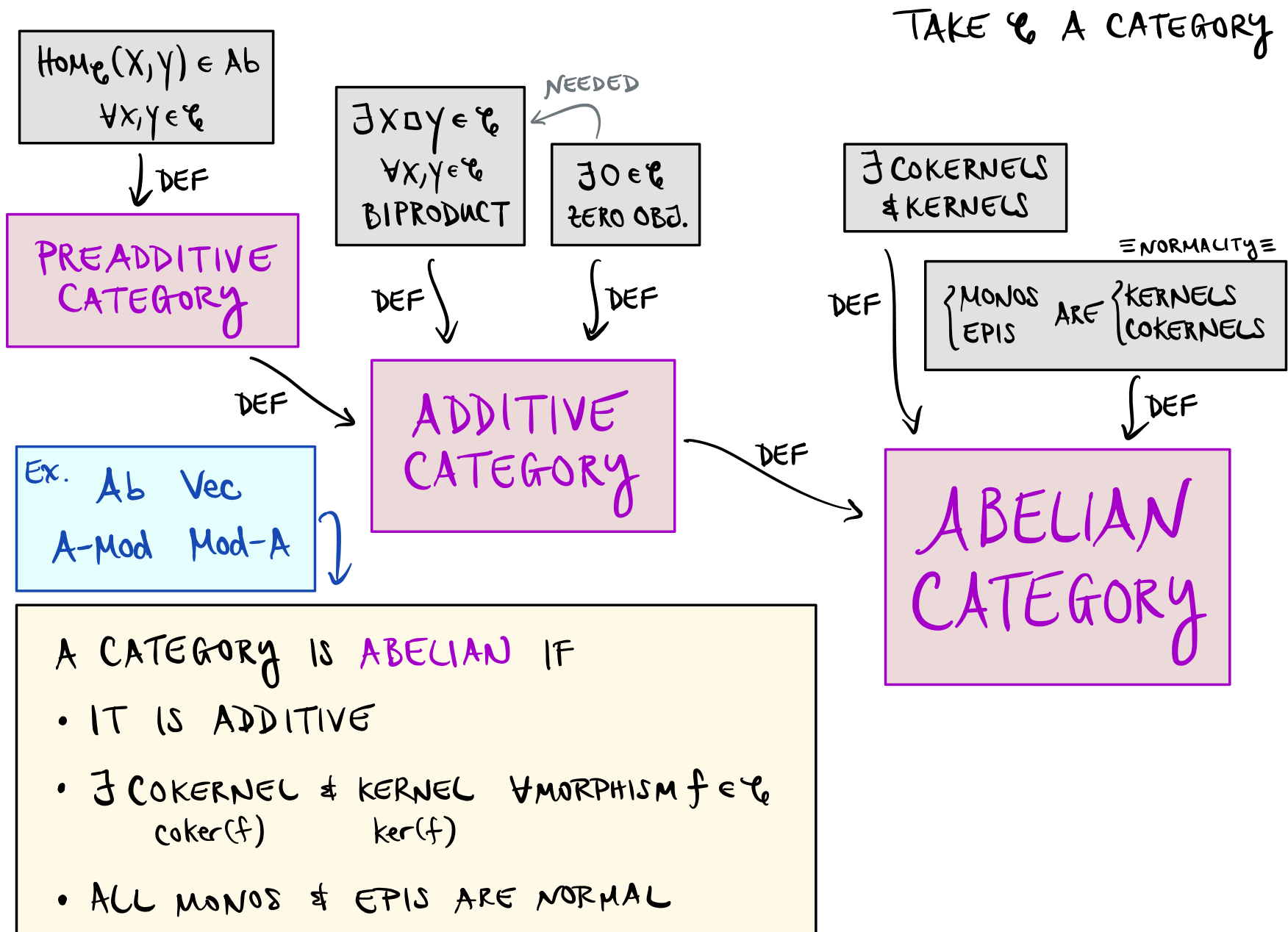


A CATEGORY IS **ABELIAN** IF

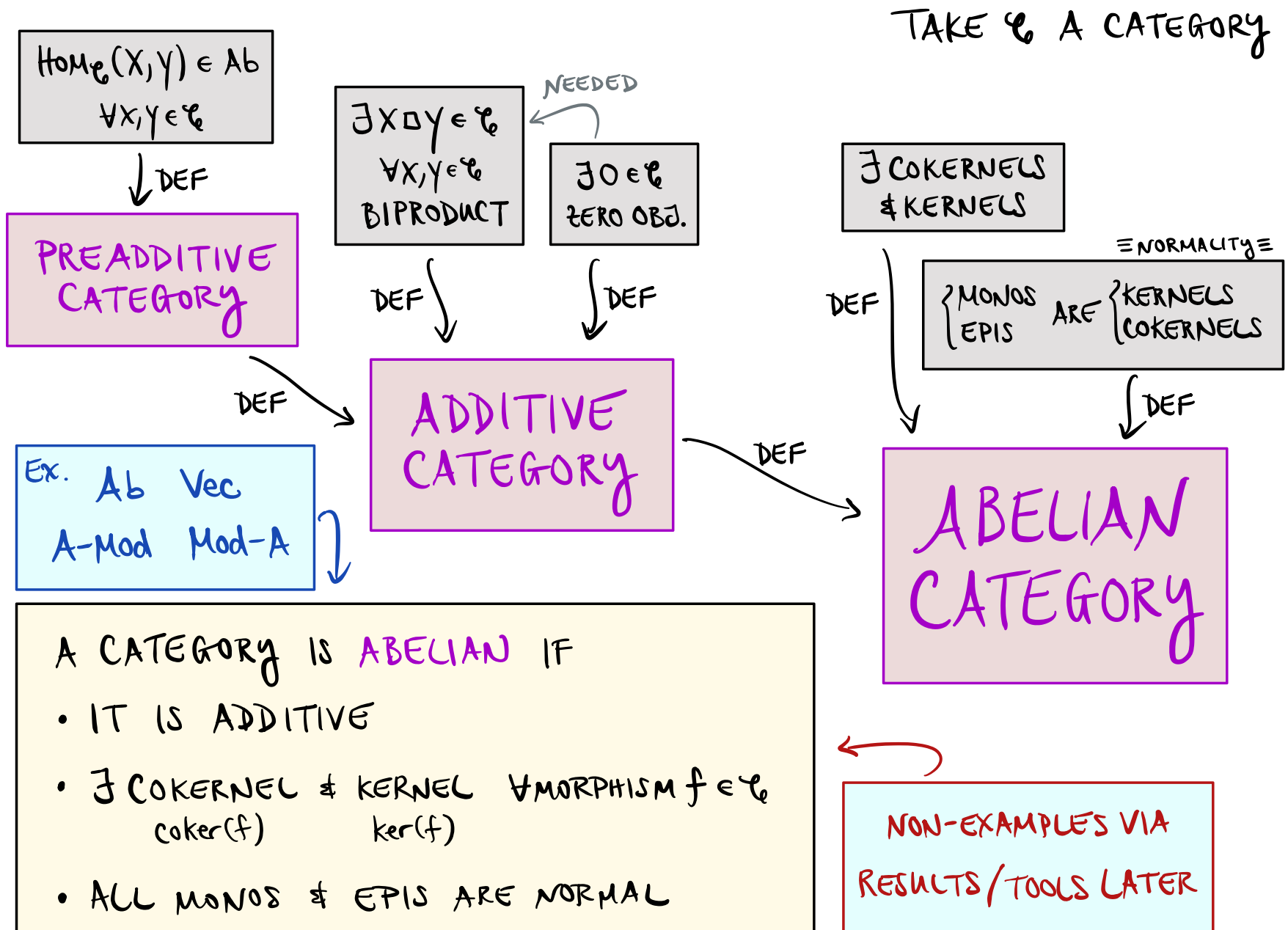
- IT IS ADDITIVE
- $\exists \text{COKERNEL} \ \& \ \text{KERNEL} \ \forall \text{MORPHISM } f \in \mathcal{C}$   
 $\text{coker}(f) \qquad \text{ker}(f)$
- ALL MONOS & EPIS ARE NORMAL



## II. ABELIAN CATEGORIES

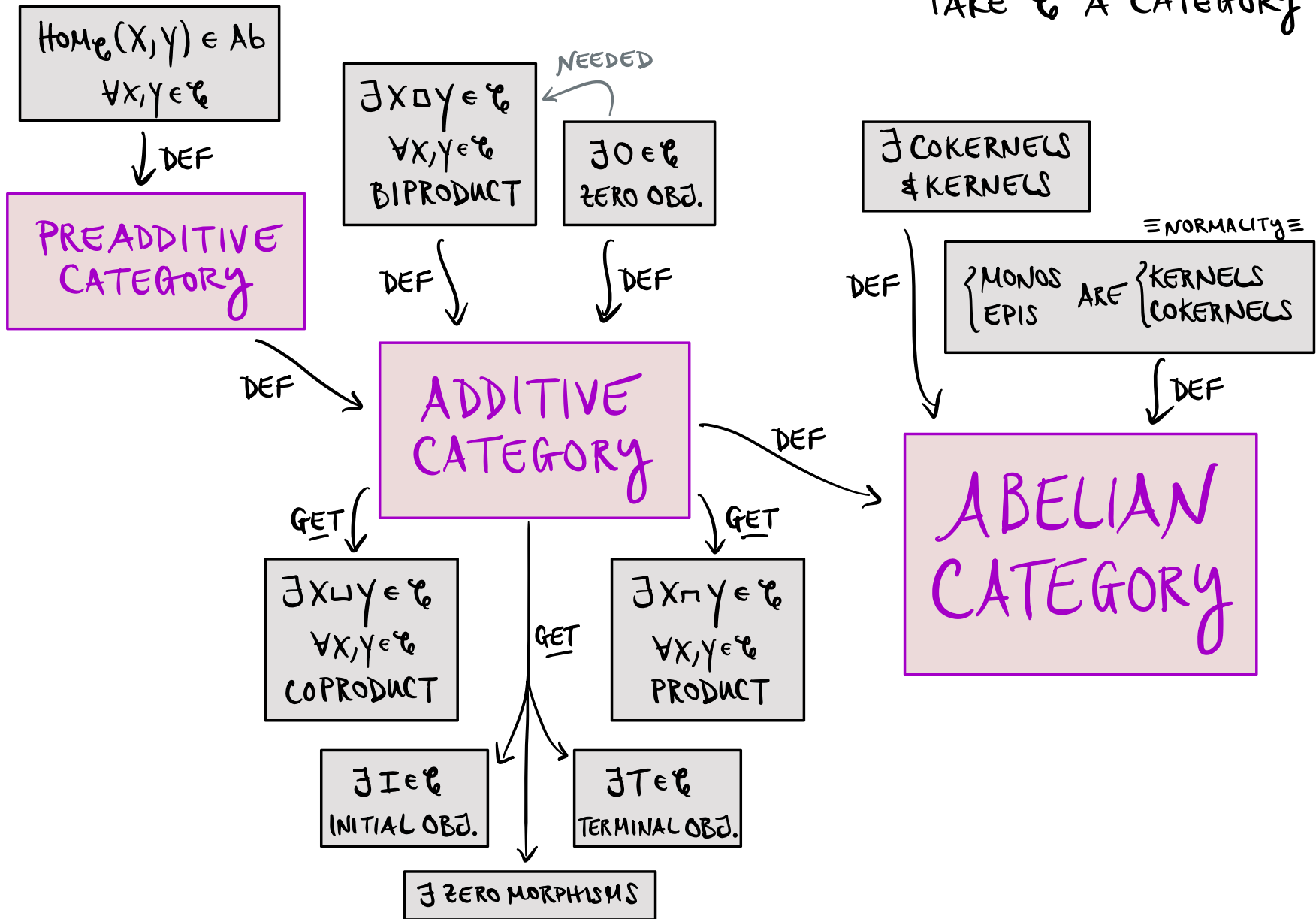


## II. ABELIAN CATEGORIES



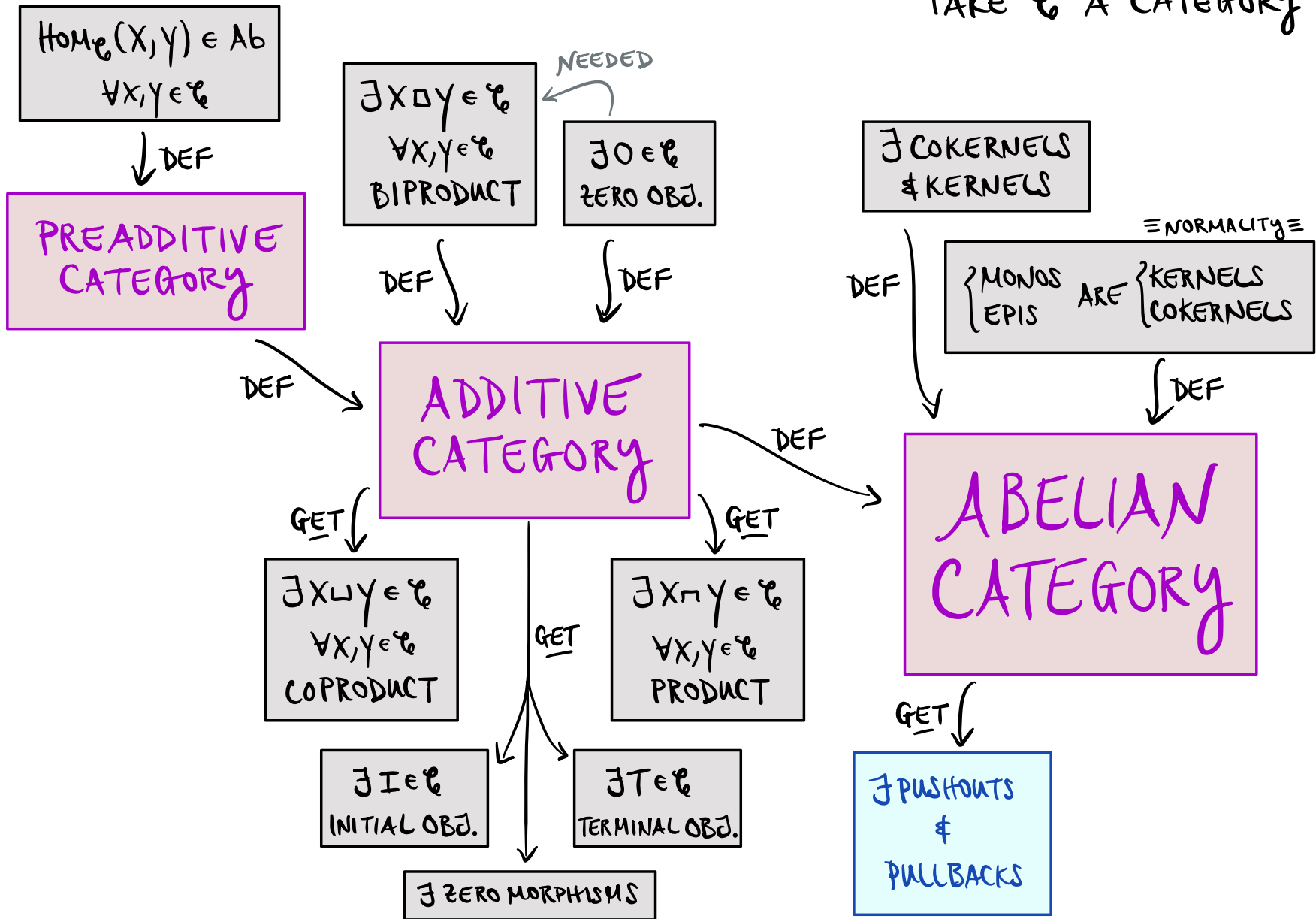
## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY



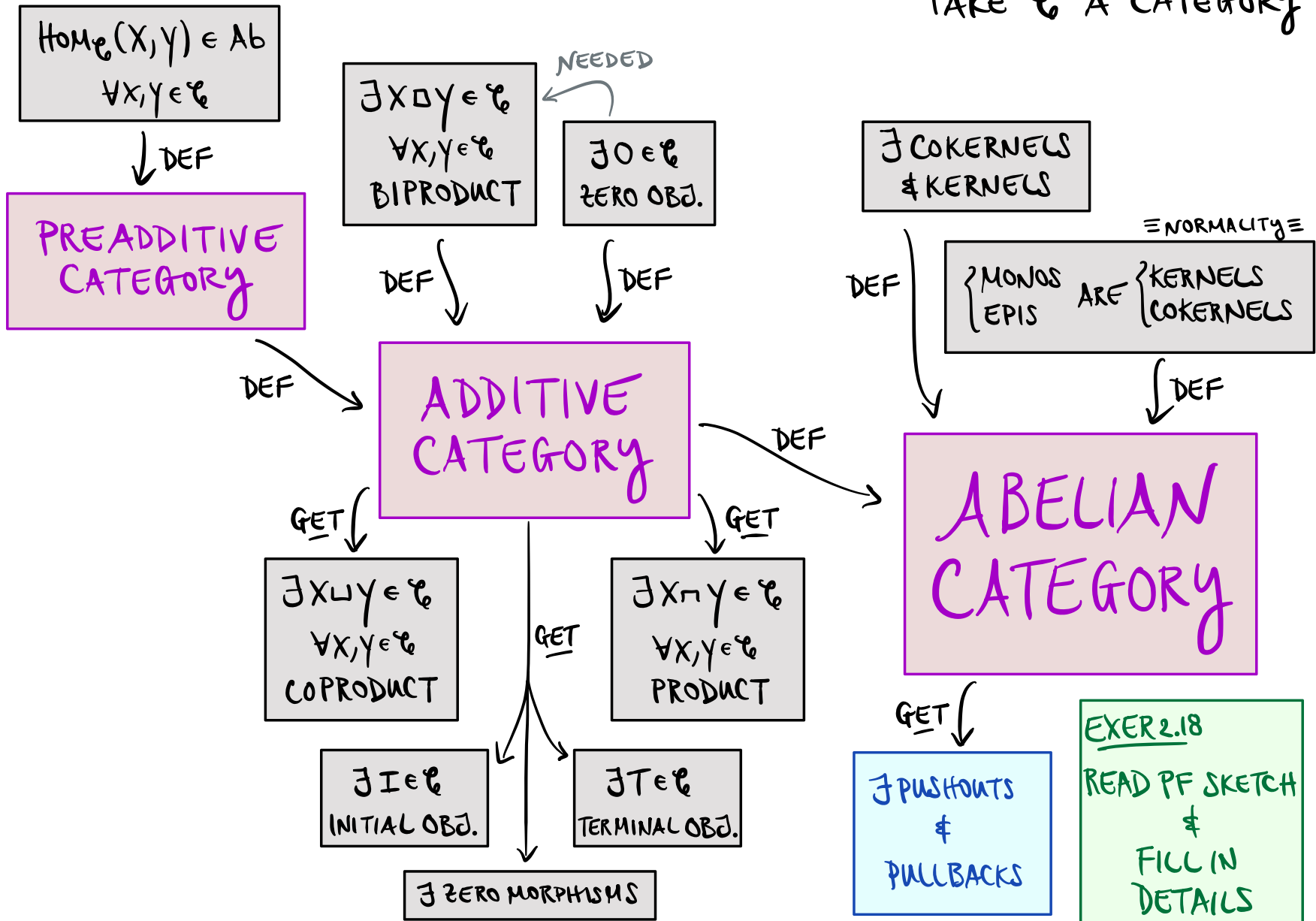
## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY

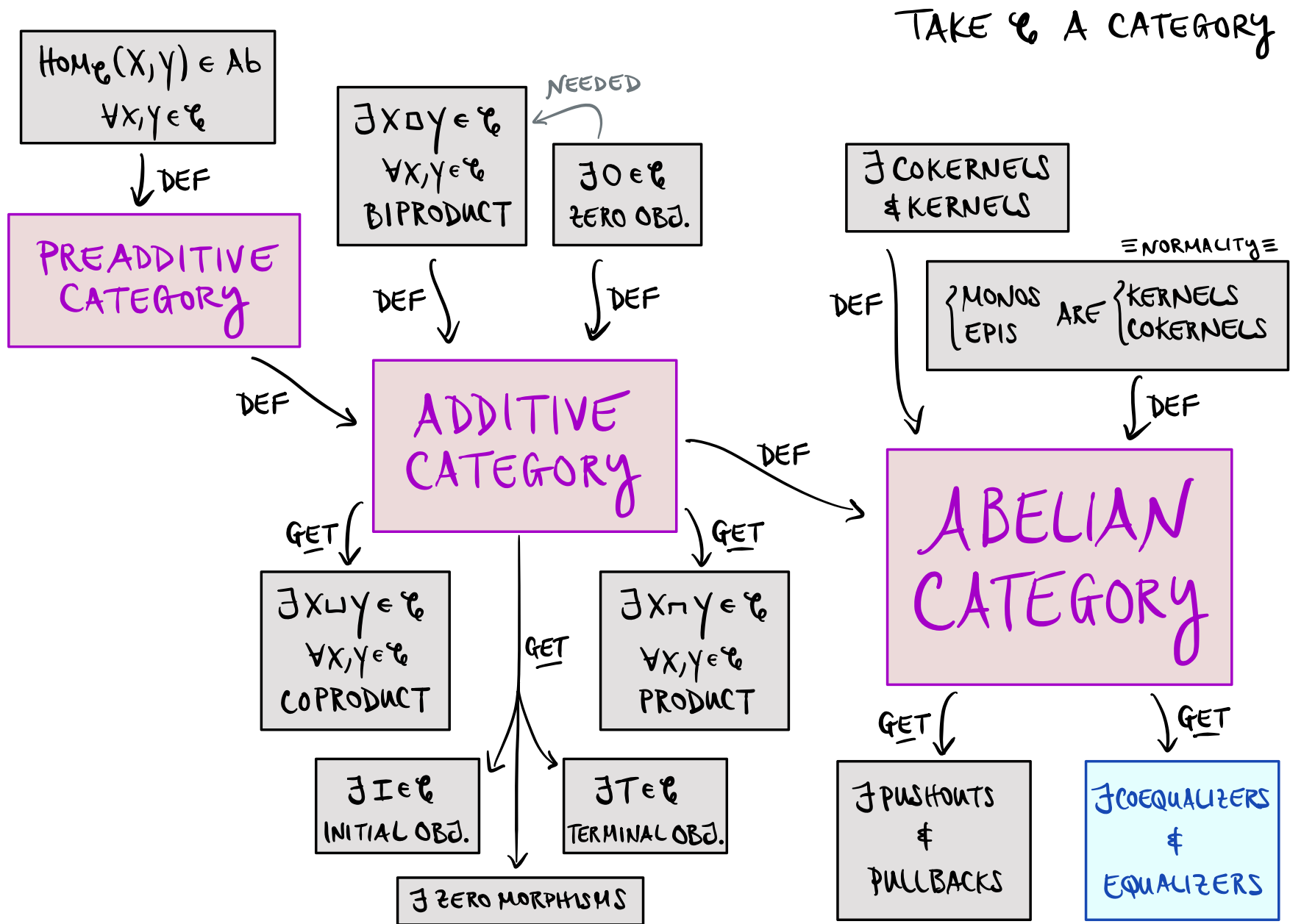


## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY

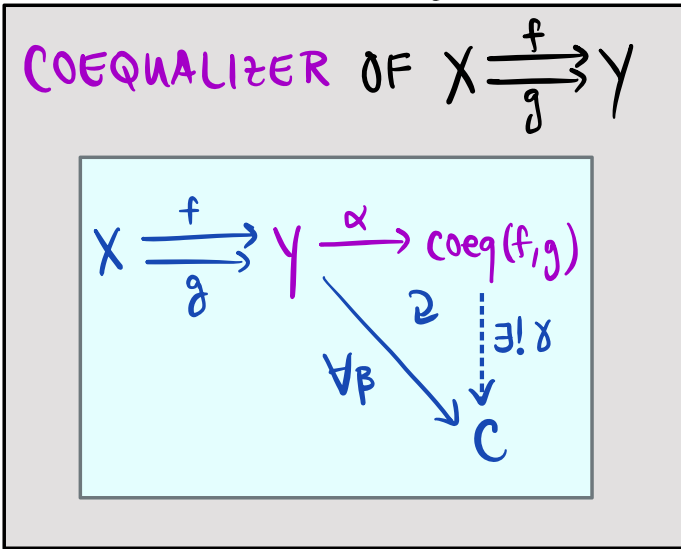


## II. ABELIAN CATEGORIES

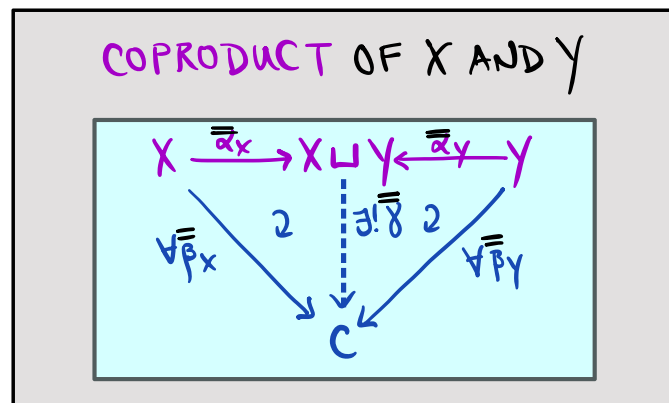
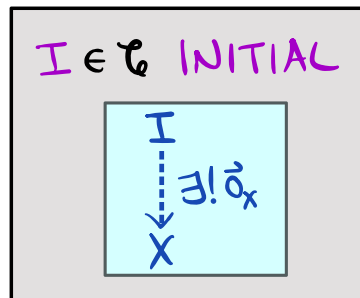
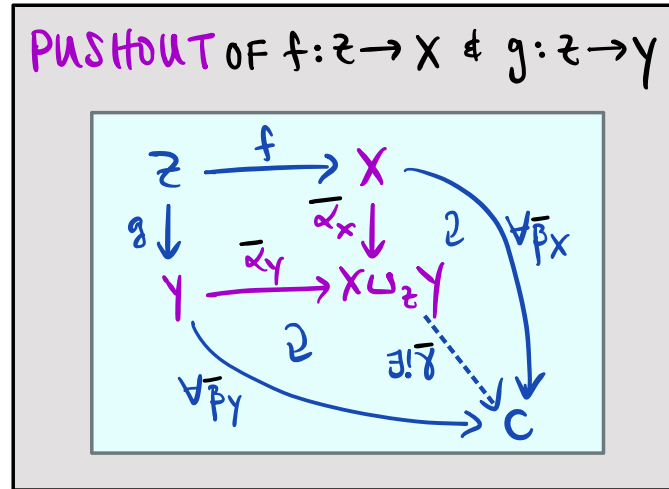


≡ RECALL ≡

GIVEN A CATEGORY  $\mathcal{C}$ :

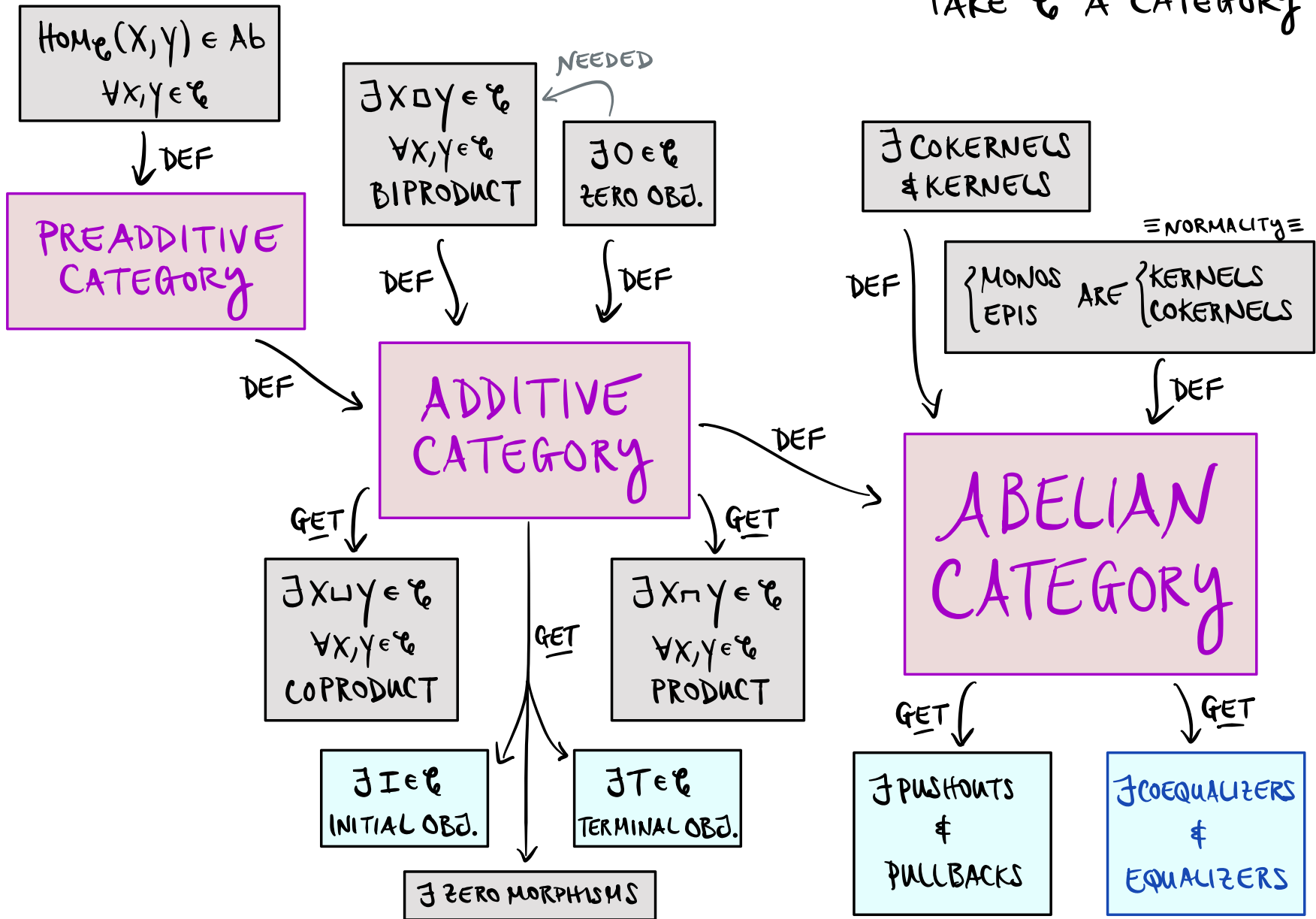


HAVE THE FOLLOWING  
EXISTENCE IMPLICATIONS



## II. ABELIAN CATEGORIES

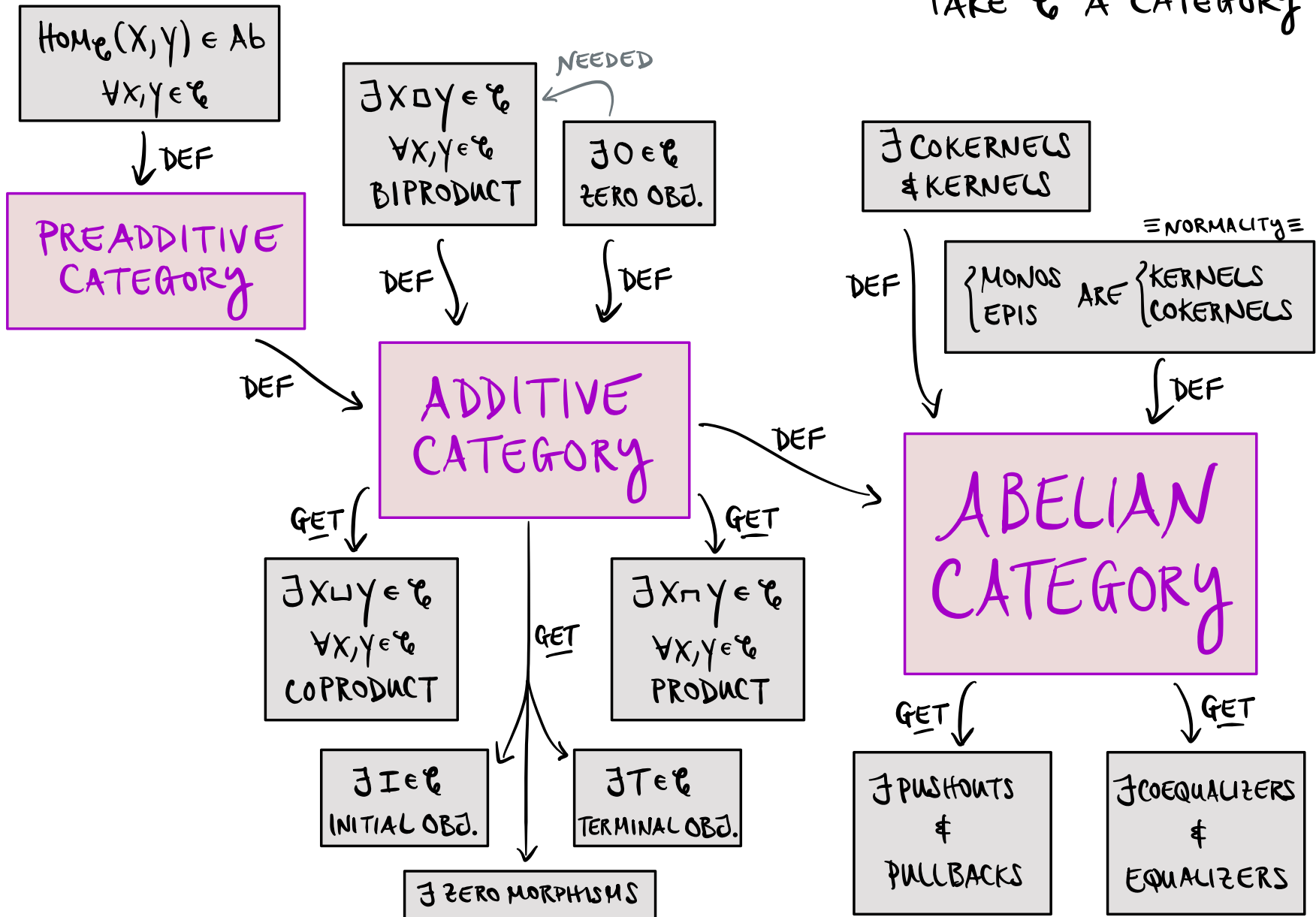
TAKE  $\mathcal{C}$  A CATEGORY





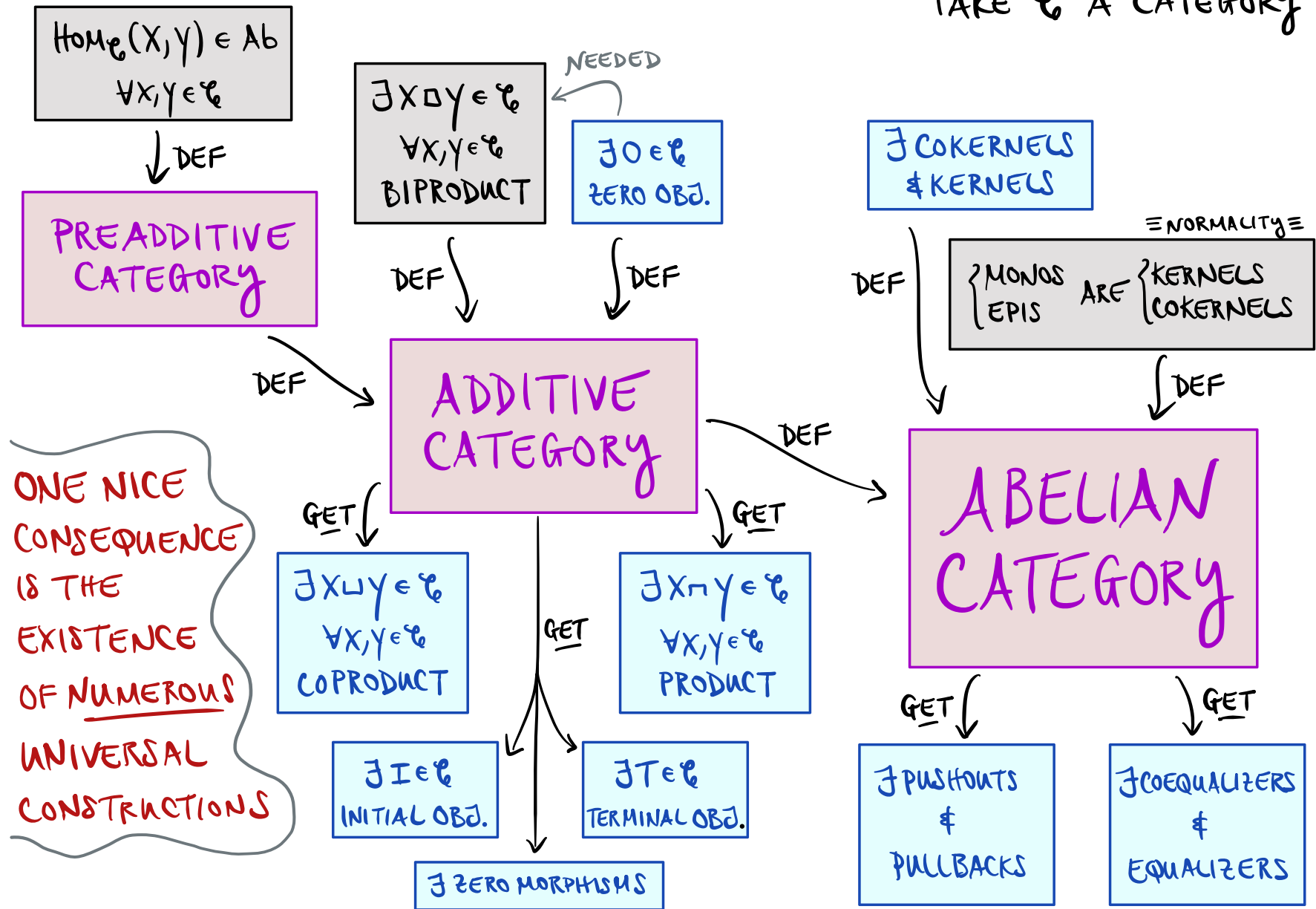
## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY



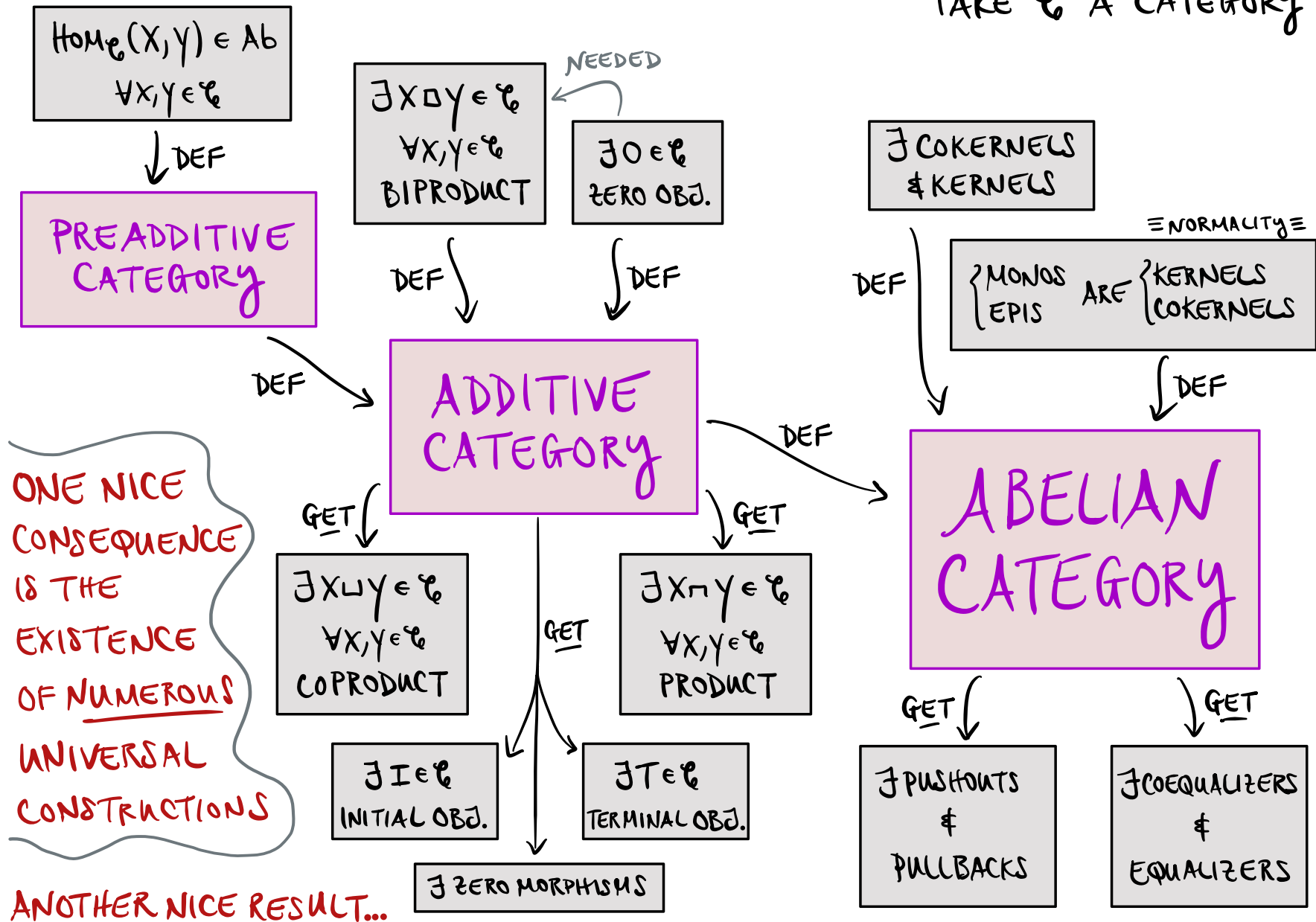
## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY

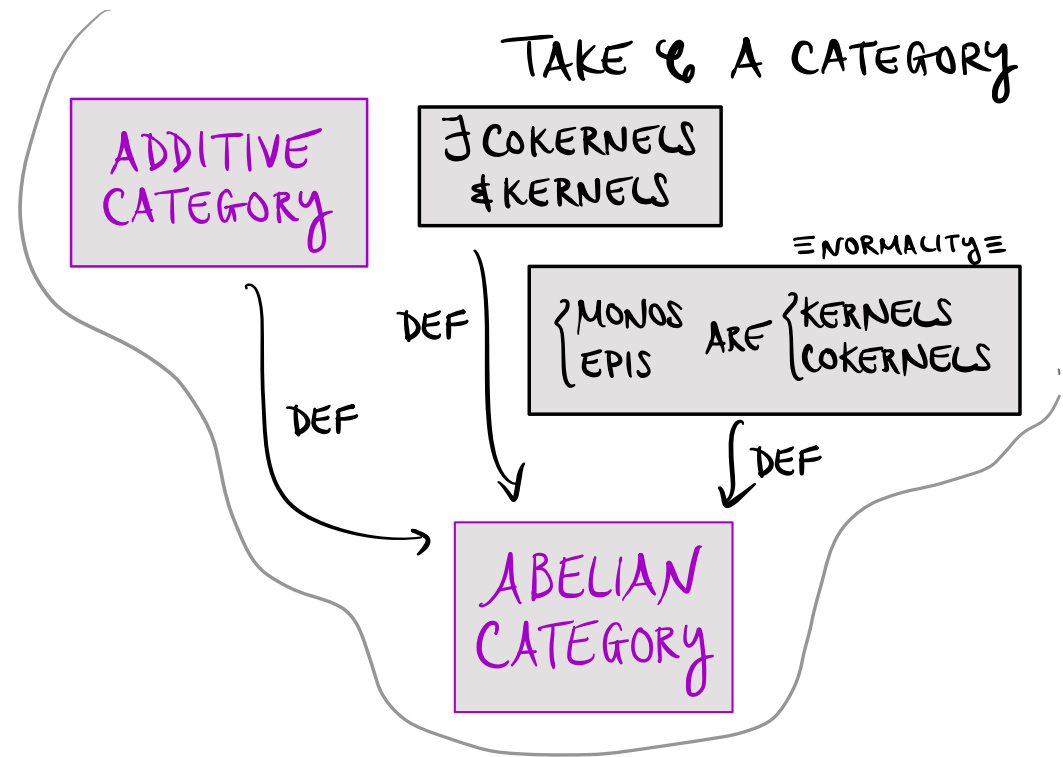


## II. ABELIAN CATEGORIES

TAKE  $\mathcal{C}$  A CATEGORY



## II. ABELIAN CATEGORIES

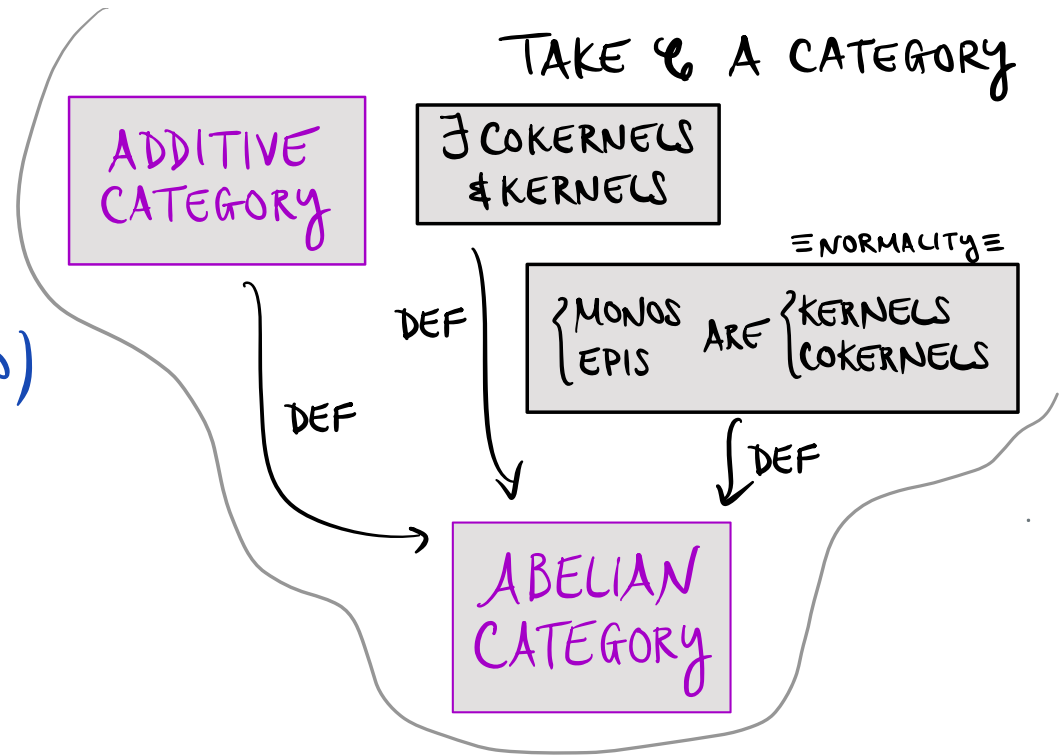


## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$\text{ISO} \Leftrightarrow \text{MONIC EPI} (\Leftrightarrow \text{EPI MONO})$



## II. ABELIAN CATEGORIES

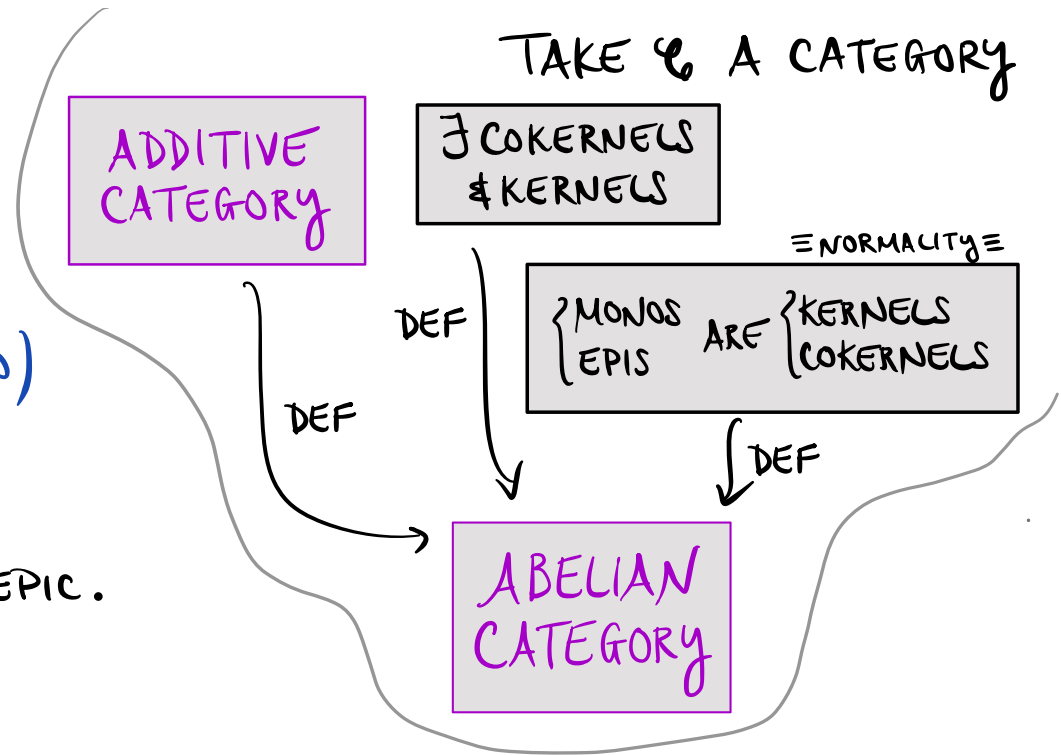
PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$\text{ISO} \Leftrightarrow \text{MONIC EPI} (\Leftrightarrow \text{EPI MONO})$

PF/  $(\Rightarrow)$   $\checkmark$

$(\Leftarrow)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.



## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

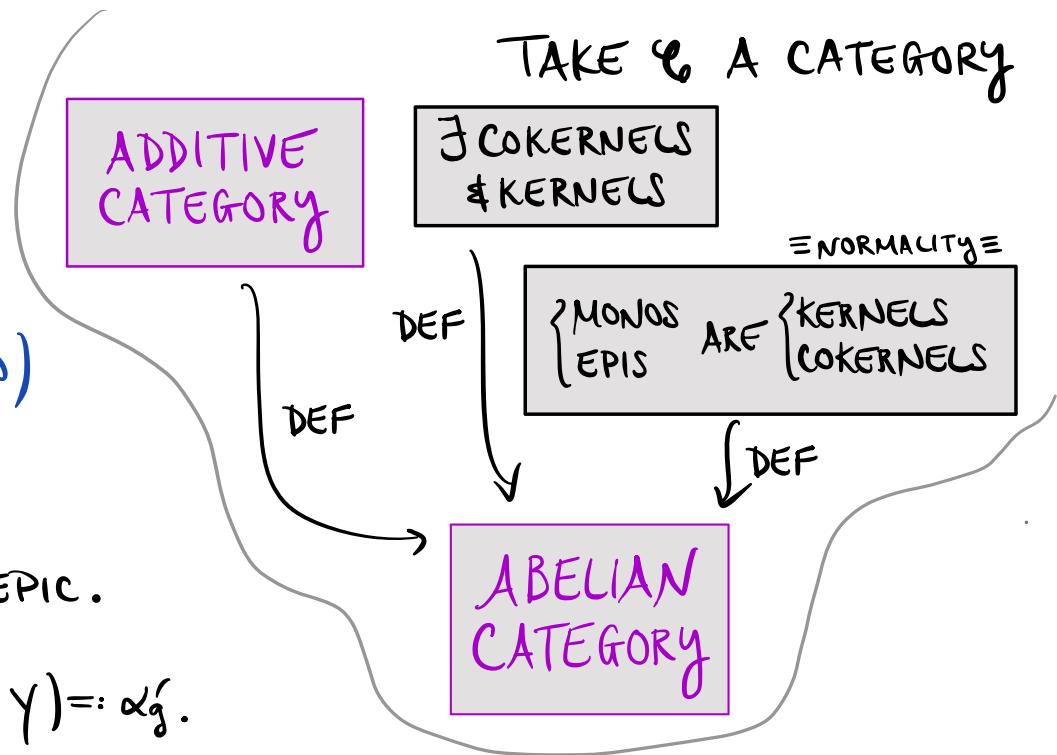
$ISO \iff MONIC \ EPI \ (\iff \ EPI \ MONO)$

PF/  $(\implies)$   $\checkmark$

$(\impliedby)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\implies f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

GET  $\ker(g) \xrightarrow[\substack{\cong \\ \downarrow \delta}]{\alpha'_g = f} X \xrightarrow{g} Y$ .



## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

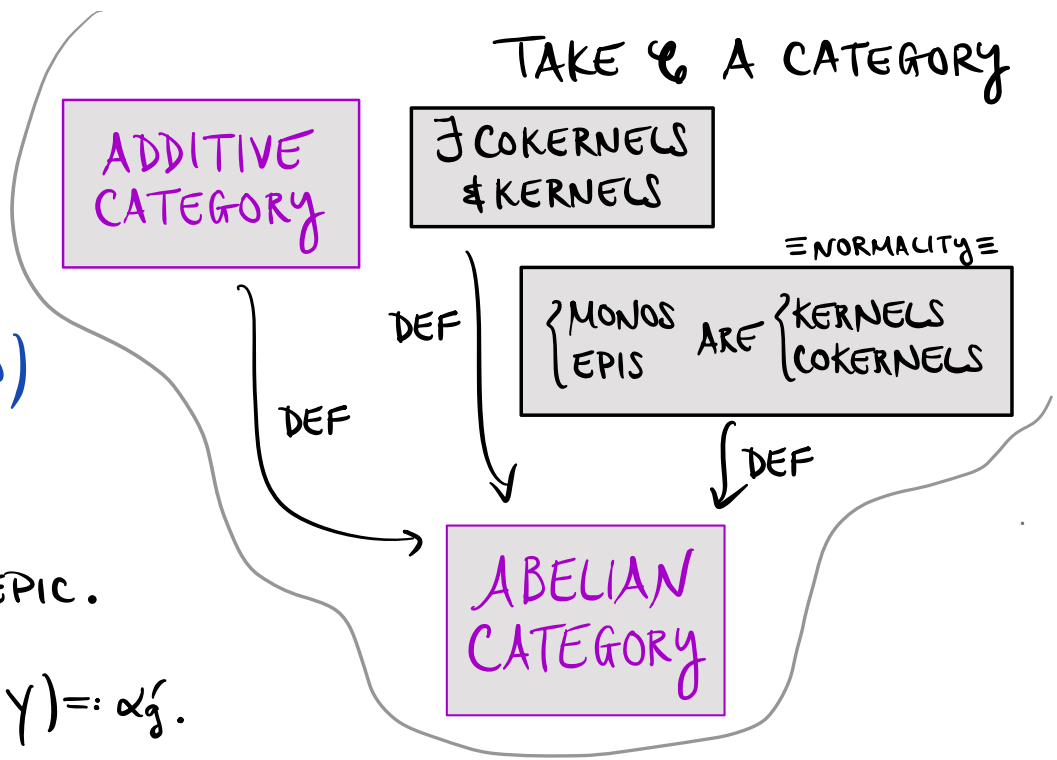
$ISO \iff MONIC \ EPI \ (\iff \ EPI \ MONO)$

PF/  $(\implies)$   $\checkmark$

$(\impliedby)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\implies f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

$$\text{GET } \ker(g) \xrightarrow[\underset{\vec{0}}{\alpha'_g = f}]{\underset{\vec{0}}{f}} X \xrightarrow{g} Y. \quad \text{ALSO } \ker(g) \xrightarrow[\underset{\vec{0}}{f}]{\underset{\vec{0}}{f}} X \xrightarrow{\vec{0}} Y. \quad \therefore gf = \vec{0}f.$$





## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$ISO \Leftrightarrow MONIC \ EPI \ (\Leftrightarrow \ EPI \ MONO)$

PF/  $(\Rightarrow)$  ✓

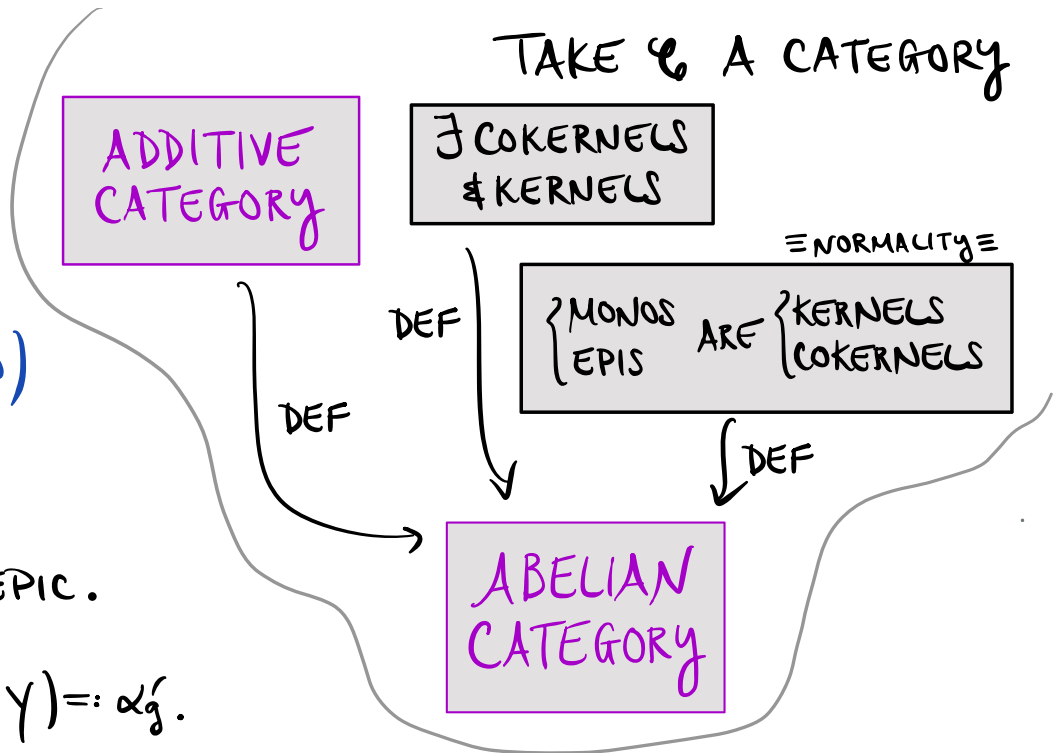
$(\Leftarrow)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\Rightarrow f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

GET  $\ker(g) \xrightarrow[\substack{\alpha'_g=f \\ \downarrow \bar{0}}]{\substack{\downarrow \bar{0} \\ \alpha'_g=f}} X \xrightarrow{g} Y$ . ALSO  $\ker(g) \xrightarrow[\substack{\downarrow \bar{0} \\ f}]{\substack{\downarrow \bar{0} \\ f}} X \xrightarrow{\bar{0}} Y$ .  $\therefore gf = \bar{0}f$ .

NOW  $f$  EPIC  $\Rightarrow g = \bar{0}$ .  $\therefore$  WE GET:

$$\ker(\bar{0}) \xrightarrow{f} X \xrightarrow[\substack{\downarrow \bar{0} \\ \bar{0}=g}]{\bar{0}=g} Y$$



## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$ISO \Leftrightarrow \text{MONIC EPI} (\Leftrightarrow \text{EPI MONO})$

PF/  $(\Rightarrow)$   $\checkmark$

$(\Leftarrow)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\Rightarrow f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

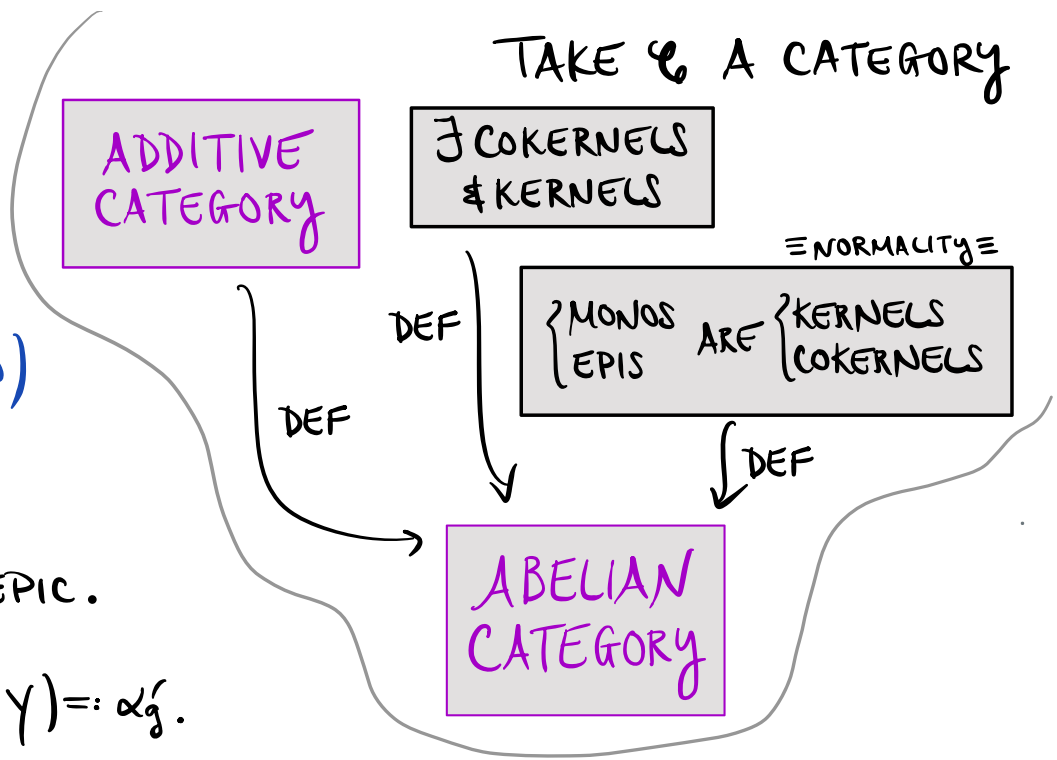
GET  $\ker(g) \xrightarrow[\alpha'_g]{\alpha'_g = f} X \xrightarrow{g} Y$ . ALSO  $\ker(g) \xrightarrow{f} X \xrightarrow{\vec{0}} Y$ .  $\therefore gf = \vec{0}f$ .

NOW  $f$  EPIC  $\Rightarrow g = \vec{0}$ .

$\therefore$  WE GET:

$$\begin{array}{ccc} X & \xrightarrow{\text{id}_X} & X \\ & \searrow & \downarrow \alpha'_g \\ \ker(\vec{0}) & \xrightarrow{f} & X \end{array} \quad \begin{array}{ccc} & \xrightarrow{\vec{0}} & Y \\ & \downarrow \alpha'_g & \\ & \xrightarrow{\vec{0}} & Y \end{array}$$

(PICK  $\beta'_X = \text{id}_X$ )



## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$ISO \iff MONIC \ EPI \ (\iff \ EPI \ MONO)$

PF/  $(\implies)$   $\checkmark$

$(\impliedby)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\implies f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

GET  $\ker(g) \xrightarrow[\substack{\cong \\ \cong}]{\alpha'_g = f} X \xrightarrow{g} Y$ . ALSO  $\ker(g) \xrightarrow[\substack{\cong \\ \cong}]{f} X \xrightarrow{\vec{0}} Y$ .  $\therefore gf = \vec{0}f$ .

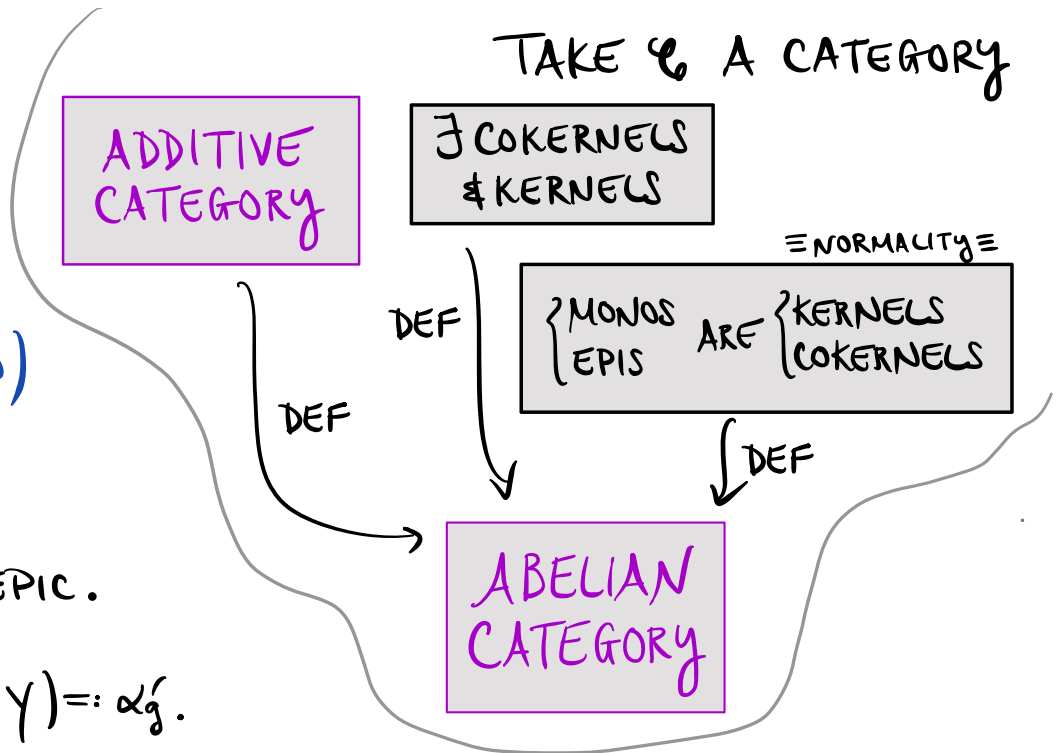
NOW  $f$  EPIC  $\implies g = \vec{0}$ .

$\therefore$  WE GET:

$$\begin{array}{c} X \xrightarrow{\text{id}_X} X \xrightarrow{\vec{0}} Y \\ \exists! \gamma \downarrow \cong \quad \cong \quad \cong \\ \ker(\vec{0}) \xrightarrow{f} X \xrightarrow{\vec{0}} Y \end{array}$$

(PICK  $\beta'_X = \text{id}_X$ )

$\therefore f\gamma = \text{id}_X$



## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$$ISO \Leftrightarrow \text{MONIC EPI} (\Leftrightarrow \text{EPI MONO})$$

PF/  $(\Rightarrow)$  ✓

$(\Leftarrow)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

$$\text{NORMALITY} \Rightarrow f = \ker(g: X \rightarrow Y) =: \alpha'_g.$$

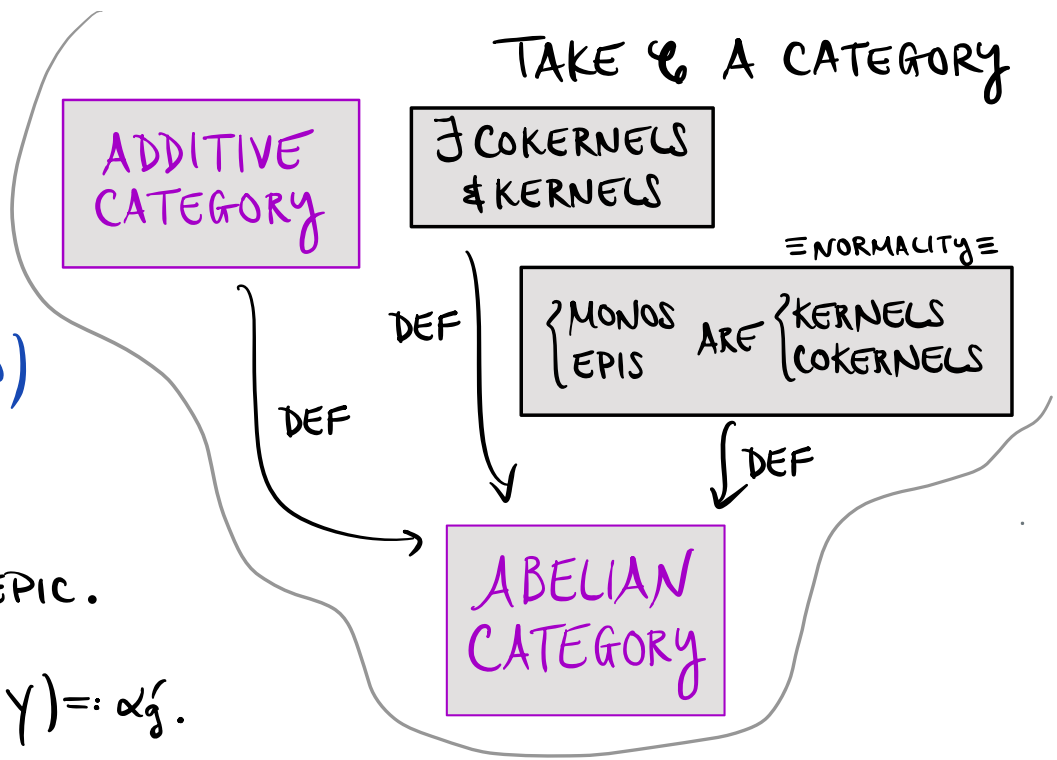
$$\text{GET } \ker(g) \xrightarrow[\substack{\alpha'_g = f \\ \downarrow \cong}}{f} X \xrightarrow{g} Y. \quad \text{ALSO } \ker(g) \xrightarrow[\substack{f \\ \downarrow \cong}}{f} X \xrightarrow[\substack{\cong \\ \downarrow \cong}]{\vec{0}} Y. \quad \therefore gf = \vec{0}f.$$

NOW  $f$  EPIC  $\Rightarrow g = \vec{0}. \quad \therefore$  WE GET:

$$\begin{array}{ccc} X & \xrightarrow{\text{id}_X} & X \\ \downarrow \exists \gamma & \searrow \cong & \downarrow \vec{0} \\ \ker(\vec{0}) & \xrightarrow{f} & X \end{array} \quad \begin{array}{c} \xrightarrow{\cong} \\ \xrightarrow{\vec{0}} \end{array} \begin{array}{c} X \\ \xrightarrow{\vec{0}} \\ Y \end{array} \quad (\text{PICK } \beta'_X = \text{id}_X)$$

$$\therefore f\gamma = \text{id}_X \quad \therefore f\gamma f = f$$

$$\text{NOW } f \text{ MONIC} \Rightarrow \gamma f = \text{id}_{\ker(\vec{0})}$$



## II. ABELIAN CATEGORIES

PROP:

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :

$ISO \Leftrightarrow MONIC \ EPI \ (\Leftrightarrow \ EPI \ MONO)$

PF/  $(\Rightarrow)$  ✓

$(\Leftarrow)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\Rightarrow f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

GET  $\ker(g) \xrightarrow[\substack{\alpha'_g = f \\ \downarrow \cong}]{\cong} X \xrightarrow{g} Y$ . ALSO  $\ker(g) \xrightarrow[\substack{\downarrow \cong}]{f} X \xrightarrow{\cong} Y$ .  $\therefore gf = \vec{0}$ .

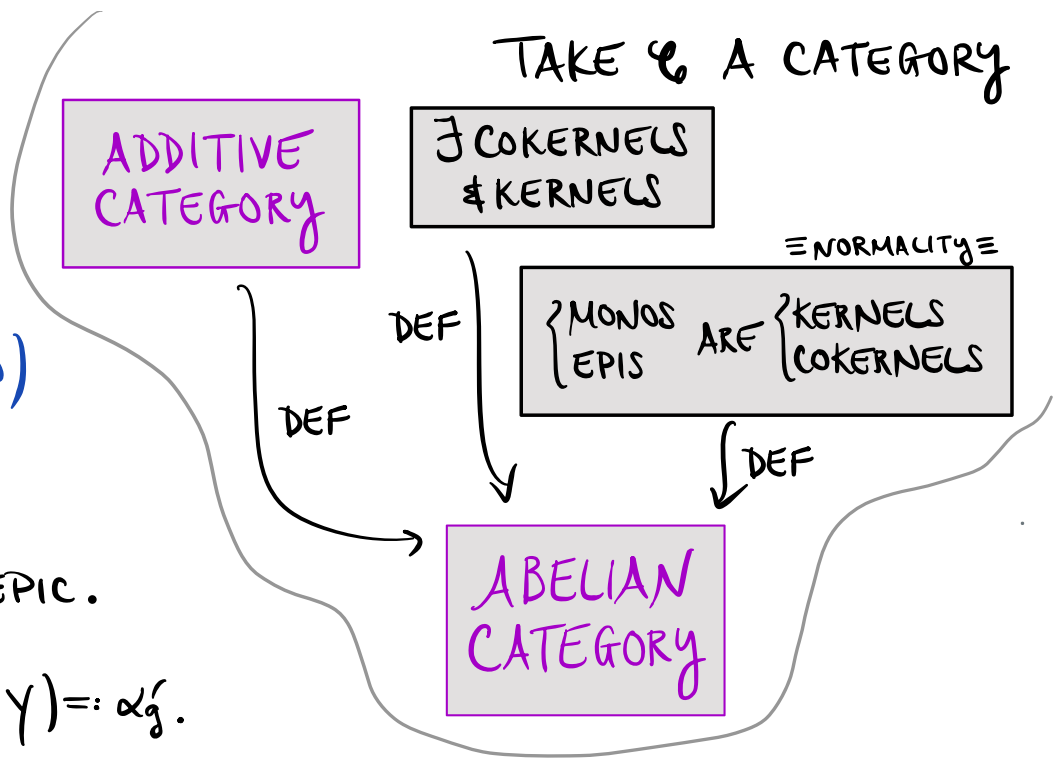
NOW  $f$  EPIC  $\Rightarrow g = \vec{0}$ .  $\therefore$  WE GET:

$$\begin{array}{c}
 X \xrightarrow{\text{id}_X} X \xrightarrow{\vec{0}} Y \\
 \exists! \gamma \downarrow \cong \quad \downarrow \cong \quad \downarrow \cong \\
 \ker(\vec{0}) \xrightarrow{f} X \xrightarrow{\vec{0}} Y
 \end{array}
 \quad
 \begin{array}{l}
 \text{(PICK } \beta'_X = \text{id}_X) \\
 \text{and } \gamma = \text{id}_X
 \end{array}$$

$\therefore f\gamma = \text{id}_X \quad \therefore f\gamma f = f$

NOW  $f$  MONIC  $\Rightarrow \gamma f = \text{id}_{\ker(\vec{0})}$

$\therefore f$  IS AN ISO w/  $f^{-1} = \gamma$  //



## II. ABELIAN CATEGORIES

PROP: Ex. Ring IS NOT ABELIAN (EXER 2.2)

IN AN ABELIAN CATEGORY  $\mathcal{C}$ :  
 $ISO \iff MONIC \ EPI \ (\iff \ EPI \ MONO)$

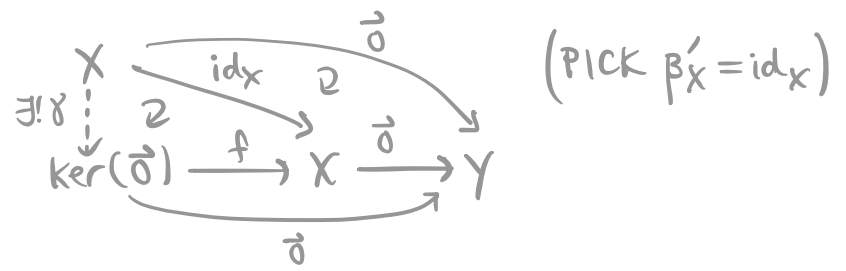
PF/  $(\implies) \checkmark$

$(\impliedby)$  LET  $f: W \rightarrow X$  BE MONIC EPIC.

NORMALITY  $\implies f = \ker(g: X \rightarrow Y) =: \alpha'_g$ .

GET  $\ker(g) \xrightarrow[\substack{\alpha'_g = f \\ \cong}]{\cong} X \xrightarrow{g} Y$ . ALSO  $\ker(g) \xrightarrow[\substack{\cong \\ \cong}]{f} X \xrightarrow{\vec{0}} Y$ .  $\therefore gf = \vec{0}f$ .

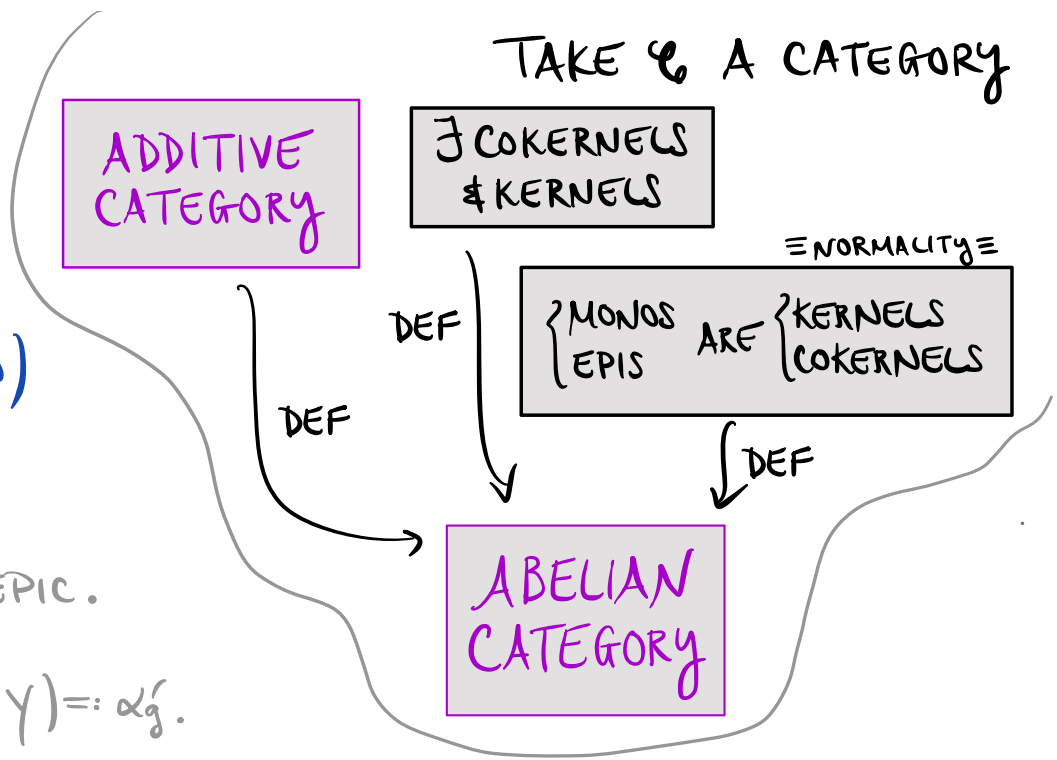
NOW  $f$  EPIC  $\implies g = \vec{0}$ .  $\therefore$  WE GET:



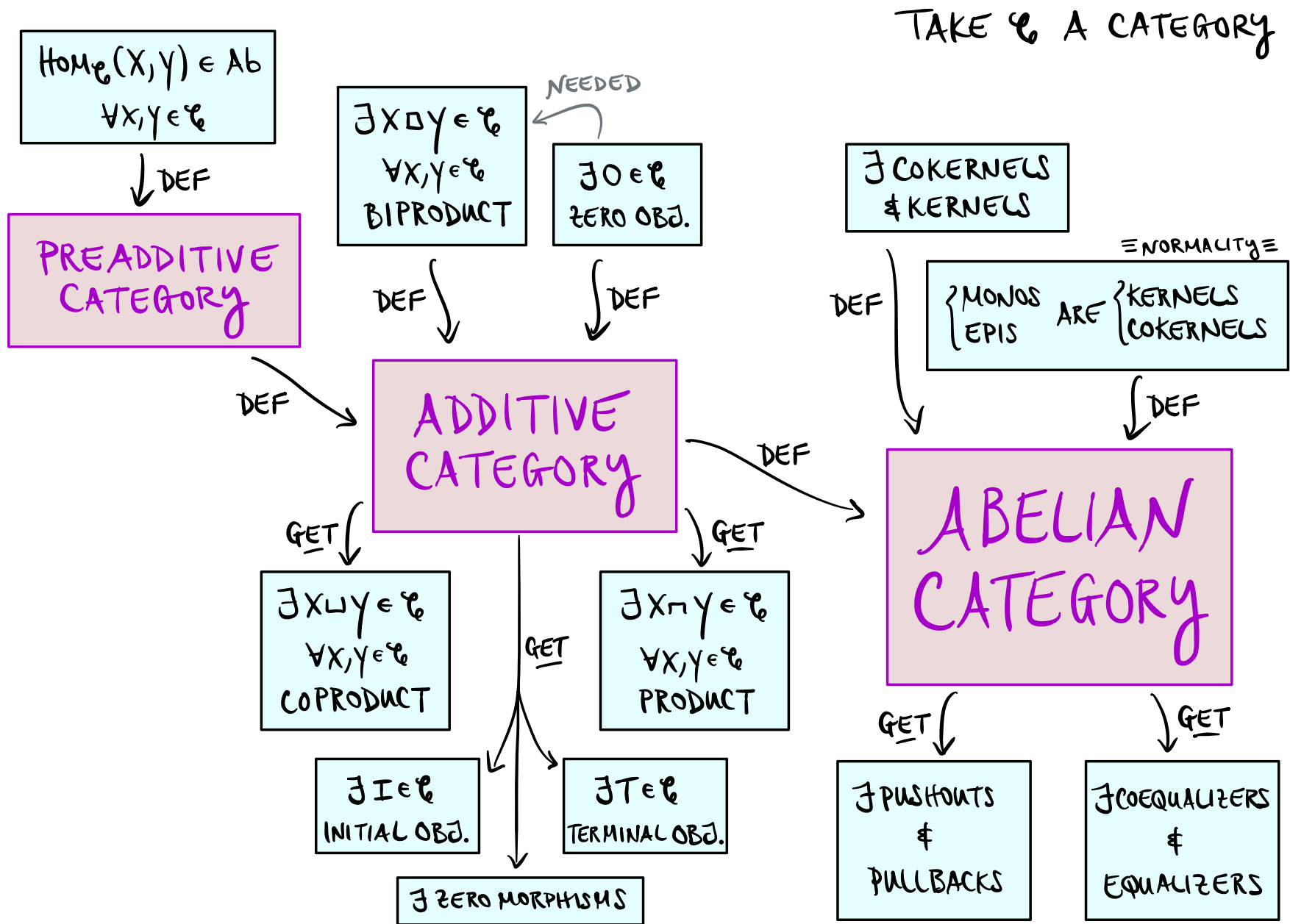
$\therefore f\gamma = id_X \quad \therefore f\gamma f = f$

NOW  $f$  MONIC  $\implies \gamma f = id_{\ker(\vec{0})}$

$\therefore f$  IS AN ISO w/  $f^{-1} = \gamma$  .//



## II. ABELIAN CATEGORIES



MATH 466/566  
SPRING 2024

CHELSEA WALTON  
RICE U.

## LECTURE #7

### TOPICS:

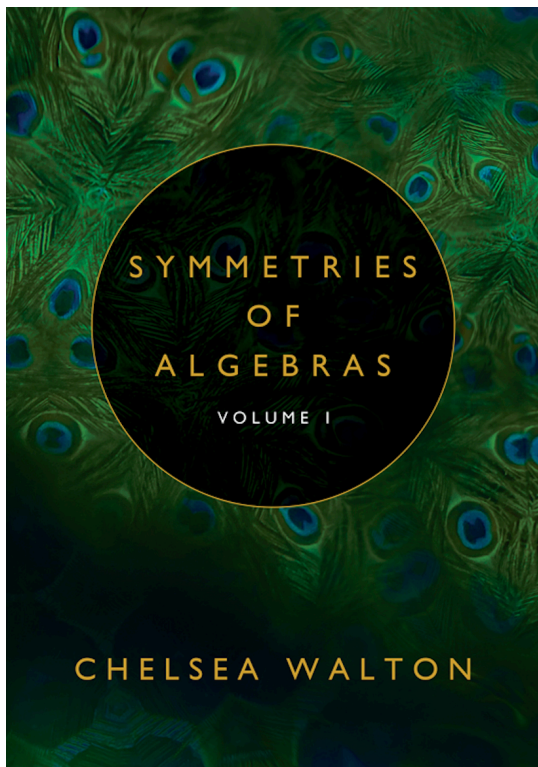
- ✓ I. UNIVERSAL CONSTRUCTIONS (§2.2.1)
- ✓ II. ABELIAN CATEGORIES (F2.2.2)

NEXT: FUNCTORS & NATURAL TRANSFORMATIONS



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Lecture #7 keywords: additive category, abelian category, biproduct of objects, coequalizer of morphisms, cokernel, equalizer of morphisms, kernel, preadditive category, pullback of morphisms, pushout of morphisms